

```

*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for *****
*** educational and research purposes only, by the faculty, students, *****
*** and staff of Yale University only. *****
***
*****

```

```

BOTTOM-UP GREEDY

```

```

(BUG.ASSIGN-MES)

```

```

This module assigns MES to each VN, using an algorithm inspired by the
top-down greedy register allocation. This algorithm is "bottom-up"
because we draw our dags right (entrance at the top, exit at the bottom).

```

```

The general algorithm is recursive -- choose functional units for the
operands of a VN, then choose a functional unit for the VN based on
those chosen for the operands. A VN is assigned only after all its
operands and constraining VNs are assigned. When choosing functional
units for the operands of a VN, we are "mindful" of the intended
destination of the VN and the possible functional-units the VN can be
done on.

```

```

When recursing up through the DAG, the VNs with greatest depth are done
first. This is similar to the "highest levels first" strategy of
list scheduling.

```

```

(eval-when (compile load)
  (include list-scheduler:declarations) )

```

```

(defun bug.assign-mes ()
  (heights-depths.assign)
  (resource.initialize-schedule 'bug)
  (registers.initialize)

  (bug.assign-likely-mes)

  (loop (for vn in (bug.depth-sort-vns +ls.exit-vns*)) (do
    (bug.vn:assign vn ( ) ) ) )

  (bug.make-defis&useis)
  ( ) )

```

```

*****
***
*** (BUG.ASSIGN-LIKELY-MES)
***
*** Initializes :LIKELY-MES of each VN to a list of feasible MES that
*** could be used for the VN. For DEF and USE VNs, the list is set to
*** the MES of the locations of the DEF or USE (if given); if a DEF has
*** no locations, but a USE that reads the DEF does, the DEF adopts the
*** :LIKELY-MES of the USE. For operations VNs, it is set to the list
*** of functional-unit MES that can compute the value of the VN.
***
*****

```

```

(defun bug.assign-likely-mes ()
  (loop (for-each-vn vn) (do
    (caseq (vn:type vn)
      (pseudo-op)

```

```

( (def use)
  (:= (vn:likely-mes vn)
      (loop (for (me register) in (vn:locations vn) ) (save
        me) ) ) )

(copy
  (:= (vn:likely-mes vn) ( ) ) )

(operation
  (:= (vn:likely-mes vn)
      (loop (for me in (operator:functional-unit-mes
        (vn:operator vn) ) )
        (when (me:datum:ok? me vn) )
        (save me) ) ) )

(t
  (error (list vn "BUG.ASSIGN-LIKELY-MES: Case error."))))))

(loop (for-each-vn vn)
  (when (&& (== 'def (vn:type vn) )
    (! (vn:likely-mes vn) ) ) )
  (bind use-vn (car (for-some (reading-vn in (vn:reading-vns vn) )
    (== 'use (vn:type reading-vn) ) ) ) )
  (when use-vn)
  (do
    (:= (vn:likely-mes vn) (vn:likely-mes use-vn) ) ) )
  ( ) )

```

```

*****
***
*** (BUG.VN:ASSIGN VN DESTINATIONS)
***
*** This is the recursive function that assigns a VN, setting :ME,
*** :LIKELY-MES, and :BUG-CYCLE of VN to reflect our choice of where to
*** compute VN. DESTINATIONS is a list of MES where we want the value
*** of VN to end up in (the value need end up in only one of them); if
*** DESTINATIONS is ( ) that means we don't care where the value goes.
***
*** After assigning VN, all the operands of VN are examined; any DEF
*** operands are reassigned based on what was chosen for VN.
***
*****

```

```

(defun bug.vn:assign ( vn destinations )
  (if (! (vn:bug-cycle vn) ) (then
    (caseq (vn:type vn)
      (pseudo-op)
      (def
        (bug.def-vn:assign      vn destinations) )
      (use
        (bug.use-vn:assign     vn destinations) )
      (copy
        (bug.copy-vn:assign    vn destinations) )
      (operation
        (bug.operation-vn:assign vn destinations)

        (loop (for operand-vn in (vn:operand-vns vn) )
          (when (== 'def (vn:type operand-vn) ) )
          (do
            (bug.def-vn:reassign operand-vn (vn:me.vn) ) ) ) )

      (t
        (error (list vn "BUG.VN:ASSIGN: Case error." ) ) ) ) ) ) )
  ( ) )

```

```

=====
***
*** (BUG.USE-VN:ASSIGN VN DESTINATIONS)
***
*** Assigns a USE VN merely by recursively assigning the operand of
*** the USE.
***
=====
(defun bug.use-vn:assign ( vn destinations )
  (bug.vn:assign (car (vn:operand-vns vn) )
                 (vn:likely-mes vn) )
  (:= (vn:bug-cycle vn) (vn:bug-cycle (car (vn:operand-vns vn) ) ) )
  ( ) )

=====
***
*** (BUG.COPY-VN:ASSIGN VN DESTINATIONS)
***
*** Assigns a COPY VN by recursively assigning its operand. It is assumed
*** that the copy operation will take one cycle and can be done in
*** whatever register bank contains the operand.
***
=====
(defun bug.copy-vn:assign ( vn destinations )
  (loop (for operand-vn
         in (bug.depth-sort-vns (append (vn:constraining-vns vn)
                                       (vn:operand-vns vn) ) ) )
    (do
      (bug.vn:assign operand-vn
                     (if (memq operand-vn (vn:operand-vns vn) )
                         destinations
                         ( ) ) ) )
      (:= (vn:bug-cycle vn) (+ 1 (bug.vn:available (car (vn:operand-vns vn) ) ) ) )
      (:= (vn:likely-mes vn) (vn:likely-mes (car (vn:operand-vns vn) ) ) )
      ( ) )

=====
***
*** (BUG.DEF-VN:ASSIGN VN DESTINATIONS)
***
*** VN is a DEF. We assume that the machine has the bandwidth to provide
*** VN to its readers on cycle 0 (see the next function).
***
=====
(defun bug.def-vn:assign ( vn destinations )
  (:= (vn:bug-cycle vn) 0) )

=====
***
*** (BUG.OPERATION-VN:ASSIGN VN DESTINATIONS)
***
*** VN is an operation VN. Picks a good ME for for VN. First all the
*** predecessor VNs (operands and constraining) are recursively assigned,
*** doing the ones of largest depth first. As we assign each operand,

```

```

=====
*** we make a "best guess" about the MEs likely to be used for VN; that
*** best guess is based on as DESTINATIONS for the recursive assignment
*** to the operand. After we've done all the operands, we then make
*** a final best guess for VN, preferring an ME that is closest to the
*** operands if we have any choice left.
***
*** After we've picked an ME and a cycle, we schedule it on the resource
*** schedule to indicate we've made a choice.
***
=====
(defun bug.operation-vn:assign ( vn destinations )
  (loop (for operand-vn
         in (bug.depth-sort-vns (append (vn:constraining-vns vn)
                                       (vn:operand-vns vn) ) ) )
    (bind cycle&me-list
          (bug.vn:destinations:cycle&me-list
           vn destinations)
      operand-destinations
      (if (not (memq operand-vn (vn:operand-vns vn) ) ) (then
        ( ) )
        (else
         (loop (for (cycle me) in cycle&me-list) (save
           me ) ) ) )
      (do
        (bug.vn:assign operand-vn operand-destinations) ) )
    (let ( (cycle&me-list (bug.vn:destinations:cycle&me-list vn destinations))
          (best-cycle ( ) )
          (best-me ( ) ) )
      (if (not (cdr cycle&me-list) ) (then
        (desetq ( (best-cycle best-me) . ( ) ) cycle&me-list) )
        (else
         (desetq (best-cycle best-me)
                 (loop (for cycle&me in cycle&me-list)
                       (bind (cycle me) cycle&me)
                       (minimize
                        cycle&me
                        (loop (for operand-vn in (vn:operand-vns vn) )
                            (reduce max 0
                                    (me:me:delay (vn:me operand-vn) me)))))))
         (bug.vn:cycle:me:schedule vn best-cycle best-me) )
      ( ) )

=====
***
*** (BUG.VN:DESTINATIONS:CYCLE&ME-LIST VN DESTINATIONS)
***
*** Given an operation VN and suggested destinations, returns a list
*** of pairs:
***
*** (CYCLE ME)
***
*** describing the best places and times where VN can be computed; any
*** pair in the list will calculate VN and move it to a closest
*** destination as early as any other pair. CYCLE specifies the time
*** at which ME can be scheduled.
***
*** The result is calculated by looking at each ME that could be used
*** for VN and calculating the cost of moving the operands to the ME,
*** doing the operation, and then moving the value to the closest

```

```

*** destination.
***
=====
(defun bug.vn:destinations:cycle&me-list ( vn destinations )
  (loop (for me in (vn:likely-mes vn)
        (when (me:datum:ok? me vn)
          (bind cycle (bug.vn:me:cycle vn me)
                cost (+ cycle
                      (+ (me:delay me)
                        (bug.me:me-list:min-delay me destinations) ) ) )
          (save-miniums '(,cycle ,me) cost) ) ) )
  )
=====
***
*** (BUG.VN:ME:CYCLE VN ME)
***
*** Returns the earliest cycle in which ME can be scheduled to compute
*** VN. If none of VN's operands have been assigned yet, then the "cycle"
*** returned is actually an estimate of the cost of moving VN's
*** operands to the ME, based on what little we know of where VN and
*** its operands can be computed.
***
=====
(defun bug.vn:me:cycle ( vn me )
  (if (for-some (operand-vn in (vn:operand-vns vn)
              (vn:bug-cycle operand-vn) )
      (then
        (bug.assigned-vn:me:cycle vn me) )
      (else
        (bug.unassigned-vn:me:cost vn me) ) ) )
  )
=====
***
*** (BUG.UNASSIGNED-VN:ME:COST VN ME)
***
*** VN has none of its operands assigned. Returns an estimate of the
*** cost of moving the operands of VN to the ME. The cost is computed
*** as follows:
***
*** The operands of VN with maximum depth are gathered. The cost
*** is the maximum over those operands of F( OPERAND-VN ), where
*** F is the delay of moving the operand to the ME.
***
=====
(defun bug.unassigned-vn:me:cost ( vn me )
  (let ( (max-depth (loop (for operand-vn in (vn:operand-vns vn)
                          (reduce max 0 (vn:depth operand-vn) ) ) ) )
        (loop (for operand-vn in (vn:operand-vns vn)
              (when (== max-depth (vn:depth operand-vn) )
                (reduce max 0 (bug.me-list:me:min-delay
                              (vn:likely-mes operand-vn)
                              me) ) ) ) )
  )
  )
=====
***
*** (BUG.ASSIGNED-VN:ME:CYCLE VN ME)
***

```

```

*** At least one of VN's operands have been assigned. Returns the first
*** cycle in which VN's operation can be scheduled using ME; the cycle
*** is the first after all the operands are available at the inputs of
*** the ME and in which the ME is free. (We also require that the cycle
*** be after all the constraining VNs have completed.)
***
=====
(defun bug.assigned-vn:me:cycle ( vn me )
  (let ( (available-cycle
        (max
          (loop (for operand-vn in (vn:operand-vns vn)
                (when (vn:bug-cycle operand-vn)
                  (reduce max 0 (+ (bug.vn:available operand-vn)
                                (bug.me-list:me:min-delay
                                  (vn:likely-mes operand-vn)
                                  me) ) ) )
          (loop (for constraining-vn in (vn:constraining-vns vn)
                (for constraining-delay
                  in (vn:constraining-delays vn) )
                (when (vn:bug-cycle constraining-vn)
                  (reduce max
                    0
                    (cycle:delays:constrained-cycle
                     (vn:bug-cycle constraining-vn)
                     (vn:delay constraining-vn)
                     (me:delay me)
                     constraining-delay) ) ) ) )
        (cycle:resource-request:first-available-cycle
         available-cycle
         (me:resources me) ) ) )
  )
  )
=====
***
*** (BUG.VN:AVAILABLE VN)
***
*** Returns the first cycle in which the value of VN is available for use.
***
=====
(defun bug.vn:available ( vn )
  (caseq (vn:type vn)
    ( (def copy
      (vn:bug-cycle vn) )
    (operation
      (+ (vn:bug-cycle vn)
        (me:delay (vn:me vn) ) ) )
    (t
      (error (list vn "BUG.VN:AVAILAABLE: Case error.") ) ) )
  )
  )
=====
***
*** (BUG.DEPH-SORT-VNS VN-LIST)
*** (BUG.HEIGHT-SORT-VNS VN-LIST)
***
*** Sorts a list of VNs in decreasing depth or height order. We break
*** ties based on the name of the VN -- this is slightly less efficient,
*** but guarantees greater uniformity in results when making slight
*** changes to the input program (useful only for debugging).

```

```

:***
:***=====
(defun bug.depth-sort-vns ( vn-list )
  (sort vn-list
    (f:l ( vn1 vn2 )
      (?( (> (vn:depth vn1) (vn:depth vn2) )
        t)
        ( (== (vn:depth vn1) (vn:depth vn2) )
          (lexorder (vn:name vn1) (vn:name vn2) ) ) )
      ( t
        ( ) ) ) ) ) )

(defun bug.height-sort-vns ( vn-list )
  (sort vn-list
    (f:l ( vn1 vn2 )
      (?( (> (vn:height vn1) (vn:height vn2) )
        t)
        ( (== (vn:height vn1) (vn:height vn2) )
          (lexorder (vn:name vn1) (vn:name vn2) ) ) )
      ( t
        ( ) ) ) ) ) )

:***=====
:***
:*** (BUG.VN:CYCLE:ME:SCHEDULE VN CYCLE ME)
:***
:*** ME has been chosen for VN. Its resources are scheduled at CYCLE,
:*** and VN:BUG-CYCLE, VN:ME, and VN:LIKELY-MES are updated. If ME is a
:*** register bank, we assumed it we're trying to read from it.
:***
:***=====
(defun bug.vn:cycle:me:schedule ( vn cycle me )
  (assert me "BUG: No ME was found!")

  (cycle:resource-request:schedule
    cycle
    (if (== 'register-bank (me:type me) )
      (me:read-resources me)
      (me:resources me) ) )

  (:= (vn:me vn) me)
  (:= (vn:likely-mes vn) '(,me) )
  (:= (vn:bug-cycle vn) cycle)
  ( ) )

:***=====
:***
:*** (BUG.ME-LIST:ME:MIN-DELAY ME)
:*** (BUG.ME:ME-LIST:MIN-DELAY ME)
:***
:*** These two functions return the minimum required delay in moving a
:*** a value between one ME and the closest ME in a list of MEs.
:***
:***=====
(defun bug.me-list:me:min-delay ( me-list me )
  (if (! me-list) (then
    0)

```

```

  (else
    (loop (for next-me in me-list)
      (reduce min 100000 (me:me:delay next-me me) ) ) ) ) )

(defun bug.me:me-list:min-delay ( me me-list )
  (if (! me-list) (then
    0)
  (else
    (loop (for next-me in me-list)
      (reduce min 100000 (me:me:delay me next-me) ) ) ) ) )

