

BALM-SETL - a simple implementation of SETL

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It is suggested that a preliminary implementation of SETL be done in BALM. Though inefficient, such an implementation would be simple, and would permit experimentation with the form of the language, and perhaps be useful as a bootstrap for later versions.

Given below are routines to implement most of the basic SETL operations. A set is implemented as a BALM list whose first item is the value of the variable SETMARK, which will print as \*\*\*. An ordered set is represented by a BALM list. All other data-objects in BALM can be members of sets or ordered sets.

Thus the SETL items are:

integer, real, octal,  
character, string,  
procedure, expression  
set, list, vector

with the following special objects

TRUE, FALSE, UNDEF, SETMARK  
NIL(empty list), NULLSET(empty set)

Undefined will print out as 000.

Operations

The following operations can be written as shown:

$\# a$	SIZ A	SETSIZE(A)
$\} a$	EL A	CHOOSEL(A)
$a \text{ less } \} a$	REM A	CHREM(A)
$x \in a$	X IN A	SETMEMB(X,A)
$a \subseteq b$	A SUBSET B	SUBSET(A,B)
$a \subset b$	A PSUBSET B	PSUBSET(A,B)
$a \text{ with } x$	A WITH X	WITH(A,X)
$a \cup b$	A UNION B	UNION(A,B)
$a \cap b$	A INT B	INTERSECT(A,B)
$a \text{ less } x$	A LESS X	LESS(A,X)
$a \text{ lessf } x$	A LESSF X	LESSF(A,X)
$f(x)$	F OF L	APPL(F,L)
$f\{x\}$	F SOF L	SAPPL(F,L)
$a \equiv b$	A EQS B	EQSET(A,B)
$\langle x,y,z \rangle$		LIST(X,Y,Z)
$\{x,y,z\}$		MAKSET(X;Y,Z)
$(\exists x \in a)p(x)$	P EXISTS A	EXISTS(P,A)
$(\exists x \in \ell)p(x)$	P EXISTL L	EXISTL(P,L)
$(\forall x \in a)p(x)$	P UNIVS A	UNIVS(P,A)
$(\forall x \in \ell)p(x)$	P UNIVL L	UNIVL(P,L)

The middle column shows the prefix or infix form, and the right column the procedural form. We have written A and B for sets; X, Y, and Z for arbitrary SETL items; and L for a list. Precedences have not been chosen, but can be established along with the definitions of the operators, as follows:

UNARY(=SIZ,=SETSIZE,750);

INFIX(=WITH,=WITH,731,730);

The procedure definitions are as follows:

SETMARK = MKVAR (<\*\*\*>);

NULLSET = SETMARK: NIL;

SETSIZE = PROC(S), LENGTH(S)-1 END;

CHOOSEL = PROC(S), HD TL S END;

CHREM = PROC(S), SETMARK: TL TL S END;

EQSET = PROC(S1, S2), IF ATOM(S1)  $\vee$  STRINGP(S1) THEN S1  $\equiv$  S2  
ELSEIF LENGTH(S1)  $\equiv$  LENGTH(S2) THEN FALSE  
ELSEIF VECTORP(S1) THEN  
    (IF VECTORP(S2) THEN EQV(S1, S2)  
    ELSE FALSE)  
ELSEIF HD S1  $\equiv$  SETMARK THEN  
    (IF HD S2  $\equiv$  SETMARK THEN EQS(TL, S1,  
    TL S2) ELSE FALSE)  
ELSE EQL(S1, S2) END;

EQL = PROC(L1, L2), IF NULL(L1) THEN TRUE  
ELSEIF EQSET(HD L1, HD L2) THEN EQL(TL L1,  
TL L2)  
ELSE FALSE END;