```

```

*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for *****
*** educational and research purposes only, by the faculty, students, *****
*** and staff of Yale University only. *****
***
*****
=====
: BOTTOM UP GREEDY, DEFS AND USES
:
: This file implements the part of the bottom-up-greedy algorithms that
: handles DEFS and USES.
:
=====
(eval-when (compile load)
  (include list-scheduler:declarations) )

*****
***
*** (BUG.DEF-VN:REASSIGN VN DESTINATION)
***
*** VN is a DEF VN. One of the readers of VN has just been assigned,
*** and this function is called to get a good guess for the location
*** of the DEF. DESTINATION is the ME just assigned to the reader.
***
*** If the DEF requires a single location and doesn't yet have one, then
*** we find a register bank connected to DESTINATION that is available
*** earliest and assign it to VN. Otherwise, the DEF is allowed multiple
*** locations and we don't assign it any location, on the assumption
*** that we'll be able to make up "good" locations later.
***
*** A DEF requires a single location if it is a variable, it doesn't
*** have any currently assigned location, and the variable is written
*** somewhere on the trace. Variables that aren't written on the trace
*** are allowed to have multiple locations, as are constants.
***
*** *** NOTE ***
*** We should really take account of conflicts between the register
*** usage here, as we do later on when we assign DEFS and USES.
***
=====
(defun bug.def-vn:reassign ( vn destination )
  (if (&& (! (vn:constant? vn) )
      (! (vn:likely-mes vn) )
      (> (dag.name:trace-write-count (vn:name vn) ) 0) )
    (then
      (assert *ls.new-defs-allowed?* "BUG: Huh?")
      (let ( (best-cycle 100000)
            (best-me ( ) ) )
          (loop (for me in (me:inputs destination) )
                (when (&& (== 'register-bank (me:type me) ) )
                    (> (me:registers-left me) 0) )
                (bind cycle (cycle:resource-request:first-available-cycle
                              0
                              (me:read-resources me) ) )
                    (when (< cycle best-cycle) )
                  (do
                    (:= best-cycle cycle)
                    (:= best-me me) ) ) ) )
            (loop (for me in (me:inputs destination) )
                  (when (&& (== 'register-bank (me:type me) ) )
                      (> (me:registers-left me) 0) )
                  (bind cycle (cycle:resource-request:first-available-cycle
                                0
                                (me:read-resources me) ) )
                      (when (< cycle best-cycle) )
                    (do
                      (:= best-cycle cycle)
                      (:= best-me me) ) ) ) ) )
  )
)

```

```

(bug.vn:cycle:me:schedule
  vn best-cycle best-me)
(:= (me:registers-left best-me) (- &&& 1) ) ) )
) )

=====
***
*** (BUG.MAKE-DEFIS&USEIS)
***
*** This procedure is called after all main recursive assignment algorithms
*** has been run. It inserts DEFis between DEFS and their readers, and
*** USEis between USEs and their operands.
***
*** First, USEis are inserted between USEs that have no given locations
*** and their operands. Then DEFis are inserted between all DEFS and
*** their readers. Finally, USEis are inserted between USEs that do
*** have given locations and their operands.
***
=====
(defun bug.make-defis&useis ( )
  ;*** Make USEis for USEs that have locations.
  ;***
  (loop (for vn in *ls.exit-vns* )
        (when (&& (== 'use (vn:type vn) )
                (vn:locations vn) ) )
        (do
          (bug.closed-use-vn:make-useis vn) ) )
  ;*** Make DEFis for DEFS.
  ;***
  (loop (for vn in (bug.height-sort-vns *ls.entry-vns* ) )
        (when (== 'def (vn:type vn) ) )
        (do
          (bug.def-vn:make-defis vn) ) )
  ;*** Make USEis for USEs that don't have locations.
  ;***
  (loop (for vn in *ls.exit-vns* )
        (when (&& (== 'use (vn:type vn) )
                (! (vn:locations vn) ) ) )
        (do
          (bug.open-use-vn:make-useis vn) ) )
  ) )

=====
***
*** (BUG.CLOSED-USE-VN:MAKE-USEIS VN)
***
*** VN is a USE VN that has given locations. A USEi for each different
*** location is inserted between VN and its single operand.
***
=====
(defun bug.closed-use-vn:make-useis ( vn )
  (let ( (usei-vns
        (loop (for (me register) in (vn:locations vn) ) (save

```

```

(vn:create (vn:new
  type      'use1
  name      (vn:name      vn)
  constant? (vn:constant? vn)
  datatype  (vn:datatype  vn)
  bug-cycle (vn:bug-cycle vn)
  height    (vn:height    vn)
  depth     (vn:depth     vn)
  likely-mes (.me)
  me        me
  register-bank-me me
  register  (register) ) ) ) )

```

```

(vn:insert-vns vn use1-vns)
( ) )

```

```

=====
***
*** (BUG.OPEN-USE-VN:MAKE-USE1S VN)
***
*** VN is a USE VN that has no given locations.  If the USE doesn't have
*** a DEF1 as its operand, that means the variable of the USE is written
*** on the trace; we splice a USE1 between the operand and the USE.
***
*** If the USE does have a DEF1 as its operand, that means the variable
*** (or constant) isn't written on the trace.  We splice a USE1 between
*** ALL the DEF1s of the variable.  The effect of this is to keep all
*** the copies of this variable live over the whole trace (good if the
*** trace is a loop -- the variable or constant is a loop invariant).
*** Perhaps we'll want to modify this to do it only when given a loop
*** trace.
***
=====

```

```

(defun bug.open-use-vn:make-use1s ( vn )
  (let* ( (operand-vn (car (vn:operand-vns vn) ) )
         (pred-vns
          (if (== 'def1 (vn:type operand-vn) ) (then
            (vn:reading-vns (car (vn:operand-vns operand-vn) ) ) )
            (else
             (vn:operand-vns vn) ) ) ) )
        (loop (for pred-vn in pred-vns)
              (bind use1-vn (vn:create (vn:new
                type      'use1
                name      (vn:name      vn)
                constant? (vn:constant? vn)
                datatype  (vn:datatype  vn)
                bug-cycle (vn:bug-cycle vn)
                height    (vn:height    vn)
                depth     (vn:depth     vn) ) ) ) )
          (do (if (memq pred-vn (vn:operand-vns vn) ) (then
                (vn:splice-vn pred-vn use1-vn (list vn) ) )
              (else
               (push (vn:reading-vns pred-vn) use1-vn)
               (push (vn:operand-vns use1-vn) pred-vn)
               (push (vn:reading-vns use1-vn) vn)
               (push (vn:operand-vns vn) use1-vn) ) ) ) )
          ( ) ) )

```

```

=====
***
*** (BUG.DEF-VN:MAKE-DEF1S VN)
***
*** VN is a DEF.  DEF1s are inserted between the DEF and all its readers.
*** picking likely locations (register banks or constant generators)
*** and assigning them to the DEF1s.
***
*** Each reader of the DEF is considered in turn, and a good location
*** for that reader is chosen.  Then all the readers with the same
*** location are grouped together, and a single DEF1 for that location
*** is spliced between the DEF and the group of readers.
***
=====

```

```

(defun bug.def-vn:make-def1s ( vn )
  (let ( (me-table ( ) )
        (all-reading-vns
         (append (bug.height-sort-vns (vn:reading-vns vn) )
                 (vn:off-live-reading-vns vn) ) ) )
        (number-reading-vns
         (length (vn:reading-vns vn) ) ) )
    ;*** Pick good locations for the DEF by picking a good location
    ;*** that each of the readers would prefer.  Keep a table
    ;*** that groups all the readers of a preferred ME together.
    ;*** As locations are picked, keep track of the number of
    ;*** registers left in register banks.  We have to treat
    ;*** off-live readers almost like normal readers, because
    ;*** they need locations for the DEF as well.
    ;***
    (loop (for reading-vn in all-reading-vns)
          (incr 1 from 1)
          (bind reading-vn? (<= 1 number-reading-vns)
                me (bug.def-vn:reading-vn:choose-me
                  vn
                  (if reading-vn? reading-vn ( ) ) ) )
          (do (if (! (memq me (vn:likely-mes vn) ) ) (then
                (push (vn:likely-mes vn) me)
                (caseq (me:type me)
                  (register-bank
                   (:= (me:registers-left me) (- &&& 1) ) )
                  (constant-generator
                   (cycle:resource-request:schedule
                    (cycle:resource-request:first-available-cycle
                     0 (me:resources me) )
                    (me:resources me) ) )
                  (t
                   (error (list me "Case error.") ) ) ) ) )
              (let ( (me-entry (assoc me me-table) ) )
                (if (! me-entry) (then
                  (:= me-entry '( .me ( ) ( ) ) )
                  (push me-table me-entry) ) )
                (if reading-vn? (then
                  (push (cadr me-entry) reading-vn) )
                (else
                 (push (caddr me-entry) reading-vn) ) ) ) ) )
              ;*** Now for each assigned location of the DEF, we have
              ;*** a group of readers.  Insert a DEF1 between the DEF and

```

```

*** readers.
***
(loop (for (me reading-vns off-live-reading-vns) in me-table)
      (bind (register) (assoc me (vn:locations vn) ) )
      (bind defn-vn (vn:create (vn:new
                               type      'defn
                               name      (vn:name      vn)
                               constant? (vn:constant? vn)
                               datatype   (vn:datatype  vn)
                               bug-cycle  (vn:bug-cycle  vn)
                               height     (vn:height    vn)
                               depth      (vn:depth     vn)
                               likely-mes '(,me)
                               me         me) ) )
      (do
        (if register (then
                     (:= (vn:register-bank-me defn-vn) me)
                     (:= (vn:register      defn-vn) register) ) )
          (vn:splice-vn vn defn-vn reading-vns off-live-reading-vns) ) )
      ) )

=====
***
*** (BUG.DEF-VN:READING-VN:CHOOSE-ME VN READING-VN)
***
*** Chooses a good location for VN, a DEF, to be read by one its children.
*** READING-VN. The choice is made depending on whether or not the DEF
*** is a constant or variable, and whether or not we are allowed to make
*** up new locations for the DEF or must stick with the locations given.
***
=====
(defun bug.def-vn:reading-vn:choose-me ( vn reading-vn )
  (if (vn:constant? vn) (then
    (if *ls.new-defs-allowed?* (then
      (bug.def-open-constant-vn:reading-vn:choose-me  vn reading-vn) )
    (else
      (bug.def-closed-constant-vn:reading-vn:choose-me vn reading-vn))))
  (else
    (if (& *ls.new-defs-allowed?*
        (== (dag.name:trace-write-count (vn:name vn) ) 0) )
      (then
        (bug.def-open-var-vn:reading-vn:choose-me  vn reading-vn) )
      (else
        (bug.def-closed-var-vn:reading-vn:choose-me vn reading-vn))))))

=====
***
*** (BUG.DEF-OPEN-CONSTANT-VN:READING-VN:CHOOSE-ME VN READING-VN)
***
*** VN is a constant DEF for which we are allowed to make up new
*** locations.  If the constant is immediate and we're supposed to load
*** immediates from constant generators, then we look for the closest
*** such generator.  Otherwise, we treat the DEF just as if it were a
*** register, assigning it a register location.
***
=====
(defun bug.def-open-constant-vn:reading-vn:choose-me ( vn reading-vn )
  (let ( (mes (constant:constant-generator-mes (vn:name vn) ) ) )
    (if (& mes

```

```

(== 'load *ls.immediate-constant-action*) )
(then
  (car (bug.def-constant-vn:reading-vn:best-cg-me&cycle
        vn reading-vn) ) )
(else
  (bug.def-open-var-vn:reading-vn:choose-me vn reading-vn) ) ) )

=====
***
*** (BUG.DEF-CLOSED-CONSTANT-VN:READING-VN:CHOOSE-ME VN READING-VN)
***
*** VN is a constant DEF for which we are not allowed to make up new locations.
*** We consider two options:
***
*** 1. Using an already assigned location.
*** 2. Using a constant generator.
***
*** We pick the one that is closest to READING-VN, preferring #1 over #2 if
*** there is a choice (i.e. preferring a register bank over a constant
*** generator, hmmm).
***
=====
(defun bug.def-closed-constant-vn:reading-vn:choose-me ( vn reading-vn )
  (let* ( (dest-me (bug.vn:dest-me reading-vn) )
         (best-delay 100000)
         (best-me ( ) ) )
    (loop (for (me register) in (vn:locations vn) )
          (bind delay (me:me:delay me dest-me) )
          (when (|| (== me dest-me)
                  (< delay best-delay) ) )
          (do
            (:= best-delay delay)
            (:= best-me me) ) )
    (let ( ( (best-cg-me best-cg-delay)
            (bug.def-constant-vn:reading-vn:best-cg-me&cycle
              vn reading-vn) ) )
      (if (& best-cg-me
          (< best-cg-delay best-delay) )
        (then
          (:= best-me best-cg-me) ) ) )
    (assert best-me
      (h vn) t (h reading-vn)
      "BUG.DEF-CLOSED-CONSTANT-VN:READING-VN:CHOOSE-ME: Wasn't able "
      "to find a register or constant generator for VN!")
    best-me ) )

(defun bug.def-constant-vn:reading-vn:best-cg-me&cycle
  ( vn reading-vn )
  (let ( (dest-me (bug.vn:dest-me reading-vn) ) )
    (loop (for me in (constant:constant-generator-mes (vn:name vn) ) )
          (bind cycle (+ (cycle:resource-request:first-available-cycle
                        0
                        (me:resources me) )

```

```

      (me:me:delay me dest-me) ) )
    (minimize '(,me ,cycle) cycle) ) )

:***
:***
:*** (BUG.DEF-OPEN-VN:READING-VN:CHOOSE-ME VN READING-VN)
:*** (BUG.DEF-CLOSED-VN:READING-VN:CHOOSE-ME VN READING-VN)
:***
:*** VN is a DEF of a variable. These two functions pick a good register
:*** bank location for it, the former considering any register bank, the
:*** latter considering only register banks already assigned to the VN.
:***
:***
:***
(defun bug.def-open-var-vn:reading-vn:choose-me ( vn reading-vn )
  (bug.def-var-vn:reading-vn:choose-me
   vn
   reading-vn
   *ls.register-bank-mes*) )

(defun bug.def-closed-var-vn:reading-vn:choose-me ( vn reading-vn )
  (bug.def-var-vn:reading-vn:choose-me
   vn
   reading-vn
   (if reading-vn
    (vn:likely-mes vn)
    *ls.register-bank-mes*) ) )
:***
:*** It's possible to have a closed VN (a variable written
:*** on the trace) that isn't assigned any likely ME at
:*** this point. This only occurs if READING-VN is ().
:*** that is, if VN doesn't have any readers, just off-live
:*** readers.
:***
:***
:***
:*** (BUG.DEF-VAR-VN:READING-VN:CHOOSE-ME VN READING-VN FEASIBLE-MES)
:***
:*** VN is a DEF of a variable. Chooses a good register bank location
:*** for it from FEASIBLE-MES, a list of register banks. We first find
:*** the register banks capable of holding VN that are closest to
:*** READING-VN. Then of those, we find the ones that conflict least
:*** with the other operands of READING-VN that already have locations.
:*** Then if there is still a choice left, we prefer a register bank that
:*** is either identical with the destination of READING-VN (if it is
:*** a USE1) or one that is already assigned to VN.
:***
:***
:***
(defun bug.def-var-vn:reading-vn:choose-me ( vn reading-vn feasible-mes )
  (let* ( (dest-me
          (bug.vn:dest-me reading-vn) )

         (close-mes
          (loop (for me in feasible-mes)
                (when (|| (memq me (vn:likely-mes vn) )
                          (> (me:registers-left me) 0) ) )
                (bind delay (me:me:delay me dest-me) )
                (save-minimums me delay) ) ) )

         (least-conflicting-mes

```

```

      (if (! reading-vn) (then
        close-mes)
        (else
         (loop (for me in close-mes) (save-minimums
           me
           (bug.vn:me:register-conflicts reading-vn me) ) ) ) ) )

(best-me
 (if (memq dest-me least-conflicting-mes) (then
  dest-me)
  (else
   (if-let ( (common-mes (intersection least-conflicting-mes
    (vn:likely-mes vn)))
            (car common-mes)
            (car least-conflicting-mes) ) ) ) ) )

(assert best-me
 (h vn) t (h reading-vn) t
 "BUG.DEF-VAR-VN:READING-VN:CHOOSE-ME: Wasn't able to "
 " find a close register bank!")

best-me) )

:***
:***
:*** (BUG.VN:DEST-ME VN)
:***
:*** Returns the destination ME that an operand of VN needs to deliver its
:*** result to. The ME depends on the type of VN:
:***
:*** OPERATION: the functional unit assigned to the VN.
:*** COPY: wherever the max height reader of VN wants its value.
:*** USE1: any register bank assigned to the use.
:***
:*** Returning () means that we don't have any particular preference for a
:*** a destination ME; this should only occur if VN is a USE1 and doesn't have
:*** a required location. VN may be (), in which case the result is
:*** ().
:***
:***
:***
(defun bug.vn:dest-me ( vn )
  (&& vn
   (caseq (vn:type vn)
    (operation
     (vn:me vn) )
    (copy
     (bug.vn:dest-me (loop (for reading-vn in (vn:reading-vns vn) )
                          (maximize reading-vn
                           (vn:height reading-vn))))))

    (use1
     (car (vn:likely-mes vn) ) )
    (use
     (car (vn:likely-mes vn) ) )
    (t
     (error (list vn "BUG.VN:DEST-ME") ) ) ) ) )

:***
:***
:*** (BUG.VN:ME:REGISTER-CONFLICTS VN ME)
:***
:***

```



```

:*** Attempts to determine whether by using ME (a register bank) as a
:*** location for one of the as-yet unassigned operands of VN, that will
:*** conflict with the already assigned operands. In general, an
:*** assignment of register banks to the operands of VN conflicts if it
:*** is not possible to read all the operands from the register banks
:*** in the same cycle (e.g. because there are not enough register ports).
:***
:*** Currently, we assume that the only possible resource that could cause
:*** conflict is the number of available ports (i.e. all ports of of a
:*** bank are connected to the identical set of destinations). So checking
:*** to see if ME conflicts with previously assigned operands is merely
:*** a matter of counting the number of ME ports required by all the
:*** operands.
:***
:*** Eventually, we might also consider conflict of ME:ME:RESOURCES as
:*** well.
:***
:*** Returns 1 if there is conflict, 0 if not (so we can add them up?).
:***
:***=====

```

```

(defun bug.vn:me:register-conflicts ( vn me )
  (let ( (required-ports
          (loop (for operand-vn in (vn:operand-vns vn) )
                (when (&& (!= 'def (vn:type operand-vn) )
                        (== me (car (vn:likely-mes vn) ) ) ) )
                (reduce + 1 1) ) ) )
        (if (> required-ports (me:read-ports me) )
            1
            0) ) )

```

```
*****  
*** Copyright (C) 1988 John R. Ellis. This file may be used for *****  
*** educational and research purposes only, by the faculty, students, *****  
*** and staff of Yale University only. *****  
*****
```

```
(:= *ls.build-module-list* '(  
  list-scheduler:resource  
  list-scheduler:machine-model  
  list-scheduler:shortest-path  
  
  list-scheduler:vn  
  list-scheduler:dag  
  list-scheduler:heights-depths  
  
  list-scheduler:bug  
  list-scheduler:bug-defs-uses  
  
  list-scheduler:scheduler  
  list-scheduler:sch-operation  
  list-scheduler:sch-copy  
  list-scheduler:sch-splits-joins  
  list-scheduler:registers  
  
  list-scheduler:mis-to-eli  
  list-scheduler:simulator  
  list-scheduler:lsex-options  
  list-scheduler:lsex  
  ) )  
(:= *build-module-list* (append *build-module-list* *ls.build-module-list* )
```

```

:***-----*****
:*** Copyright (C) 1988 John R. Ellis. This file may be used for *****
:*** educational and research purposes only, by the faculty, students, *****
:*** and staff of Yale University only. *****
:***-----*****
:*** Build the latest version of the list-scheduling compiler.
:***

(remprop 'loop 'build-information) ;*** Get rid of Yale Loop info.

(:= *build-module-list* () )

(load 'utilities:build) ;*** This must go first.
(load 'interpreter:build)
(load 'trace:build)
(load 'diophantine:build)
(load 'flow-analysis:build)
(load 'experiments:build)
(load 'list-scheduler:build)

(build)

(load 'list-scheduler:eli-model)

```

```

*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for *****
*** educational and research purposes only, by the faculty, students, *****
*** and staff of Yale University only. *****
*** *****
*****

```

## DAG Making

This module constructs a dag of VNs from the trace.

### (DAG.INITIALIZE)

Initializes this module.

### (DAG.MAKE LIVE-BEFORE SOURCE-RECORD-LIST LIVE-AFTER)

Constructs a dag of VNs from this information (parameters identical in meaning to those of GENERATE-CODE).

### \*LS.ENTRY-VNS\*

### \*LS.EXIT-VNS\*

Lists of all VNs with no predecessors and no successors, respectively, of any kind.

### \*LS.NEW-DEFS-ALLOWED?\*

True if on this trace we are allowed to make up new locations for DEFS (that is, if we weren't given any DEF at the beginning of the trace or if this trace is not the beginning trace of a program.)

### \*LS.LOOP-TRACE?\*

True if this trace begins with a LOOP-START and ends with a TRACE-FENCE pseudo-op, indicating the trace is a body of a loop (hack hack).

### (DAG.NAME:TRACE-WRITE-COUNT NAME)

Returns the number of times NAME was written by an operation on the trace.

### (DAG.NAME:USE-VN NAME)

Returns a USE VN of the given NAME if such a VN exists, () otherwise.

### (DAG.SET-COUNTS)

Sets the predecessors-left and readers-left counts of each VN.

### (CYCLE:DELAYS:CONSTRAINED-CYCLE CYCLE PD SD CD)

Returns the cycle that a constrained successor can earliest be scheduled, given the CYCLE of its constraining predecessor, the predecessor's delay PD, the successor's delay SD, and the constraining delay between them CD. The result cycle is defined as:

$$\text{MAX}(\text{CYCLE} + \text{PD} - \text{SD} + \text{CD}, \text{CYCLE} + \text{CD})$$

This guarantees that the first cycle of the successor must be at least CD cycles after the first cycle of the predecessor, and the last cycle of the successor at least CD cycles after the last cycle of the predecessor. These requirements arise from the nature of partial schedules -- a write of a variable must not be allowed to be copied up into a rejoin above a predecessor read of the same variable that is in the rejoin partial schedule (an analogous situation holds for splits).

### (CYCLE:DELAYS:CONSTRAINING-CYCLE CYCLE PD SD CD)

Returns the cycle that a constraining predecessor can latest be

```

: scheduled given the CYCLE of its constrained successor, the delay
: of the predecessor, the delay of the successor, and the constraining
: delay between them.
:
:=====

```

```

(eval-when (compile load)
  (include list-scheduler:declarations) )

```

```

(declare (special

```

```

  *dag.name:vn*           ;*** Hash table mapping names to the most recent
                          ;*** VN containing that name. Uses EQUALT so
                          ;*** that numbers are considered equal only if they
                          ;*** are of the same type.
  *dag.name:write-count* ;*** Hash table mapping names to the number of times
                          ;*** they were written on the trace.
  *dag.name:locations*   ;*** Hash table mapping names to locations given in
                          ;*** a DEF pseudo-op at the beginning of the trace.
  *dag.name:use-vn*      ;*** Hash table mapping names to USE VNs.
  *dag.live-before*      ;*** The live-before variables.
  *dag.live-after-constants* ;*** Constants read on the trace which we want to
                          ;*** treat as live-after.
  *dag.loop-start?*      ;*** True if this trace contains a LOOP-START?
  *dag.use?*             ;*** True if we saw a non-empty USE.
) )

```

```

(defun dag.initialize ()

```

```

  (:= *ls.entry-vns*           ( ) )
  (:= *ls.exit-vns*           ( ) )
  (:= *ls.new-defs-allowed?*   t )
  (:= *ls.loop-trace?*        ( ) )
  (:= *dag.name:vn*           ( ) )
  (:= *dag.name:write-count*   ( ) )
  (:= *dag.name:locations*     ( ) )
  (:= *dag.live-before*       ( ) )
  (:= *dag.live-after-constants* ( ) )
  (:= *dag.loop-start?*       ( ) )
  (:= *dag.use?*              ( ) )
) )

```

```

(defun dag.make ( live-before source-record-list live-after )
  (dag.initialize)
  (start-trace)

```

```

  (:= *dag.name:vn*           (hash-table:create 'equalt ( ) ( ) ) )
  (:= *dag.name:use-vn*       (hash-table:create 'equalt ( ) ( ) ) )
  (:= *dag.name:write-count*  (hash-table:create ( ) ( ) 0 ) )
  (:= *dag.name:locations*    (hash-table:create 'equalt ( ) ( ) ) )
  (:= *dag.live-before*       live-before)

```

```

  ;*** Make VNs for the trace operations.
  ;***

```

```

  (loop (for source-record in source-record-list) (do
    (dag.source-record:create-vn source-record) ) )

```

```

*** Make VNs for the live after variables and constants, including
*** the constants read on the trace if there isn't a non-empty
*** USE at the end.
***
(if (! *dag.use?*) (then
  (:= live-after (unionq &&& *dag.live-after-constants*) ) ) )
(loop (for name in live-after)
  (when (!= '%trace name) )
  (do
    (dag.live-name:create-use name) ) )

*** Set the successor fields of the VNs.
***
(dag.set-successor-vns)

*** Delete dead and useless VNs.
***
(dag.remove-dead-vns)

*** Gather the lists of entry and exit VNs.
***
(loop (for-each-vn vn) (do
  (if (&& (! (vn:operand-vns vn) )
    (! (vn:constraining-vns vn) ) )
    (then
      (push *ls.entry-vns* vn) ) )

  (if (&& (! (vn:reading-vns vn) )
    (! (vn:constrained-vns vn) ) )
    (then
      (push *ls.exit-vns* vn) ) ) ) )
) )

```

```

=====
***
*** (DAG.SOURCE-RECORD:CREATE-VN
*** (OPER TRACE-DIRECTION DATUM OFF-LIVE) )
***
*** Processes one source-record from the trace, creating an OPER, DEF,
*** or, USE VN for that record.
***
=====
(defun dag.source-record:create-vn ( source-record )
  (let* ( ( oper trace-direction datum off-live)
    source-record)
    ( operator (oper:operator oper) )
    ( group (oper:group oper) ) )

  (?
    *** A DEF?
    ***
    ( (= 'def operator)
      (loop (for (name re register) in (oper:part oper 'body) )
        (when (!= '%trace name) )
        (do
          (:= *ls.new-defs-allowed?* ( ) )
          (push ([[] *dag.name:locations* name)
            '(.(name:re re) .register) ) ) ) )

    *** A USE?
  )
)

```

```

***
( (= 'use operator)
  (loop (for use-spec in (oper:part oper 'body) )
    (when (!= '%trace (car use-spec) ) )
    (do
      (:= *dag.use?* t)
      (dag.use-spec:create-use use-spec) ) ) )

*** A DEF-BLOCK?
***
( (= 'def-block operator)
  (:= *ls.new-defs-allowed?* ( ) )
  (dag.source-record:create-pseudo-op source-record) )

*** A LOOP-START?
***
( (= 'loop-start operator)
  (:= *dag.loop-start?* t)
  (predecessors oper trace-direction ( ) ) )

*** A TRACE-FENCE?
***
( (= 'trace-fence operator)
  (:= *ls.loop-trace?* *dag.loop-start?*)
  (predecessors oper trace-direction ( ) ) )

*** An IN-PARAMETER?
***
( (= 'in-parameter group)
  (dag.source-record:create-in-parameter source-record) )

*** An OUT-PARAMETER?
***
( (= 'out-parameter group)
  (dag.source-record:create-out-parameter source-record) )

*** Some random pseudo op?
***
( (oper:property? oper 'pseudo-op)
  (if (memq operator '(def-block dcl esc assert) )
    (then
      (dag.source-record:create-pseudo-op source-record) )
    (else
      (predecessors oper trace-direction ( ) ) ) ) )

*** An "real" normal operation.
***
( t
  (dag.source-record:create-operation source-record) ) )
) )

```

```

=====
***
*** (DAG.SOURCE-RECORD:CREATE-PSEUDO-OP
*** (OPER TRACE-DIRECTION DATUM OFF-LIVE) )
***
*** Processes one pseudo-operation source-record from the trace, making
*** a PSEUDO-OP VN for that record.
***
=====

```

```

(defun dag.source-record:create-pseudo-op
  ( (oper trace-direction datum off-live) )

  (let ( (vn (vn:create (vn:new
                        type      'pseudo-op
                        oper      oper
                        datum     datum) ) ) )

    (predecessors oper trace-direction vn)
    vn ) )

;====
;***
;*** (DAG.SOURCE-RECORD:CREATE-IN-PARAMETER
;*** (OPER TRACE-DIRECTION DATUM OFF-LIVE) )
;***
;*** Processes an xIN operation representing an input parameter by generating
;*** an FPLOAD that reads in the operand.
;***
;====

(defun dag.source-record:create-in-parameter
  ( (oper trace-direction datum off-live) )

  (if (> (name:rank (oper:part oper 'written) ) 0) (then
    ( ) )
    (else
     (let* ( (datatype (oper:dest-datatype oper) )
             (written (oper:part oper 'written) )
             (address-vn (dag:operand:vn '(address .written) ) )
             (load-op (caseq datatype
                       (float 'fpload)
                       (integer 'ipload) ) )

             (vn (vn:create (vn:new
                            type      'operation
                            name      written
                            oper      '(,load-op () () )
                            operand-vns '(,address-vn)
                            datum     datum) ) ) )

       (:= ([]h *dag.name:vn* written) vn)
       vn ) ) ) )

;====
;***
;*** (DAG.SOURCE-RECORD:CREATE-OUT-PARAMETER
;*** (OPER TRACE-DIRECTION DATUM OFF-LIVE) )
;***
;*** Processes an xOUT operation representing an output parameter by
;*** generating an FPSTORE that stores the operand.
;***
;====

(defun dag.source-record:create-out-parameter
  ( (oper trace-direction datum off-live) )

  (if (> (name:rank (oper:part oper 'read1) ) 0) (then
    ( ) )
    (else
     (let* ( (datatype (oper:dest-datatype oper) )
             (read1 (oper:part oper 'read1) )

             (address-vn (dag:operand:vn '(address .read1) ) )
             (store-op (caseq datatype
                       (float 'fpstore)
                       (integer 'ipstore) ) )

             (vn (vn:create (vn:new
                            type      'operation
                            oper      '(,store-op () () )
                            operand-vns '(,address-vn ,read-vn)
                            datum     datum) ) ) )

       (assert read-vn "DAG: Operand doesn't have a VN.")
       vn ) ) ) )

```

5

PS:&lt;C.S.BULLDOG.LIST-SCHEDULER.TEST&gt;DAG.LSP.70

```

(address-vn (dag:operand:vn '(address .read1) ) )
(read-vn (dag:operand:vn read1) )
(store-op (caseq datatype
          (float 'fpstore)
          (integer 'ipstore) ) )

(vn (vn:create (vn:new
               type      'operation
               oper      '(,store-op () () )
               operand-vns '(,address-vn ,read-vn)
               datum     datum) ) ) )

(assert read-vn "DAG: Operand doesn't have a VN.")
vn ) ) )

;====
;***
;*** (DAG.SOURCE-RECORD:CREATE-OPERATION
;*** (OPER TRACE-DIRECTION DATUM OFF-LIVE) )
;***
;*** Processes one operation source-record from the trace, making an
;*** OPERATION VN for that record.
;***
;====

(defun dag.source-record:create-operation
  ( (oper trace-direction datum off-live) )

  (let* ( (vn (vn:create (vn:new
                        type      (if (memq (oper:operator oper)
                                           '(iassign fassign vbase iconstant
                                           fconstant) )
                                     'copy
                                     'operation)
                        datum     datum
                        name      (if (== 'vstore (oper:group oper) )
                                      ( )
                                      (oper:part oper 'written) )
                        oper      (oper:part oper 'written) )

                        (pred-list (predecessors oper trace-direction vn) ) )

         (:= (vn:operand-vns vn) (dag:operand:vn (oper:operand-vns oper) ) )

         (:= (vn:constraining-vns vn) (dag:pred-list:constraining-vns&delays pred-list) )
         (:= (vn:constraining-delays vn) (vn:constraining-delays vn) )

         (:= (vn:off-live-vns vn) (loop (for name in off-live)
                                       (when (!= '%trace name)
                                           (bind off-live-vn (dag:operand:vn name) )
                                           (save off-live-vn) ) ) )

         (if (vn:name vn) (then