EQV = PROC(V1, V2), BEGIN(L, I)  
L = LENGTH(V1), I = 1

NXT = IF I GT L THEN RETURN TRUE;  
IF (EQSET(V1[I], V2[I])) THEN RETURN FALSE,  
I = I+1, GO TO NXT  
END END;

```
EQS      = PROC(S1, S2), IF NULL(S1) THEN TRUE
          ELSEIF SETMEMB(HD S1, S2) THEN EQS(TL S1, S2)
          ELSE FALSE      END;

SETMEMB = PROC(E, S), IF NULL(S) THEN TRUE
          ELSEIF EQSET(E, HD S) THEN TRUE
          ELSE SETMEMB(E, TL S)      END;

WITH     = PROC(S, E), IF SETMEMB(E, S) THEN S ELSE SETMARK:E:TL S
          END;

LESS    = PROC(S, E), IF NULL(S) THEN NIL
          ELSEIF EQSET(HD S, E) THEN TL S
          ELSE HD S:LESS(TL S, E)      END;

LESSF   = PROC(S, E), IF NULL(S) THEN NIL
          ELSEIF EQSET(HD HD S, E) THEN TL S
          ELSE HD S:LESSF(TL S, E)      END;

UNDEF   = MKVAR(<000>);

APPL    = PROC(FN, A), IF NULL(F) THEN UNDEF
          ELSEIF EQL(A, HD FN) THEN LAST(HD FN)
          ELSE APPL(TL FN, A)      END;

SAPPL   = PROC(FN, A), SETMARK:SAPPL1(FN, A)
          END;

SAPPL1  = PROC(FN, A), IF NULL(F) THEN NIL
          ELSEIF EQL(A, HD FN), THEN LAST(HD FN):
          SAPPL1(TL FN, A)
          ELSE SAPPL1(TL FN, A)      END;
```

```

MAKSET      = NPROC(L), SETMARK:SETC(L)  END;
SETC        = PROC(L), SETMARK:REMUDS(L) END;
REMUDS      = PROC(L), IF NULL(L) THEN NIL
              ELSEIF HD L  $\equiv$  UNDEF THEN REMUDS(TL L)
              ELSE WITH(REMUDS(TL L), HD L)  END;
EXISTS      = PROC(A,P), EXIST(TL A,P)  END;
EXISTL      = PROC(L,P), IF NULL(L) THEN FALSE
              ELSEIF P(HL D) THEN TRUE
              ELSE EXISTL(TL L,P)  END;
UNIVS      = PROC(A,P), UNIVL(TL A,P)  END;
UNIVL      = PROC(L,P), IF NULL(L) THEN TRUE
              ELSEIF P(HD L) THEN UNIVL(TL L,P)
              ELSE FALSE  END;
UNION       = PROC(A,B), IF B  $\equiv$  NULLSET THEN A
              ELSE UNION(A WITH EL B, REM B)  END;
INTERSECT   = PROC(A,B), IF A  $\equiv$  NULLSET THEN NULLSET
              ELSEIF EL A IN B THEN INTERSECT(REM A,B)
              WITH EL A
              ELSE INTERSECT(REM A,B)  END;
SUBSET      = PROC(A,B), IF SIZ A GT SIZ B THEN FALSE
              ELSE SUBSET1(A,B)  END;
SUBSET1     = PROC(A,B), IF A  $\equiv$  NULLSET THEN TRUE
              ELSEIF EL A IN B THEN SUBSET1(REM A,B)
              ELSE FALSE  END;
PSUBSET     = PROC(A,B) IF SIZ A GE SIZ B THEN FALSE
              ELSE SUBSET1(A,B)  END;

```

Notes:

SETC converts a list to a set, using REMUDS.

REMUDS removes duplicated SETL items and UNDEFs from a list.

MAKSET uses SETC, so can have UNDEFs and duplicated items.

### Useful expression forms

The expression  $(\forall x \in a)(\text{block})$  can be implemented in BALM-SETL without any additional extensions as:

```
MAPX(TL A, PROC(X), block END)
```

where block is any valid BALM expression or command. This can easily be improved by operator and macro definitions to a form such as:

```
FOREACH X IN A REPEAT block
```

The expression  $\{f(x,y) \mid x \in a, y \in b, g(x,y)\}$  could be implemented as:

```
SETC(MAPX(TL A,  
        PROC(X),  
        MAPX(TL B, PROC(Y), IF G(X,Y) THEN F(X,Y) ELSE UNDEF  
        END)  
        ) )
```

which looks nasty, but is actually a mechanical translation.

This can also be improved by operator and macro definitions to a form such as:

```
CONSET(IF G(X,Y) THEN F(X,Y), X IN A, Y IN B)
```

The general form of this would be:

```
CONSET(expr, x1 in a1, x2 in a2, ...)
```

where expr is any expression which may be a conditional or may be a block of code. For example, we could also define INTERSECT by:

```
INTERSECT = PROC(A,B), CONSET(IF X IN B THEN X,X IN A) END;
```

Example:

Suppose we consider Boolean expressions in disjunctive normal form. Because  $\wedge$  and  $\vee$  are commutative, such expressions can be expressed as sets of sets of variables. Thus the expression  $(A \wedge B \wedge C) \vee (B \wedge C) \vee (C \wedge D)$  can be expressed as  $\{\{A, B, C\}, \{B, C\}, \{C, D\}\}$ , or in BALM-SETL as:

```
S = SETC(SETC(=A,=B,=C), SETC(=B,=C), SETC(=C,=D))
```

Such an expression can be partially simplified by deleting any term which includes the variables of another term. This can be expressed as  $\{x \mid x \in S, \neg(\exists y \in S)(x \supset y)\}$ , which can be written in BALM-SETL as a procedure defined by:

```
SIMPL = PROC(S)
CONSET(IF  $\neg$ (PROC(Y), Y PSUBSET X) EXISTS S THEN X,X IN S)
END;
```

For example, if S is the example above, the command:

```
PRINT(SIMPL(S));
```

will print:

```
(***(**B C) (**C D))
```

the BALM-SETL representation of  $\{\{B, C\}, \{C, D\}\}$ .

Comments:

This implementation is highly inefficient with respect to computation time, essentially because of the time required to determine if an element is a member of a set, which itself may require the invocation of the routine to test two items for

equality. Other schemes for overcoming this disadvantage are being considered. However, note that sets with common subsets are permitted to share memory in the obvious cases.