```

6

```

( := ( []h *dag.name:vn* (vn:name vn) ) vn
( := ( []h *dag.name:write-count* (vn:name vn)
      (+ 1 &&&) ) )
() ) )

=====
***
*** (DAG.OPER:OPERAND-VNS OPER)
***
*** Returns the list of operand VNs read by OPER.
***
=====
(defun dag.oper:operand-vns (oper)
  (caseq (oper:operator oper)
    ( (ipload fpload)
      '( (dag.operand:vn (oper:part oper 'index) ) ) )
    ( (ipstore fpstore)
      '( (dag.operand:vn (oper:part oper 'index) )
        (dag.operand:vn (oper:part oper 'value) ) ) )
    ( (ivload fvload)
      '( (dag.operand:vn '(vbase ,(oper:part oper 'vector) ) )
        (dag.operand:vn (oper:part oper 'index) ) ) )
    ( (ivstore fvstore)
      '( (dag.operand:vn '(vbase ,(oper:part oper 'vector) ) )
        (dag.operand:vn (oper:part oper 'index) )
        (dag.operand:vn (oper:part oper 'value) ) ) )
    (vbase
      '( (dag.operand:vn '(vbase ,(oper:part oper 'vector) ) ) ) )
    ( t
      (loop (for operand in (oper:part oper 'read) ) (save
        (dag.operand:vn operand) ) ) ) ) )

=====
***
*** (DAG.OPERAND:VN OPERAND)
***
*** Returns the VN corresponding to OPERAND, which is either a variable
*** or a constant. If a previous VN for the variable or constant exists,
*** it is returned. Otherwise, one is created.
***
=====
(defun dag.operand:vn (operand)
  (if-let ( (vn ( []h *dag.name:vn* operand ) ) ) (then
    vn)
    (else
      (if (litatom operand) (then
        (assert (memq operand *dag.live-before*)
          (h operand) " DAG: Name not mentioned in LIVE-BEFORE list.")
        ( := ( []h *dag.name:vn* operand )
          (vn:create (vn:new
            type 'def
            name operand
            locations ( []h *dag.name:locations* operand ) ) ) ) )
        (else

```

```

(dag.constant:create-vn operand) ) ) ) )

=====
***
*** (DAG.CONSTANT:CREATE-VN CONSTANT)
***
*** Creates a new VN for a constant (we know there isn't currently one).
***
=====
(defun dag.constant:create-vn (constant)
  (let ( (datatype (caseq (typep constant)
    ( (double flonum) 'float
    ( t 'integer ) ) )
    (locations ( []h *dag.name:locations* constant )
    (immediate? (constant:constant-generator-mes constant) ) )
    ;*** If desired, add this constant to the live-after list of
    ;*** variables to cause it to be retained for the whole trace.
    ;***
    (if (if immediate?
      (== 'retained-register *ls.immediate-constant-action*)
      (== 'retained-register *ls.memory-constant-action*) )
      (then
        (if (for-every (x in *dag.live-after-constants*
          (! (eql x constant) ) )
          (then
            (push *dag.live-after-constants* constant) ) ) ) )
      ( ? ;*** Any locations for the constant given in the DEF at
        ;*** the beginning of the trace? If so, use them.
        ;***
        ( locations
          ( := ( []h *dag.name:vn* constant )
            (vn:create (vn:new
              type 'def
              constant? t
              name constant
              locations locations ) ) ) )
          ;*** No constant generator can make this constant and
          ;*** we're not allowed to make up new DEF locations? Make
          ;*** an XPLoad of the constant from memory.
          ;***
          ( (& (| (! *ls.new-defs-allowed*)
            (== 'load *ls.memory-constant-action*) )
            (! immediate? ) )
            (let ( (address-vn (dag.operand:vn '(address ,constant) ) )
              (load-op (caseq datatype
                (float 'fpload)
                (integer 'ipload) ) ) )
              ( := ( []h *dag.name:vn* constant )
                (vn:create (vn:new
                  type 'operation
                  name constant
                  oper '(,load-op () () )
                  operand-vns '(,address-vn) ) ) ) ) )
              ;*** Just make a normal constant VN.
              ;***
              ( t
                ( := ( []h *dag.name:vn* constant )

```

```

(vn:create (vn:new
            type      'def
            constant? t
            name      constant) ) ) ) ) )

:***
:***
:*** (DAG.LIVE-NAME:CREATE-USE NAME)
:***
:*** Makes a USE VN for a name that is live on exit from the trace. A
:*** USE VN may already exist for NAME, created by a USE on the trace.
:***
:***
:***
(defun dag.live-name:create-use ( name )
  (if-let ( (vn ([])h *dag.name:use-vn* name) ) ) (then
    vn)
  (else
    (:= ([])h *dag.name:use-vn* name)
        (vn:create (vn:new
                    type      'use
                    name      name
                    operand-vns '(.(dag.operand:vn name) ) ) ) ) ) ) ) )

:***
:***
:*** (DAG.USE-SPEC:CREATE-USE (NAME ME REGISTER) )
:***
:*** The list (NAME ME REGISTER) occurred in a USE in the trace. A USE
:*** VN for NAME is created if one doesn't exist for it already, and it
:*** is updated by the specified value location.
:***
:***
:***
(defun dag.use-spec:create-use ( (name me register) )
  (if-let ( (vn ([])h *dag.name:use-vn* name) ) ) (then
    (push (vn:locations vn) '(.(name:me me) .register) )
    vn)
  (else
    (:= ([])h *dag.name:use-vn* name)
        (vn:create (vn:new
                    type      'use
                    name      name
                    operand-vns '(.(dag.operand:vn name) )
                    locations  '(.(name:me me) .register) ) ) ) ) ) )

:***
:***
:*** (DAG.PRED-LIST:CONSTRAINING-VNS&DELAYS PRED-LIST)
:***
:*** PRED-LIST is the result returned by a call to PREDECESSORS that
:*** describes the predecessors of an operation. This returns the
:*** VNs that constrain the scheduling of an operation (e.g. conditional-
:*** conflict VNs or write-after-write VNs); operand VNs are not included.
:***
:*** The result has the form:
:***
:*** (VNS DELAYS)
:***
:*** where VNS is a list of predecessor VNs and DELAYS is a corresponding

```

```

:*** list of numbers saying how much later the current VN must be done
:*** after the predecessor VN.
:***
:***
:***
:***
(defun dag.pred-list:constraining-vns&delays ( pred-list )
  (loop (for (pred reason source-operand source-type pred-operand pred-type)
          in pred-list)
        (initial vns ()
              delays () )
        (do
          (caseq reason
            ( (conditional-conflict possible-operand-conflict)
              (push vns pred)
              (push delays 1) )
            (operand-conflict
              (?( (& (== 'written source-type)
                       (== 'read pred-type) )
                (push vns pred)
                (push delays 0) )
              ( (& (== 'written source-type)
                   (== 'written pred-type) )
                (push vns pred)
                (push delays 1) )
              ( (& (== 'read source-type)
                   (== 'written pred-type) )
                (let ( (pred-oper (vn:oper pred) ) )
                  (if (& (oper:property? pred-oper 'vector-reference)
                       (== (oper:part pred-oper 'vector)
                           (nth-elt (cdr pred-oper) pred-operand)))
                    (then
                     (push vns pred)
                     (push delays 1) ) ) ) ) ) )
              (t
                (error (list pred-list "DAG: Invalid pred-list.") ) ) ) )
          (result
            '(.(dreverse vns)
              .(dreverse delays) ) ) ) ) )

:***
:***
:*** (DAG.SET-SUCCESSOR-VNS)
:***
:*** Sets the successor fields of each VN from its successors' predecessor
:*** fields:
:***
:*** :READING-VNS      from :OPERAND-VNS
:*** :CONSTRAINED-VNS from :CONSTRAINING-VNS
:*** :OFF-LIVE-READING-VNS from :OFF-LIVE-VNS
:***
:***
:***
(defun dag.set-successor-vns ()
  (loop (for-each-vn vn) (do
    (loop (for operand-vn in (vn:operand-vns vn) ) (do
      (push (vn:reading-vns operand-vn) vn) ) )
    )
  )

```



```

(loop (for off-live-vn in (vn:off-live-vns vn) (do
  (push (vn:off-live-reading-vns off-live-vn) vn) ) )
(loop (for constraining-vn in (vn:constraining-vns vn) (do
  (push (vn:constrained-vns constraining-vn) vn) ) )
  ) )

(loop (for-each-vn vn) (do
  (:= (vn:reading-vns      vn) (reverse &&&))
  (:= (vn:constrained-vns  vn) (reverse &&&))
  (:= (vn:off-live-reading-vns vn) (reverse &&&)) ) )
  ) )

;====
;***
;*** (DAG.NAME:TRACE-WRITE-COUNT NAME)
;***
;====

(defun dag.name:trace-write-count ( name )
  ([[]h *dag.name:write-count* name) )

;====
;***
;*** (DAG.NAME:USE-VN NAME)
;***
;====

(defun dag.name:use-vn ( name )
  ([[]h *dag.name:use-vn* name) )

;====
;***
;*** (DAG.REMOVE-DEAD-VNS)
;***
;*** Removes dead code from the DAG (excessive copying by the bookkeeper
;*** can make dead code). A VN is dead if it has no readers or off-live
;*** readers. removing a dead VN can make its operands newly dead, so
;*** we use a standard iterative algorithm to remove them all.
;***
;====

(defun dag.remove-dead-vns ()
  (let ( (to-do
        (loop (for-each-vn vn)
              (when (dag.vn:dead? vn) )
              (save vn) ) ) )
    (loop (while to-do)
          (bind vn (pop to-do)
            operand-vns (vn:operand-vns vn) )
          (do
            (vn:delete vn)

            (loop (for operand-vn in operand-vns)
                  (when (dag.vn:dead? operand-vn) )
                  (do
                    (push to-do operand-vn) ) ) ) )
          (loop (for-each-vn vn) (do
            (assert (! (dag.vn:dead? vn) )

```

```

  (h vn) t "DAG: A dead VN is still left.") ) )
  ) )

(defun dag.vn:dead? ( vn )
  (&& (vn:name vn)
  (memq (vn:type vn) '(operation copy def) )
  (! (vn:reading-vns      vn) )
  (! (vn:off-live-reading-vns vn) ) ) )

;====
;***
;*** (DAG.SET-COUNTS)
;***
;====

(defun dag.set-counts ()
  (loop (for-each-vn vn) (do
    (:= (vn:predecessors-left vn)
      (+ (length (vn:operand-vns      vn) )
        (length (vn:constraining-vns vn) ) ) )
    (:= (vn:readers-left vn)
      (+ (length (vn:reading-vns      vn) )
        (length (vn:off-live-reading-vns vn) ) ) ) ) ) )
  ) )

;====
;***
;*** (DAG.PRINT [DEPTH-FIRST?])
;***
;*** Dumps out the dag in pretty debugging format; if DEPTH-FIRST? then
;*** the dag is printed out in a depth-first order. If DEPTH-FIRST? is a VN
;*** then only the VN and its descendents are printed.
;***
;====

(declare (special
  *dag.visited-vns*
  ) )

(defun dag.print (&optional depth-first?)
  (let ( (*dag.visited-vns* *vn-set.empty-set* )
    (if depth-first? (then
      (if (vn:is depth-first?) (then
        (dag.vn:depth-first-print depth-first?) )
      (else
        (loop (for vn in *ls.entry-vns*) (do
          (dag.vn:depth-first-print vn) ) ) ) ) )
    (else
      (loop (for-each-vn vn) (do
        (vn:print vn) ) ) ) )
  ) ) )

(defun dag.vn:depth-first-print ( vn )
  (if (! (vn-set:member? *dag.visited-vns* vn) ) (then
    (:= *dag.visited-vns* (vn-set:union1 &&& vn) )
    (vn:print vn)

```

```
(loop (for reading-vn in (vn:reading-vns vn) ) (do
  (dag.vn:depth-first-print reading-vn) ) )
() ) )
```

```
=====  
***  
*** (DAG.CHECK-CONSISTENCY)  
***  
*** Checks to make sure that the DAG is consistently constructed (all  
*** double links properly maintained, correct entrance and exit nodes,  
*** etc).  
***  
=====
```

```
(defun dag.check-consistency ()
  (loop (for vn in *ls.entry-vns*) (do
    (assert (&& (== 'def (vn:type vn) )
      (! (vn:operand-vns vn) )
      (! (vn:constraining-vns vn) ) )
      (h vn) ) ) )
  (loop (for vn in *ls.exit-vns*) (do
    (assert (&& (| (== 'use (vn:type vn) )
      (! (vn:name vn) ) )
      (! (vn:reading-vns vn) )
      (! (vn:constrained-vns vn) ) )
      (h vn) ) ) )
  (loop (for-each-vn vn) (do
    (if (&& (! (vn:operand-vns vn) )
      (! (vn:constraining-vns vn) ) )
      (then
        (assert (&& (memq vn *ls.entry-vns*)
          (== 'def (vn:type vn) ) )
          (h vn) ) ) )
      (loop (for operand-vn in (vn:operand-vns vn) ) (do
        (assert (memq vn (vn:reading-vns operand-vn) )
          (h vn) t (h operand-vn) ) ) )
      (loop (for constraining-vn in (vn:constraining-vns vn) ) (do
        (assert (memq vn (vn:constrained-vns constraining-vn) )
          (h vn) t (h constraining-vn) ) ) )
      (loop (for off-live-vn in (vn:operand-vns vn) ) (do
        (assert (memq vn (vn:off-live-reading-vns off-live-vn) )
          (h vn) t (h off-live-vn) ) ) )
      (if (&& (! (vn:reading-vns vn) )
        (! (vn:constrained-vns vn) ) )
        (then
          (assert (&& (memq vn *ls.exit-vns*)
            (== 'use (vn:type vn) ) )
            (h vn) ) ) )
        (loop (for reading-vn in (vn:reading-vns vn) ) (do
          (assert (memq vn (vn:operand-vns reading-vn) )
            (h vn) t (h reading-vn) ) ) )
        (loop (for constrained-vn in (vn:constrained-vns vn) ) (do
          (assert (memq vn (vn:constraining-vns constrained-vn) )
            (h vn) t (h constrained-vn) ) ) )
```

```
(loop (for off-live-reading-vn in (vn:off-live-reading-vns vn) ) (do
  (assert (memq vn (vn:off-live-vns off-live-reading-vn) )
    (h vn) t (h off-live-reading-vn) ) ) ) )
() )
```

```
=====  
***  
*** (CYCLE:DELAYS:CONSTRAINED-CYCLE CYCLE PD SD CD)  
***  
=====
```

```
(defun cycle:delays:constrained-cycle ( cycle pd sd cd )
  (+ cycle (+ cd (max 0 (- pd sd) ) ) ) )
```

```
=====  
***  
*** (CYCLE:DELAYS:CONSTRAINING-CYCLE CYCLE PD SD CD)  
***  
=====
```

```
(defun cycle:delays:constraining-cycle ( cycle pd sd cd )
  (- cycle (+ cd (max 0 (- pd sd) ) ) ) )
```

```
*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for *****
*** educational and research purposes only, by the faculty, students, *****
*** and staff of Yale University only. *****
*** *****
```

```
=====
: This file should be included at compile time by all modules of
: LIST-SCHEDULER.
:=====
```

```
(declare (special
  *hash-table-not-found*

  ;*** MACHINE-MODEL
  ;***
  *ls-mc-universe*
  *ls-register-bank-mes*
  *ls-constant-generator-mes*
  *ls-functional-unit-mes*

  ;*** VM
  ;***
  *ls-vn-universe*
  *vn-set.empty-set*

  ;*** DAG
  ;***
  *ls.entry-vns*
  *ls.exit-vns*
  *ls.new-defs-allowed?*
  *ls.loop-trace?*

  ;*** SCHEDULER
  ;***
  *ls.max-schedule-size*
  *ls.schedule*
  *ls.last-cycle*
  *ls.cycle*

  ;*** Options
  ;***
  *ls.immediate-constant-action*
  *ls.memory-constant-action*
  *ls.trace-information?*
  *ls.timed-functions*
) )
```

```
(eval-when (compile)
  (build '(
    utilities:visible-fields
    utilities:sharp-sharp
    utilities:bit-set
    utilities:object-universe
    utilities:array
    utilities:loop-min-max

    interpreter:naddr

    list-scheduler:machine-model
    list-scheduler:vn
```

```
) ) )
```

```

*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for
*** educational and research purposes only, by the faculty, students,
*** and staff of Yale University only.
*****

=====
: The ELI model used by Rutt's CG experiments.
=====

(alias "" 'atomconcat)

(defmacro def-bus ( me1 me2 number-of-values )
  (let ( (resource-name (atomconcat me1 '<->' me2) ) )
    (eval-when (eval load)
      (def-connection ,me1 ,me2 ( (,resource-name) ) )
      (def-connection ,me2 ,me1 ( (,resource-name) ) )
      (def-resource-class ,resource-name ,number-of-values) ) ) )

(defun em.cluster ( 1 (unique-name unique-fields) )
  (let ( (ci (atomconcat 'c 1) )
        (ci-8 (atomconcat 'c (mod (- 1 3) 8) ) )
        (ci-1 (atomconcat 'c (mod (- 1 1) 8) ) )
        (ci+1 (atomconcat 'c (mod (+ 1 1) 8) ) )
        (ci+8 (atomconcat 'c (mod (+ 1 3) 8) ) ) )
    (

      *** Register bank CiR
      ***
      (def-me (name (,ci 'r) )
              (type register-bank)
              (size 64)
              (inputs ( (,ci-8 'r)
                        (,ci-1 'r)
                        (,ci 'r)
                        (,ci+1 'r)
                        (,ci+8 'r)
                        (,ci 'c)
                        (,ci 'm)
                        (,ci '+)
                        (,ci '=)
                        ..(if unique-name
                            '( (,ci unique-name) )
                            ( ) ) ) )
              (read-ports 8)
              (write-ports 8)
              (read-resources ( (,ci 'r-read) ) )
              (write-resources ( (,ci 'r-write) ) ) )
      (def-resource-class ,(,ci 'r-read) 8)
      (def-resource-class ,(,ci 'r-write) 8)

      *** Integer adder Ci+
      ***
      (def-me (name (,ci '+) )
              (type functional-unit)
              (delay 0)
              (inputs ( (,ci 'r)
                        (,ci 'c) ) )
              (resources ( (,ci '+) ) )
              (operators (float fix inot idiv isub leq imax imin

```

```

iadd ineg ior ige ilt ile ine iexp isul igt
iand iabs bitrev iland ilor) ) )
(def-resource-class ,(,ci '+) 1)

*** Integer tester Ci=
***
(def-me (name (,ci '=) )
        (type functional-unit)
        (delay 0)
        (inputs ( (,ci 'r)
                  (,ci 'c) ) )
        (resources ( (,ci '=) ) )
        (operators (if-true if-false if-ilt if-igt if-leq if-ine
                     if-ile if-ige if-leq if-ine if-ile if-ige)))
(def-resource-class ,(,ci '=) 1)

*** Memory CiM
***
(def-me (name (,ci 'm) )
        (type functional-unit)
        (delay 8)
        (inputs ( (,ci 'r)
                  (,ci 'c) ) )
        (resources ( (,ci 'm) ) )
        (operators (vbase ivload fvload ipload fpload ivstore
                    fvstore ipstore fpstore) ) )
(def-resource-class ,(,ci 'm) 1)

*** Constant generator CiC
***
(def-me (name (,ci 'c) )
        (type constant-generator)
        (resources ( (,ci 'c) ) )
        (constraint-function em.immediate-constant?) )
(def-resource-class ,(,ci 'c) 1)

*** Unique functional-unit
***
..(if unique-name (then
  ( (def-me (name (,ci unique-name) )
            (type functional-unit)
            (inputs ( (,ci 'r)
                      (,ci 'c) ) )
            (resources ( (,ci unique-name) ) )
            ..unique-fields)
      (def-resource-class ,(,ci unique-name) 1) )
  (else
   ( ) ) )

*** Bus connections
***
(def-bus ,(,ci-8 'r) ,(,ci 'r) 1)
(def-bus ,(,ci-1 'r) ,(,ci 'r) 1)
) )

(:= em.floating-test
  '(f=
    ( (delay 0)
      (operators (if-flt if-fgt if-feq if-fne if-file if-ige
                  flt fgt feq fne file fgo) ) ) ) )

```

```

(:= em.floating-multiply
  *(f+
    ( (delay 3)
      (operators (fsul imul) ) ) ) )

(:= em.floating-add
  *(f+
    ( (delay 2)
      (operators (fsub fadd fneg fain fmax fabs fdiv cos sin sqrt flt fgt
                  feq fne fle fge frecip frsqrt) ) ) )

(defun em.immediate-constant? ( constant )
  (? ( inusp constant)
     (& (< constant 2047)
         (> constant -2048) ) )
  ( (= 0 constant)
    t )
  ( (consp constant)
    (== 'address (car constant) ) )
  ( t
    ) ) )

(defun load-constant? ( constant )
  (if (em.immediate-constant? constant) (then
      (!= 'load *ls.immediate-constant-action*) )
      (else
        t ) ) )

(eval '(def-machine-model
  ..(em.cluster 0 em.floating-add)
  ..(em.cluster 1 em.floating-add)
  ..(em.cluster 2 em.floating-multiply)
  ..(em.cluster 3 em.floating-add)
  ..(em.cluster 4 em.floating-add)
  ..(em.cluster 5 ( ) )
  ..(em.cluster 6 em.floating-multiply)
  ..(em.cluster 7 em.floating-add)
  ) )

```

```

***-----*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for *****
*** educational and research purposes only, by the faculty, students, *****
*** and staff of Yale University only. *****
***-----*****

```

#### GENERATE CODE HOOK

This module provides a hook to save a snapshot of a trace and the disambiguator information in a file, so that we can exercise the code generator without needing to run the whole compiler.

Load this file into the compiler running the ideal code generator. The `*TR.GENERATE-CODE-HOOK*` will automatically be set. A breakpoint will occur during each trace. To save that trace in a file do:

```
(SAVE-TRACE name)
```

A snapshot of the trace will be saved in the file `name-n.TRACE`, where "n" is the number of the trace.

The file written looks like:

```

(:= name-n
  symbol-table
  live-before
  source-record-list
  live-after
  predecessors-results)

```

where "live-before", "source-record-list", and "live-after" are the arguments to `GENERATE-CODE`, and "predecessors-results" is a list of the results of `PREDECESSORS` called on each of the operations in the trace in turn (not including any `DEFS` or `USES`); the datum handed to `PREDECESSORS` is the integer position of the operation in the trace. "symbol-table" is a list of tuples of the form:

```
(name datatype rank)
```

This file can then be used later on with `TEST-BED` to simulate the compiler interface without needing the whole compiler present.

By setting the global `*GCH.AUTOMATIC?*` to some name `XXX`, then `SAVE-TRACE` will automatically be called with argument `'XXX` on each trace.

```

(declare (special
  *gch.live-before*
  *gch.source-record-list*
  *gch.live-after*
  *gch.automatic?*
  ) )

```

```
(:= *tr.generate-code-hook* 'ls.generate-code-hook)
```

```
(:= *gch.trace-number* 0)
(:= *gch.automatic?* ())
```

```

(defun ls.generate-code-hook
  (*gch.live-before* *gch.source-record-list* *gch.live-after* )

```

```

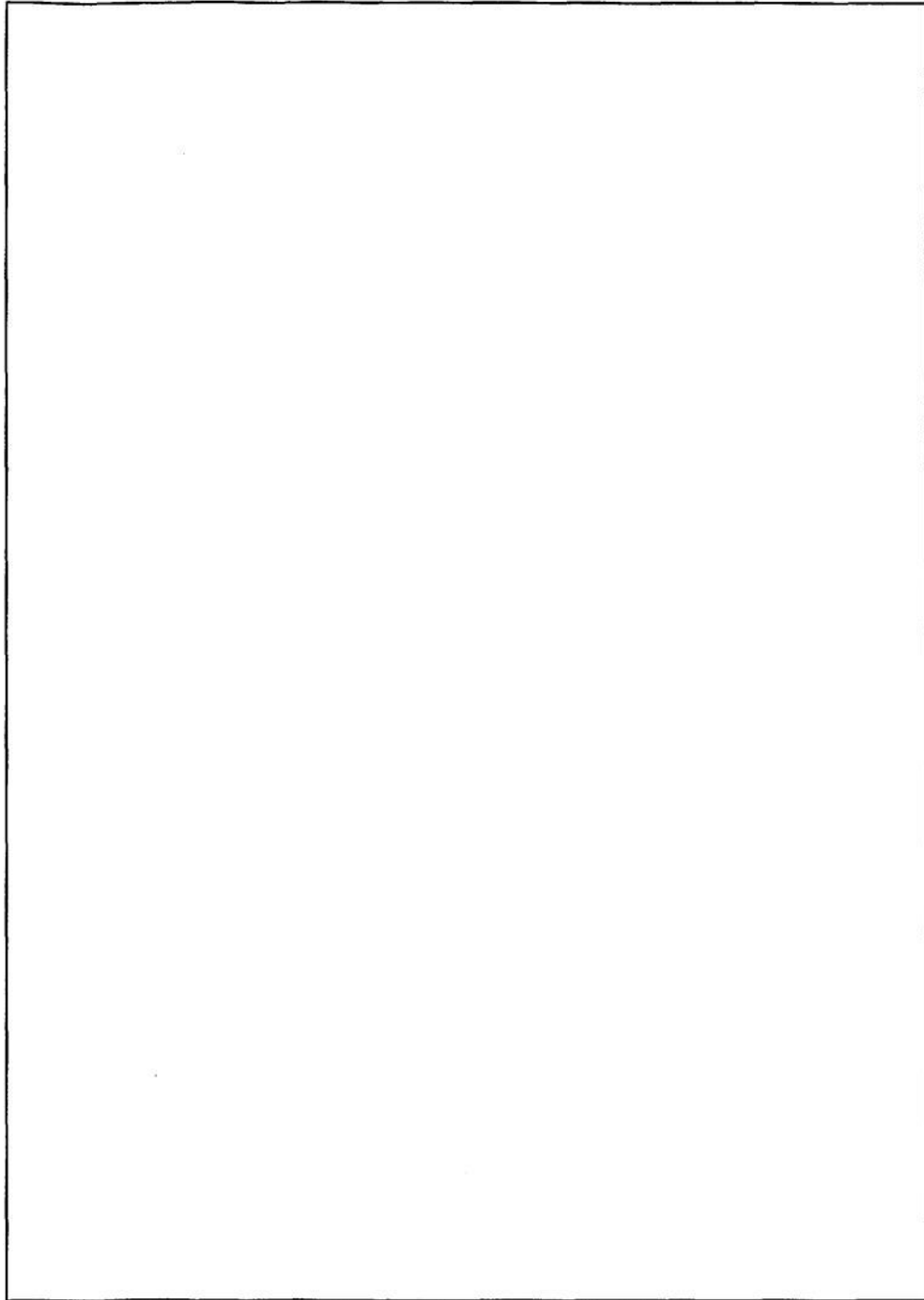
(if (! *gch.source-record-list*) (then
  (:= *gch.trace-number* 0) )
  (else
    (:= *gch.trace-number* (+ &&& 1) )
    (if *gch.automatic?* (then
      (save-trace *gch.automatic?*) )
      (else
        (break-point ls.generate-code-hook) ) ) ) )
  ) )

```

```

(defun save-trace ( name )
  (let ( (trace-name (atomconcat name "-" *gch.trace-number*)) )
    (iota ( file (stconcat trace-name ".trace") '(out newversion) ) )
      (let ( (new-source-record-list ())
            (predecessors-list ())
            (first-op
              (oper:operator (car (car *gch.source-record-list*))) )
            (last-op
              (oper:operator
                (car (last-elt *gch.source-record-list*)) ) ) )
        (if (!= 'def first-op) (then
          (push *gch.source-record-list* '( (def) () () () ) ) )
          (if (!= 'use last-op) (then
            (:= *gch.source-record-list*
              (append! &&& '( (use) () () () ) ) ) )
            (start-trace)
            (loop (for (oper trace-direction datum off-live)
              in *gch.source-record-list*
                (incr i from 0)
              (do
                (push new-source-record-list
                  '(,oper ,trace-direction ,i ,off-live) )
                (push predecessors-list
                  (if (oper:property? oper 'pseudo-op)
                    ()
                    (predecessors oper trace-direction i) ) ) )
              (result
                (:= new-source-record-list (dreverse &&&))
                (:= predecessors-list (dreverse &&&)) ) )
            (fmsg file
              "(:= " trace-name " "(" t t
                " " t
                (e (loop (for-each-name name) (do
                  (msg " "
                    name " "
                    (name:datatype name) " "
                    (name:rank name) name) " t) ) ) )
              )" t t
              (h *gch.live-before* 10000 10000) t t
              (h new-source-record-list 10000 10000) t t
              (h *gch.live-after* 10000 10000) t t
              (h predecessors-list 10000 10000) t t
              )" t)
            (filename file) ) ) ) )

```



```

*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for
*** educational and research purposes only, by the faculty, students,
*** and staff of Yale University only.
*****
=====
(HEIGHTS-DEPTHS.ASSIGN)
Sets the :DEPTH of each VN to be the minimum estimated time that the
VN's operation can be scheduled, using the :DELAY of operators to
make the estimate (and assuming infinite resources).
Sets :HEIGHT to be the minimum time from an exit of the DAG.
=====

(eval-when (compile load)
  (include list-scheduler:declarations) )

(defun heights-depths.assign ()
  (loop (for vn in *ls.exit-vns*) (do
    (hd.vn:assign-depth vn) ) )

  (loop (for vn in *ls.entry-vns*) (do
    (hd.vn:assign-height vn) ) )
  () )

(defun hd.vn:assign-depth ( vn )
  (if (vn:depth vn) (then
    (vn:depth vn) )

    (else
      (let ( (delay (hd.vn:estimated-delay vn) ) )
        (:= (vn:depth vn)
          (loop (for operand-vn in (vn:operand-vns vn) )
            (reduce max 0 (+ (hd.vn:estimated-delay operand-vn)
              (hd.vn:assign-depth operand-vn))))))

        (:= (vn:depth vn)
          (loop (for constraining-vn in (vn:constraining-vns vn) )
            (for constraining-delay in (vn:constraining-delays vn) )
              (reduce max
                (vn:depth vn)
                (cycle:delays:constrained-cycle
                  (hd.vn:assign-depth constraining-vn)
                  (hd.vn:estimated-delay constraining-vn)
                  delay
                  constraining-delay) ) ) )

          (vn:depth vn) ) ) ) )

(defun hd.vn:assign-height ( vn )
  (if (vn:height vn) (then
    (vn:height vn) )

    (else
      (let ( (delay (hd.vn:estimated-delay vn) ) )
        (:= (vn:height vn)
          (loop (for reading-vn in (vn:reading-vns vn) )

```

```

(reduce max 0 (+ delay
  (hd.vn:assign-height reading-vn))))))

(:= (vn:height vn)
  (loop (for constrained-vn in (vn:constrained-vns vn) )
    (reduce max
      (vn:height vn)
      (cycle:delays:constraining-cycle
        (hd.vn:assign-height constrained-vn)
        delay
        (hd.vn:estimated-delay constrained-vn)
        (vn:constrained-vn:delay vn constrained-vn)
        ) ) ) )

(vn:height vn) ) ) )

(defun hd.vn:estimated-delay ( vn )
  (caseq (vn:type vn)
    ( (def defl use use1 copy pseudo-op)
      0)
    ( operation
      (+ 1 (operator:min-delay (vn:operator vn) ) ) )
    ( t
      (error (list vn "HD.VN:ESTIMATED-DELAY: Case error.") ) ) ) )

```



```

*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for
*** educational and research purposes only, by the faculty, students,
*** and staff of Yale University only.
*****
*****
=====
: Top level driver for invoking the list-scheduling compiler.
=====
(eval-when (compile load)
  (include list-scheduler:declarations) )

(declare (special
  *lsex.code*           ;*** The most recent TinyLisp code.
  *lsex.actual-params* ;*** The most recent actual parameters.
  *lsex.unoptimized-naddr* ;*** The NADDR straight out of TinyLisp.
  *lsex.optimized-naddr*  ;*** The flow-analysis-optimized NADDR.
  *lsex.compacted-program* ;*** The compacted ELI program.

  *gch.automagic?*      ;*** Hack debugging
) )

(declare
  (lexpr time-functions)
  (lexpr options.print)
  (lexpr print-functions-times)
)

(options.set 'timed-functions '(
  compile-tiny-lisp
  fa.analyze&optimize
  compact
  generate-code
  bug.assign-mes
  sch.schedule
  simulate
) )

(defun lsex ( &optional (code *lsex.code*)
              (actual-params *lsex.actual-params*) )
  (lsex.initialize ( ) )

  (:= *lsex.code* code)
  (:= *lsex.actual-params* actual-params)

  (msg 0 t (e (options.print) ) 0 t)

  (unwind-protect
    (let ( )
      (if *ls.timed-functions* (then
        (apply 'time-functions *ls.timed-functions*) ) )

      (if (atom *lsex.code*) (then
        (:= *lsex.unoptimized-naddr*
          (skex.filename:read-naddr *lsex.code*) )
        (:= *lsex.optimized-naddr*
          (fa.analyze&optimize *lsex.unoptimized-naddr* ( ) ) ) )
        (else

```

```

(:= *lsex.unoptimized-naddr*
  (compile-tiny-lisp *lsex.code*) )
(:= *lsex.optimized-naddr*
  (fa.analyze&optimize *lsex.unoptimized-naddr* t) ) ) )

(:= *lsex.compacted-program*
  (mis-to-eli (compact *lsex.optimized-naddr*) ) )

(msg 0 t)
(simulate *lsex.compacted-program* actual-params)
(simulator.print-execution-statistics)

(if *ls.timed-functions* (then
  (msg 0 t (e (print-function-times) ) t) ) )

( )

(if *ls.timed-functions* (then
  (apply 'untime-functions *ls.timed-functions*) ) )
(:= *gch.automagic?* ( ) )

( ) )

(defun lsex.re-cg ( )
  (lsex.initialize t)
  (unwind-protect
    (let ( )
      (if *ls.timed-functions* (then
        (apply 'time-functions *ls.timed-functions*) ) )

      (:= *lsex.compacted-program*
        (mis-to-eli (compact *lsex.optimized-naddr*) ) )

      (msg 0 t)
      (simulate *lsex.compacted-program* *lsex.actual-params*)
      (simulator.print-execution-statistics)

      (if *ls.timed-functions* (then
        (msg 0 t (e (print-function-times) ) t) ) ) )

      (if *ls.timed-functions* (then
        (apply 'untime-functions *ls.timed-functions*) ) )
      (:= *gch.automagic?* ( ) )

    ( ) )

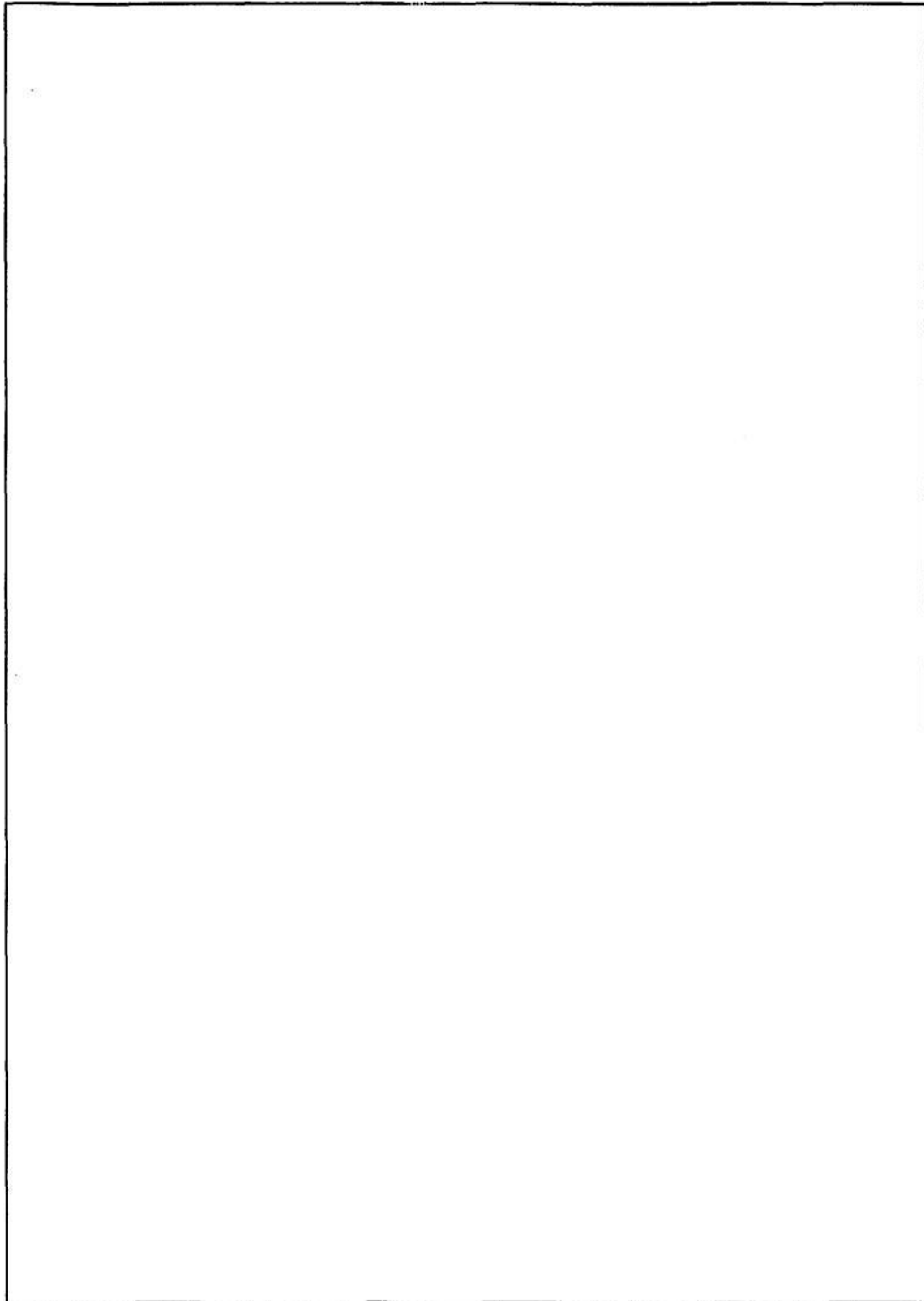
  ( ) )

(defun lsex.initialize ( re-cg? )
  (tr.initialize)
  (initialize-code-generator)
  (if (not re-cg?) (then
    (fa.initialize)
    (de.initialize)

    (:= *lsex.code* ( ) )
    (:= *lsex.actual-params* ( ) )
    (:= *lsex.unoptimized-naddr* ( ) )
    (:= *lsex.optimized-naddr* ( ) ) ) ) )

  (let ( (old (gcgag t) ) )
    (gc)
    (gcgag old)
  ( ) )

```



```
***-----*****
*** Copyright (C) 1983 John R. Ellis. This file may be used for *****
*** educational and research purposes only, by the faculty, students, *****
*** and staff of Yale University only. *****
***-----*****
```

```
=====
: LIST-SCHEDULER OPTIONS
:
:=====
```

```
(eval-when (compile)
  (build '(utilities:options) ) )
```

```
(def-option *ls.immediate-constant-action* 'retained-register list-scheduler: "
Determines the strategy for handling immediate constants. One of:
```

LOAD

The constant is always taken from the nearest constant generator.

REGISTER

The constant is assigned one or more registers at the beginning of the trace.

RETAINED-REGISTER

The constant is assigned one or more registers which are kept live for the entire trace. Approximates loop-invariant motion on constants.

")

```
(def-option *ls.memory-constant-action* 'retained-register list-scheduler: "
Determines the strategy for handling non-immediate constants that must
be loaded from memory. One of:
```

LOAD

The constant is always loaded from memory on the current trace.

REGISTER

The constant is assigned one or more registers at the beginning of the trace.

RETAINED-REGISTER

The constant is assigned one or more registers which are kept live for the entire trace. Approximates loop-invariant motion on constants.

")

```
(def-option *ls.trace-information?* t list-scheduler: "
If true, then various markers are left laying around in compacted code
giving the origins of the code.
```

")

```
(def-option *ls.timed-functions* () list-scheduler: "
If non-(), then this is the list of functions that will be timed during
each LSEX run.
```

")

```

***-----*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for *****
*** educational and research purposes only, by the faculty, students, *****
*** and staff of Yale University only. *****
***-----*****

```

MACHINE MODEL

This module implements the basic machine model used for the list scheduler. A machine consists of several machine elements (MEs) forming a connected graph.

ME

```

=====
(def-struct me

```

```

  name      : The print-name of this element.
  number    : A unique number for bit-sets and vector
             : manipulation.
  type      : One or REGISTER-BANK, FUNCTIONAL-UNIT,
             : CONSTANT-GENERATOR.
  inputs    : The MEs connected to the inputs and outputs
  outputs   : of this ME.
  operators : FUNCTIONAL-UNIT: List of NADDR operators
             : implemented by this ME.
  delay     : FUNCTIONAL-UNIT: The time it takes for the fu
             : to complete, minus 1. I.e. a 1-cycle integer
             : adder has 0 delay.
  constraint-function : FUNCTIONAL-UNIT, CONSTANT-GENERATOR: A boolean
             : function that takes one argument, a VN, returning
             : true if the VN can be calculated on this ME.
  resources : FUNCTIONAL-UNIT, CONSTANT-GENERATOR: The resource
             : request needed to use the ME.
  read-resources : REGISTER-BANK: The resource requests needed
  write-resources : to do a read or write operation on the bank.
  size        : REGISTER-BANK: The number of registers in the
             : bank.
  read-ports  : REGISTER-BANK: The numbers of read and write
  write-ports : ports.
  registers-left : REGISTER-BANK: The number of unallocated registers.
  values      : REGISTER-BANK: A vector of the current contents
             : of the registers.
  cycles-free : REGISTER-BANK: A vector giving the first cycle that
             : an unused register is available.
  value-names : REGISTER-BANK: A vector of the current names
             : of the register contents.

```

```

avoid?      : REGISTER-BANK: A vector of booleans; Ith element
             : is true if register I should be avoided if possible
             : during allocation.
)

```

```

=====
(DEF-MACHINE-MODEL
  (DEF-ME (NAME ...) (TYPE ...) ...)
  (DEF-ME ...)
  (DEF-ME ...)
  (DEF-CONNECTION ...)
)

```

Defines a machine model. Undocumented.

```

(MACHINE-MODEL.INITIALIZE)
  Initializes this module.

```

```

(NAME:ME NAME)
  Returns the ME that has a given name, raising an error if not found.

```

```

(NAME-LIST:ME-LIST NAME-LIST)
  Maps a list of ME names onto a corresponding list of MEs.

```

```

(ME-LIST:NAME-LIST ME-LIST)
  Returns the list of names corresponding to a list of MEs.

```

```

$#ME name
  Interactive syntax for referencing an ME by name.

```

```

(ME:ME:RESOURCES ME1 ME2)
  ME2 should be one of the outputs of ME1; returns the resource request
  needed to move a value from the output of ME1 to the input of ME2.

```

```

(ME:DATUM:OK? ME DATUM)
  True if DATUM satisfies the :CONSTRAINT-FUNCTION of ME. If ME is
  a constant generator, DATUM is a constant; otherwise DATUM is a VN.

```

```

*LS.REGISTER-BANK-MES*
*LS.FUNCTIONAL-UNIT-MES*
*LS.CONSTANT-GENERATOR-MES*
  Lists of the various types of MEs.

```

```

(OPERATOR:FUNCTIONAL-UNIT-MES OPERATOR)
  Returns the list of functional unit MEs that implement a NADDR operator.

```

```

(OPERATOR:MIN-DELAY OPERATOR)
  Returns the minimum delay (over all functional units) needed to compute
  a NADDR operator.

```

```

(CONSTANT:CONSTANT-GENERATOR-MES CONSTANT)
  Returns the constant generators capable of generating CONSTANT;
  returns () if no generator generates the constant. Uses a hash
  table that lives for the life of the machine model to remember
  previous results.

```

```

=====
(visible-fields me name)

```

```

(eval-when (compile load)
  (include list-scheduler:declarations) )

```

```

(declare (special
  *m.name:me*
  *m.me:me:resources*
  *m.connections*
  *m.operator:mes*
  *m.operator:min-delay*
  *m.constant:mes*
  ))

(eval-when (eval compile load)
  (def-object-universe *ls.me-universe*
    (object-name me)
    (mapping-type numbered-objects)
    (object-number-function
      (lambda (object)
        (me:number object) ) )
    (set-object-number-function
      (lambda (object number)
        (:= (me:number object) number) ) )
    (initial-size 100) ) )

(defun machine-model.initialize ()
  (resource.initialize)
  (shortest-path.initialize)
  (me-universe.initialize)

  (:= *m.name:me* (hash-table:create () () () ) )
  (:= *m.operator:mes* (hash-table:create () () () ) )
  (:= *m.constant:mes* (hash-table:create 'equal () ) )
  (:= *m.operator:min-delay* (hash-table:create () () 100000) )

  (:= *m.me:me:resources* ())
  (:= *m.connections* ())
  (:= *ls.register-bank-mes* ())
  (:= *ls.constant-generator-mes* ())
  (:= *ls.functional-unit-mes* ())
  ())

(defun name:me (name)
  (if ([h *m.name:me* name]
      (error (list name "NAME:ME: Invalid ME name.")) ) ) )

(defun name-list:me-list (name-list)
  (loop (for name in name-list) (save
    (name:me name) ) ) )

(defun me-list:name-list (me-list)
  (loop (for me in me-list) (save
    (me:name me) ) ) )

(defun me:me:resources (me1 me2)
  (? ( ( me1
        () )
      ( ( me2
        () )
      )
  ) )

```

```

( ( [h *m.me:me:resources* (me:number me1) (me:number me2) ] ) ) )

(defun me:create (me)
  (me-universe:add me)
  (:= ([h *m.name:me* (me:name me) ] me) ) )

(defun machine-model.finalize ()

  *** Construct the lists of the different type MEs.
  ***
  (loop (for-each-me me) (do
    (caseq (me:type me)
      (register-bank (push *ls.register-bank-mes* me) )
      (constant-generator (push *ls.constant-generator-mes* me) )
      (functional-unit (push *ls.functional-unit-mes* me) )
      (t (error (list me "Case error.") ) ) ) ) ) )

  *** Initialize the register bank MEs.
  ***
  (loop (for me in *ls.register-bank-mes* (do
    (:= (me:values me) (makevector (me:size me) ) )
    (:= (me:cycles-free me) (makevector (me:size me) ) )
    (:= (me:value-names me) (makevector (me:size me) ) )
    (:= (me:avoid? me) (makevector (me:size me) ) )
    (:= (me:registers-left me) (me:size me) ) ) ) )

  *** Replace ME names by the actual MEs, and construct
  *** ME:OUTPUTS from :INPUTS.
  ***
  (loop (for-each-me me) (do
    (:= (me:inputs me) (name-list:me-list $$$) )
    (:= (me:outputs me) (name-list:me-list $$$) ) ) )

  (loop (for-each-me me) (do
    (loop (for input-me in (me:inputs me) ) (do
      (:= (me:outputs input-me) (unionqi $$$ me) ) ) ) ) )

  (loop (for-each-me me) (do
    (:= (me:inputs me) (dreverse $$$) )
    (:= (me:outputs me) (dreverse $$$) ) ) )

  (loop (for-each-me me) (do
    (loop (for input-me in (me:inputs me) ) (do
      (assert (memq me (me:outputs input-me) ) ) ) )
    (loop (for output-me in (me:outputs me) ) (do
      (assert (memq me (me:inputs output-me) ) ) ) ) ) )

  *** Instantiate the names of resource classes within resource
  *** requests.
  ***
  (loop (for-each-me me) (do
    (:= (me:resources me) (resource-request:instantiate $$$) )
    (:= (me:read-resources me) (resource-request:instantiate $$$) )
    (:= (me:write-resources me) (resource-request:instantiate $$$) ) ) )

  *** Construct the ME->ME connection-resource table.
  ***
  (:= *m.me:me:resources*
    (array:new '(, (me-universe:size) , (me-universe:size) ) ) )

```

```

(loop (for (name1 name2 resource-request) in *mm.connections*)
      (bind me1 (name:me name1)
            me2 (name:me name2) )
      (do
        (assert (memq me2 (me:outputs me1) ) )
        (:= ([h *mm.me:me:resources* (me:number me1) (me:number me2) ]
             (resource-request:instantiate resource-request) ) ) )

        ;** Construct the operator tables.
        ;**
        (loop (for-each-me me) (do
              (loop (for operator in (me:operators me) ) (do
                    (push ([h *mm.operator:mes* operator) me)
                    (:= ([h *mm.operator:min-delay* operator]
                        (min &&& (me:delay me) ) ) ) ) ) )
              ;** Compute the shortest path table.
              ;**
              (shortest-path.compute)
              ( ) )
        )
      )
)

```

```

(defmacro def-me clauses
  '(me:create (me:new
               ..(loop (for (clause-name clause-body) in clauses) (append
                        '(,clause-name ',clause-body) ) ) ) ) )
)

```

```

(defmacro def-machine-model me-defs
  '(eval-when (eval load)
    (machine-model.initialize)
    ..me-defs
    (machine-model.finalize) ) )
)

```

```

(defmacro def-connection ( me1 me2 resource-request )
  '(push *mm.connections* '(,me1 ,me2 ,resource-request) ) )
)

```

```

(def-sharp-sharp me
  '(name:me ',(read) ) )
)

```

```

(defun me:datum:ok? ( me datum )
  (|| (! (me:constraint-function me) )
      (funcall (me:constraint-function me) datum) ) )
)

```

```

(defun operator:min-delay ( operator )
  ([h *mm.operator:min-delay* operator] )
)

```

```

(defun operator:functional-unit-mes ( operator )
  ([h *mm.operator:mes* operator] )
)

```

```

(defun constant:constant-generator-mes ( constant )
  (let ( (mes ([h *mm.constant:mes* constant] ) )
        (if (== mes *hash-table.not-found*) (then
            (:= ([h *mm.constant:mes* constant]
                (loop (for me in *ls.constant-generator-mes*
                      (when (me:datum:ok? me constant) )
                    (save me) ) ) )
            )
          )
        )
  )
)
)

```

```

(else
  mes) ) ) )
)

```

```

*****
*** Copyright (C) 1983 John R. Ellis. This file may be used for
*** educational and research purposes only, by the faculty, students,
*** and staff of Yale University only.
*****

```

```

=====
This module convert compacted microinstructions into ELI code.

To convert a list of MI's (such as returned by the compacter) into
ELI code:

```

```

(MIS-TO-ELI MI-LIST)

```

```

The conversion is very simple minded-- preorder traversal of the flow
graph, starting at the one node that has the (START) source. Each node
is visited only once (using the MI:TRANSLATED-TO-SOURCE flag to remember
visits).

```

```

=====
(eval-when (compile load)
  (include trace:declarations) )

```

```

(declare (special
  *mte.mis-to-do*      ;*** stack of mi's possibly not yet processed
) )

```

```

=====
***
*** (MIS-TO-ELI MI-LIST)
***
=====

```

```

(defun mis-to-eli ( mi-list )
  (assert (listp mi-list) )
  (:= *mte.mis-to-do* ( ) )

```

```

  (let ( (instr-stream ( ) )
        (instr ( ) ) )

```

```

    ;*** Clear the "translated" flag of every mi, and if it has
    ;*** no predecessors, push it on our to-do stack.
    ;***

```

```

    (loop (for mi in mi-list) (do
      (:= (mi:translated-to-source mi) ( ) )
      (if (not (mi:preds mi) ) (then
        (push *mte.mis-to-do* mi) ) ) ) )

```

```

    ;*** While there are untranslated mi's, do
    ;*** convert each one to instr
    ;***

```

```

    (loop (while *mte.mis-to-do*)
      (bind mi (pop *mte.mis-to-do*))
      (when (not (mi:translated-to-source mi) ) )
      (bind instr (mte.mi:instr mi) )
      (when instr)

```

```

    (do
      (push instr-stream instr) ) )

```

```

  (dreverse instr-stream) ) )

```

```

=====
***
*** (MTE.MI:INSTR MI)
***
*** Translates the single mi MI into an ELI instruction. As a side
*** effect, it pushes on *MTE.MIS-TO-DO* the successors of this MI. Only
*** the labels of jumps are changed. The new instruction is also placed
*** in the translated-to-source field of the mi. It is possible that
*** MIs won't have any source (they are a result of bookkeeping) and
*** they are totally ignored here.
***
=====

```

```

(defun mte.mi:instr ( mi )
  (:= (mi:translated-to-source mi) t)

```

```

  (let ( (popers ( ) )
        (cond-oper ( ) )
        (trace-dir (mi:trace-direction mi) )
        (succs (mi:succs mi) ) )

```

```

    ;*** Generate a label for this MI if this isn't the first MI.
    ;***

```

```

    (if (mi:preds mi) (then
      (push popers '( ( ) label ,(mte.mi:instr-label mi) ) ) ) ) )

```

```

    ;*** Translate each of the machine operations:
    ;***

```

```

    (loop (for oper in (mi:source mi) )
      (bind (fu operator . operands) oper)

```

```

    (do
      (if (== fu 'stop) (then ;*** Gross...
        (push popers '( ( ) stop) ) )
      (else

```

```

        (caseq operator
          ( ( ( ) trace copy noop)
            (push popers oper) )
          (t
            (caseq (operator:group operator)
              (goto)

```

```

              ( (if-compare if-boolean)
                (push popers
                  '(,fu
                    .operator
                    ..operands
                    ,(if (== (car trace-dir) 'right) (then
                      (mte.mi:instr-label
                        (car succs))))
                    ,(if (== (car trace-dir) 'left) (then
                      (mte.mi:instr-label
                        (car succs))))))
                (pop succs)
                (pop trace-dir) )
              (t
                (push popers oper) ) ) ) ) ) ) )

```

```

    ;*** Push the successors on the stack.
    ;***

```

```

    (loop (for succ in (mi:succs mi) )
      (when (not (mi:translated-to-source succ) ) )

```

```

(do
  (push *mte.mis-to-do* succ) ) )

;*** At this point, SUCCS contains one MI, the fall-through.
;*** If it's not the first thing on the to-do stack, then
;*** we need to do an explicit GOTO, because the next label
;*** generated (the top of the stack) isn't the one we want
;*** for the fall-through.
;***
(if (! succs) (then
  (assert (! (mi:succs m1) ) ) ) ;*** This must be the last MI.
  (else
    (assert (== 1 (length succs) ) )
    (if (|| (! *mte.mis-to-do*)
      (!= (car succs) (car *mte.mis-to-do*) ) )
      (then
        (push popers
          '( () goto ,(mte.mi:instr-label (car succs) ) ) ) ) ) ) ) )
  (:= (mi:translated-to-source m1) (dreverse popers) ) ) )

```

```

;*****
;***
;*** (mte.MI:INSTR-LABEL MI)
;***
;*** Returns the label of an MI.
;***
;*****

```

```

(defun mte.mi:instr-label ( mi )
  (atomconcat 'l (mi:number mi) ) )

```



```

***-----*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for *****
*** educational and research purposes only, by the faculty, students, *****
*** and staff of Yale University only. *****
***-----*****
=====
: The ELI model used by Rutt's CG experiments, modified to have
: multiple register banks per cluster with a partial crossbar.
=====

(alias '^^ 'atomconcat)

(defmacro def-bus ( me-list1 me-list2 number-of-values )
  (let ( (resource-name (atomconcat (car me-list1) '<-> (car me-list2) ) ) )
    (eval-when (eval load)
      (def-resource-class ,resource-name ,number-of-values)
      ..(loop (for me1 in me-list1) (splice
        (loop (for me2 in me-list2) (splice
          '( (def-connection ,me1 ,me2 ( (,resource-name) ) )
            (def-connection ,me2 ,me1 ( (,resource-name))))))))))

(defun em.cluster ( 1 (unique-name unique-fields) )
  (let ( (ci (atomconcat 'c 1) )
        (ci-3 (atomconcat 'c (mod (- 1 3) 8) ) )
        (ci-1 (atomconcat 'c (mod (- 1 1) 8) ) )
        (ci+1 (atomconcat 'c (mod (+ 1 1) 8) ) )
        (ci+3 (atomconcat 'c (mod (+ 1 3) 8) ) ) )
    '(
      ;*** Register bank CiR1
      ;***
      (def-me (name ,(^^ ci 'r1) )
              (type register-bank)
              (size 18)
              (inputs (.(^^ ci-3 'r1)
                       .(^^ ci-3 'r2)
                       .(^^ ci-3 'r3)
                       .(^^ ci-1 'r1)
                       .(^^ ci-1 'r2)
                       .(^^ ci-1 'r3)
                       .(^^ ci 'r1)
                       .(^^ ci 'r2)
                       .(^^ ci 'r3)
                       .(^^ ci+1 'r1)
                       .(^^ ci+1 'r2)
                       .(^^ ci+1 'r3)
                       .(^^ ci+3 'r1)
                       .(^^ ci+3 'r2)
                       .(^^ ci+3 'r3)
                       (^^ ci 'm' ) ) )
              (read-ports 2)
              (write-ports 1)
              (read-resources ( (,^^ ci 'r1-read) ) )
              (write-resources ( (,^^ ci 'r1-write) ) ) )
              (def-resource-class ,(^^ ci 'r1-read) 2)
              (def-resource-class ,(^^ ci 'r1-write) 1)

      ;*** Register bank CiR2
      ;***
      (def-me (name ,(^^ ci 'r2) )

```

```

              (type register-bank)
              (size 18)
              (inputs (.(^^ ci-3 'r1)
                       .(^^ ci-3 'r2)
                       .(^^ ci-3 'r3)
                       .(^^ ci-1 'r1)
                       .(^^ ci-1 'r2)
                       .(^^ ci-1 'r3)
                       .(^^ ci 'r1)
                       .(^^ ci 'r2)
                       .(^^ ci 'r3)
                       .(^^ ci+1 'r1)
                       .(^^ ci+1 'r2)
                       .(^^ ci+1 'r3)
                       .(^^ ci+3 'r1)
                       .(^^ ci+3 'r2)
                       .(^^ ci+3 'r3)
                       .(^^ ci '+)
                       (^^ ci 'c' ) ) )
              (read-ports 2)
              (write-ports 1)
              (read-resources ( (,^^ ci 'r2-read) ) )
              (write-resources ( (,^^ ci 'r2-write) ) ) )
              (def-resource-class ,(^^ ci 'r2-read) 2)
              (def-resource-class ,(^^ ci 'r2-write) 1)

      ;*** Register bank CiR3
      ;***
      (def-me (name ,(^^ ci 'r3) )
              (type register-bank)
              (size 18)
              (inputs (.(^^ ci-3 'r1)
                       .(^^ ci-3 'r2)
                       .(^^ ci-3 'r3)
                       .(^^ ci-1 'r1)
                       .(^^ ci-1 'r2)
                       .(^^ ci-1 'r3)
                       .(^^ ci 'r1)
                       .(^^ ci 'r2)
                       .(^^ ci 'r3)
                       .(^^ ci+1 'r1)
                       .(^^ ci+1 'r2)
                       .(^^ ci+1 'r3)
                       .(^^ ci+3 'r1)
                       .(^^ ci+3 'r2)
                       .(^^ ci+3 'r3)
                       ..(if unique-name
                          '( (,^^ ci unique-name) )
                          ( ) ) )
              (read-ports 2)
              (write-ports 1)
              (read-resources ( (,^^ ci 'r3-read) ) )
              (write-resources ( (,^^ ci 'r3-write) ) ) )
              (def-resource-class ,(^^ ci 'r3-read) 2)
              (def-resource-class ,(^^ ci 'r3-write) 1)

      ;*** Integer adder Ci+
      ;***
      (def-me (name ,(^^ ci '+) )
              (type functional-unit)
              (delay 0)
              (inputs (.(^^ ci 'r2)
                       .(^^ ci 'r3)

```

```

        (resources ( ( (~ ci 'c) ) )
        operators ( ( (~ ci '+) ) ) )
    (float fix inot idiv isub leq isax imin
      iadd ior ige ilt ile ine lexp inul igt
      iand iabs bitrev) )
(def-resource-class ( (~ ci '+) 1)

;*** Integer tester Ci=
;***
(def-me (name ( (~ ci '=) )
  (type functional-unit)
  (delay 0)
  (inputs ( ( (~ ci 'r2)
              ( (~ ci 'r3)
              ( (~ ci 'c) ) )
            (resources ( ( (~ ci '=) ) ) )
            operators (if-true if-false if-ilt if-igt if-ieq if-ine
                        if-ile if-ige if-ieq if-ine if-ile if-ige)))
  (def-resource-class ( (~ ci '=) 1)

;*** Memory CIM
;***
(def-me (name ( (~ ci 'm) )
  (type functional-unit)
  (delay 8)
  (inputs ( ( (~ ci 'r1)
              ( (~ ci 'r2)
              ( (~ ci 'c) ) )
            (resources ( ( (~ ci 'm) ) ) )
            operators (vbase ivload fvload ipload fpload ivstore
                        fvstore ipstore fpstore) ) )
  (def-resource-class ( (~ ci 'm) 1)

;*** Constant generator CiC
;***
(def-me (name ( (~ ci 'c) )
  (type constant-generator)
  (resources ( ( (~ ci 'c) ) ) )
  (constraint-function em.immediate-constant?) )
  (def-resource-class ( (~ ci 'c) 1)

;*** Unique functional-unit
;***
..(if unique-name (then
  '( (def-me (name ( (~ ci unique-name) )
    (type functional-unit)
    (inputs ( ( (~ ci 'r3) ;*** Order important!
                ( (~ ci 'r1)
                ( (~ ci 'c) ) )
    (resources ( ( (~ ci unique-name) ) ) )
    ..unique-fields)
    (def-resource-class ( (~ ci unique-name) 1) ) )
  (else
    ( ) ) )

;*** Bus connections
;***
(def-bus ( ( (~ ci-3 'r1) ( (~ ci-3 'r2) ( (~ ci-3 'r3) )
  ( (~ ci 'r1) ( (~ ci 'r2) ( (~ ci 'r3) )
  1)
  (def-bus ( ( (~ ci-1 'r1) ( (~ ci-1 'r2) ( (~ ci-1 'r3) )
  ( (~ ci 'r1) ( (~ ci 'r2) ( (~ ci 'r3) )

```

```

        ( ) )
        1)
    ) ) )
(= em.floating-test
  '(f=
  ( (delay 0)
    operators (if-flt if-fgt if-feq if-fne if-file if-fige
               flt fgt feq fne file fige) ) ) ) )

(= em.floating-multiply
  '(f*
  ( (delay 3)
    operators (fmul inul) ) ) ) )

(= em.floating-add
  '(f+
  ( (delay 2)
    operators (fsub fadd fmin fmax fabs fdiv cos sin sqrt) ) ) ) )

(defun em.immediate-constant? ( constant )
  (? ( inump constant)
    (& (< constant 2047)
       (> constant -2048) ) )
  (= 0 constant)
  t)
  ( consp constant)
  (= 'address (car constant) ) )
  ( t ) ) )

(defun load-constant? ( constant )
  (if (em.immediate-constant? constant) (then
    (:= 'load *ls.immediate-constant-action*) )
  (else
    t ) ) )

(eval '(def-machine-model
  ..(em.cluster 0 em.floating-test)
  ..(em.cluster 1 em.floating-add)
  ..(em.cluster 2 em.floating-multiply)
  ..(em.cluster 3 em.floating-add)
  ..(em.cluster 4 em.floating-test)
  ..(em.cluster 5 ( ) )
  ..(em.cluster 6 em.floating-multiply)
  ..(em.cluster 7 em.floating-add)
  ) )

```

```

*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for
*** educational and research purposes only, by the faculty, students,
*** and staff of Yale University only.
*****

```

## REGISTERS

This module handles the allocation and deallocation of registers from a register bank.

### (REGISTERS.INITIALIZE)

Initializes all the register banks to be empty.

### (CYCLE:ME:REGISTER:OCCUPIED? CYCLE ME REGISTER)

Returns the VN occupying a register if it is not free, () otherwise.

### (CYCLE:ME:VN:FREE-REGISTER? CYCLE ME VN)

Returns true if a ME has at least one register free at CYCLE or later. Eventually, we'll want to account for the operands of VN that will be freed up this cycle once it is scheduled; but now it's ignored.

### (CYCLE:ME:VN:ALLOCATE-REGISTER CYCLE ME VN REQUIRED-REGISTER)

Allocates and returns a free register from ME for the given CYCLE. The register is marked as holding VN. If non-(), REQUIRED-REGISTER is a specific register that must be allocated for VN. An error is raised if there are no free registers or if REQUIRED-REGISTER is specified but not free.

When choosing a free register, we first look to see if VN is of the form X := f(X) and if X is currently residing in ME. If so, we use X's current register (constraining delays guarantee that the register will be free).

Then we look in the preferred locations of the name of VN (see below) to see if there's a preferred register in bank ME. Then we look for a free register that hasn't been marked as "avoid" (see below). Finally, we look for any register.

### (CYCLE:ME:REGISTER:DEALLOCATE CYCLE ME REGISTER)

Deallocates REGISTER in ME, marking it as free at time CYCLE and later.

### (VN:RECORD-PREFERRED-LOCATION VN)

VN should already have an assigned location. Its location is recorded keyed under the name of VN; when we need to allocate a register for another VN of the same name, we'll prefer registers so recorded.

Also, the location of VN is marked as "avoid", so that we won't allocate it for another VN unless we have to.

### Implementation notes:

The field ME:VALUES is a vector mapping registers onto VNs occupying the registers. ME:CYCLES-FREE gives the first cycle in which the register can be allocated; it is a very large number if the register is not free.

Because CYCLE:ME:VN:FREE-REGISTER? and CYCLE:ME:VN:ALLOCATE are always called with the current cycle during list-scheduling, we know that the

CYCLE argument is monotonically increasing. The ME:REGISTERS-LEFT gives only a hint about how many registers left in the current cycle (it is <= the actual number left in the current cycle). So when we're checking to see if there are any free registers, if ME:REGISTERS-LEFT tells us there are registers left, we can believe it. If ME:REGISTERS-LEFT is 0, then we'll have to recompute it. Gross but relatively simple and efficient.

```

(eval-when (compile load)
  (include list-scheduler:declarations) )

```

```

(declare (special
  *reg.name:locations*      ;** Hash table mapping names to pairs
                           ;** (ME REGISTER) that are to be preferred
                           ;** by the register allocator.
) )

```

```

(defvar *reg.max-integer* 100000)

```

```

(defun registers.initialize ()
  (:= *reg.name:locations* (hash-table:create 'equal () ())) )

```

```

(loop (for me in *ls.register-bank-mes*) (do
  (:= (me:registers-left me) (me:size me) )
  (vector:initialize (me:values me) () )
  (vector:initialize (me:cycles-free me) 0)
  (vector:initialize (me:value-names me) () )
  (vector:initialize (me:avoid? me) () ) ) )

```

```

(defun cycle:me:register:occupied? ( cycle me register )
  (if (< cycle ([]) (me:cycles-free me) register) (then
    ([]) (me:values me) register) )
  (else
    () ) )

```

```

(defun cycle:me:vn:free-register? ( cycle me vn )
  (if (> (me:registers-left me) 0) (then
    t)
  (else
    ;** Recompute the :REGISTERS-LEFT for this cycle. See notes
    ;** above.
    ;**
    (:= (me:registers-left me)
      (loop (incr i from 0 to (- (me:size me) 1) )
        (when (>= cycle ([]) (me:cycles-free me) i) ) )
      (reduce + 0 i) ) )
    (> (me:registers-left me) 0) ) ) )

```

```

(defun cycle:me:vn:allocate-register ( cycle me vn required-register )
  (let* ( (values (me:values me) )
    (avoid? (me:avoid? me) )
    (cycles-free (me:cycles-free me) )
    (free-register

```

```

(|| required-register

  ;*** If VN is of the form X := X op Y and X is
  ;*** currently residing in ME, use its register.
  ;***
  (if (not (= 'operation (vn:type vn) ) (then
    ( )
  )
  (else
    (loop (for operand-vn in (vn:operand-vns vn) )
      (when (&& (= (vn:name vn) (vn:name operand-vn))
        (= me (vn:register-bank-me operand-vn))
      (do
        (let ( (register (vn:register operand-vn) ) )
          (:= ([] cycles-free register) cycle);*** Hack.
          (return register) ) )
        (result ( ) ) ) )

    ;*** Try for a preferred location for the VN.
    ;***
    (loop (for (use-me use-register)
      in ([]h *reg.name:locations* (vn:name vn) ) )
      (when (&& (= me use-me)
        (>= cycle ([] cycles-free use-register))))
    (do
      (return use-register) )
    (result ( ) ) )

    ;*** Try for any free register, avoiding ones marked
    ;*** as "avoid" if possible.
    ;***
    (loop (incr i from 0 to (- (me:size me) 1) )
      (when (>= cycle ([] cycles-free i) ) )
      (initial avoid-register ( ) )
    (do
      (if (not ([] avoid? i) ) (then
        (return i) ) )
      (if (not avoid-register) (then
        (:= avoid-register i) ) ) )
    (result avoid-register) )

    0) ) )

  (if (< cycle ([] cycles-free free-register) ) (then
    (error (list me "No available register." ) ) )

  (:= ([] values free-register) vn)
  (:= ([] cycles-free free-register) *reg.max-integer*)
  (:= (me:registers-left me) (- &&& 1) )

  free-register) )

(defun cycle:me:register:deallocate ( cycle me register )
  (:= ([] (me:cycles-free me) register) cycle)
  ;***
  ;*** Sic -- we don't touch :REGISTERS-FREE, since we could be
  ;*** freeing up a register "in the future". See above.
  ( ) )

(defun vn:record-preferred-location ( vn )
  (assert (vn:register vn) )

```

```

(let ( (me (vn:register-bank-me vn) )
  (register (vn:register vn) ) )
  (:= ([] (me:avoid? me) register) t)
  (push ([]h *reg.name:locations* (vn:name vn) ) '(me .register) )
  ( ) )

```

```

***-----*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for *****
*** educational and research purposes only, by the faculty, students, *****
*** and staff of Yale University only. *****
***-----*****

=====
RESOURCES

This module implements the operations on resources used in the machine
model. It also implements a single (global) schedule of resources.

A RESOURCE-CLASS defines a set of resources that are all identical as
far as the machine model is concerned. Any resource in the set can
be used in place of any other.

A RESOURCE-SET is a set of unused resources in the machine (for one
cycle). RESOURCE-SETS are represented as bit-sets.

A RESOURCE-REQUEST is a request for particular resources across
successive cycles. It is represented as a list of lists of
RESOURCE-CLASSES, each sublist representing the resources requested
for a cycle (the first sublist is for the first cycle, the second sublist
for the second cycle...)

(RESOURCE.INITIALIZE)
  Initializes this module.

(RESOURCE.INITIALIZE-SCHEDULE)
  Re-initializes the schedule of used resources to be all available.

(RESOURCE-CLASS:CREATE NAME SIZE)
  Creates a resource class with symbolic NAME and SIZE resources.

(DEF-RESOURCE-CLASS 'NAME 'SIZE)
  Same as above, except the arguments aren't evaluated.

(NAME:RESOURCE-CLASS NAME)
  Maps a symbolic name onto the corresponding RESOURCE-CLASS, raising an
  error if there is no such class.

(RESOURCE-REQUEST:INstantiate RESOURCE-REQUEST)
  Takes a resource request consisting of symbolic names of
  RESOURCE-CLASSES and returns a new proper request with the names
  replaced by the corresponding RESOURCE-CLASSES.

(CYCLE:RESOURCE-REQUEST:AVAILABLE? CYCLE RESOURCE-REQUEST)
  True if RESOURCE-REQUEST can be scheduled in the given cycle.

(CYCLE:RESOURCE-REQUEST:FIRST-AVAILABLE-CYCLE CYCLE RESOURCE-REQUEST)
  Returns the first cycle after CYCLE (inclusive) in which the resource
  request can be scheduled.

(CYCLE:RESOURCE-REQUEST:SCHEDULE? CYCLE RESOURCE-REQUEST)
  Tries to schedule the resources of RESOURCE-REQUEST beginning at
  CYCLE, returning true if successful, false o.w. If false, the
  schedule is left in garbaged state; use transactions below to restore
  the state correctly.

(CYCLE:RESOURCE-REQUEST:SCHEDULE CYCLE RESOURCE-REQUEST)
  Same as above, but raises an error if the request can't be scheduled.

(RESOURCE.START-TRANSACTION)

```

```

(RESOURCE.COMMIT-TRANSACTION)
(RESOURCE.ABORT-TRANSACTION)
  These three procedures let the client tentatively schedule resource
  requests, but then undo them if he wants. Such a tentative attempt
  is called a "transaction"; only one transaction may be in progress
  at a time. After calling RESOURCE.START-TRANSACTION, all calls
  to CYCLE:RESOURCE-REQUEST:SCHEDULE? and
  CYCLE:RESOURCE-REQUEST:SCHEDULE will "save away" the old values of
  the scheduled resources. Calling RESOURCE.ABORT-TRANSACTION will
  restore the resource schedule back to its state at the time of the
  call to RESOURCE.START-TRANSACTION. Calling
  RESOURCE.COMMIT-TRANSACTION signifies that all the resources scheduled
  are "ok" and to forget the "saved away" old schedule.

(RESOURCE.PRINT-SCHEDULE)
  Prints the current resource schedule in a pretty form.

=====

*****
***
*** RESOURCE-CLASS
***
*****
(def-struct resource-class ***
  name *** The printed name of the class.
  size *** The number of resources in the class.
  position *** Position in the bytevector.
  ls-failures ***
  bug-failures *** The sum of the sizes of requests for
  ) *** unavailable resources from this class;
  *** separate totals for list-scheduling and
  *** BUG.
*****
*****

(visible-fields resource-class name size)

(eval-when (compile load)
  (include list-scheduler:declarations) )
(eval-when (compile)
  (build '(list-scheduler:byte-vector) ) )

(declare (special
  *res.all-resource-classes*
  *res.name:resource-class*
  *res.total-classes*
  *res.max-class-size*
  *res.cycle:resource-set*
  *res.transaction-in-progress*
  *res.transaction-size*
  *res.saved-bytes*
  *res.saved-positions*
  *res.saved-cycles*
  *res.failure-type*
  ) )

(defun resource.initialize (
  (: *res.all-resource-classes* ( ) )
  (: *res.total-classes* 0)
  (: *res.max-class-size* 0)

```

```

(= *res.name:resource-class* (hash-table:create () () () ) )
(= *res.cycle:resource-set* () )
(= *res.transaction-in-progress?* () )
(= *res.transaction-size* 0 )
(= *res.saved-bytes* () )
(= *res.saved-positions* () )
(= *res.saved-cycles* () )
(= *res.failure-type* 'ls)
() )

(defun resource.initialize-schedule ( failure-type )
  (:= *res.cycle:resource-set* (makevector *ls.max-schedule-size* ) )
  (let ( (byte-size (ceiling (/ (log (+ 1 *res.max-class-size*)
                                (log 2) ) ) ) )
        (loop (incr 1 from 0 to (- *ls.max-schedule-size* 1) ) (do
          (:= ( [] *res.cycle:resource-set* 1)
              (byte-vector:create *res.total-classes* byte-size))))))

    (:= *res.transaction-in-progress?* () )
    (:= *res.transaction-size* 0)
    (:= *res.saved-bytes* (makevector 100) ) :*** Big enough?
    (:= *res.saved-positions* (makevector 100) ) :*** Big enough?
    (:= *res.saved-cycles* (makevector 100) ) :*** Big enough?

    (assert (memq failure-type '(ls bug) ) )
    (:= *res.failure-type* failure-type)
    () )

(defun resource-class:create ( name size )
  (let ( (resource-class
        (resource-class:new
         name name
         size size
         position *res.total-classes*
         ls-failures 0
         bug-failures 0) ) )

    (:= *res.total-classes* (+ &@& 1) )
    (:= *res.max-class-size* (max &@& size) )

    (:= *res.all-resource-classes* (appendi &@& resource-class) )
    (:= ( []h *res.name:resource-class* name) resource-class)

    resource-class) )

(defmacro def-resource-class ( name size )
  '(resource-class:create ',name ',size) )

(defun name:resource-class ( name )
  (|| ( []h *res.name:resource-class* name)
      (error (list name " isn't a RESOURCE-CLASS." ) ) ) )

(defun resource-request:instantiate ( resource-request )
  (loop (for cycle-request in resource-request) (save
    (loop (for item in cycle-request) (save
      (if (consp item)
          '(.(name:resource-class (car item) ) ,(cadr item) )
          '(.(name:resource-class item) 1) ) ) ) ) ) ) )

```

3

```

(defun resource.start-transaction ()
  (assert (! *res.transaction-in-progress?*) )
  (:= *res.transaction-in-progress?* t)
  (:= *res.transaction-size* 0)
  () )

(defun resource.commit-transaction ()
  (assert *res.transaction-in-progress?*)
  (:= *res.transaction-in-progress?* () )
  (:= *res.transaction-size* 0)
  () )

(defun resource.abort-transaction ()
  (assert *res.transaction-in-progress?*)
  (loop (decr 1 from (- *res.transaction-size* 1) to 0) (do
    (:= ( []b ( [] *res.cycle:resource-set* ( [] *res.saved-cycles* 1) )
          ( [] *res.saved-positions* 1) )
        ( [] *res.saved-bytes* 1) ) ) )

  (:= *res.transaction-in-progress?* () )
  (:= *res.transaction-size* 0)
  () )

(defun cycle:resource-request:available? ( cycle resource-request )
  (loop (incr 1 from cycle)
    (for cycle-request in resource-request)
    (when (! (resource-set:cycle-request:available?
              ( [] *res.cycle:resource-set* 1)
              cycle-request) ) )

    (do
      (return () ) )
    (result t) ) )

(defun cycle:resource-request:first-available-cycle
  ( first-cycle resource-request )
  (loop (incr cycle from first-cycle)
    (until (cycle:resource-request:available?
          cycle
          resource-request) )

    (result cycle) ) )

(defun cycle:resource-request:schedule? ( cycle resource-request )
  (loop (incr 1 from cycle)
    (for cycle-request in resource-request)
    (when (! (cycle:cycle-request:schedule?
              1
              cycle-request) ) )

    (do
      (return () ) )
    (result t) ) )

(defun cycle:resource-request:schedule ( cycle resource-request )
  (if (! (cycle:resource-request:schedule? cycle resource-request) ) (then
    (error (list cycle resource-request
      "CYCLE:RESOURCE-REQUEST:SCHEDULE: Couldn't schedule."))))

```

4

```

() )

(defun resource-set:cycle-request:available? ( resource-set cycle-request )
  (loop (for (resource-class request-size) in cycle-request)
    (when (> request-size
      (- (resource-class:size resource-class)
        ([[]b resource-set
          (resource-class:position resource-class))))))
  (do
    (if (== 'ls *res.failure-type*) (then
      (:= (resource-class:ls-failures resource-class) (+ &&& 1) ) )
      (else
        (:= (resource-class:bug-failures resource-class) (+ &&& 1) ) ) )
    (return () ) )
  (result t) ) )

(defun cycle:cycle-request:schedule? ( cycle cycle-request )
  (let ( (resource-set ([[] *res.cycle:resource-set* cycle) ) )
    (loop (for (resource-class request-size) in cycle-request)
      (bind old-resources
        ([[]b resource-set
          (resource-class:position resource-class) )
        new-resources
          (+ request-size old-resources) )
      (do
        (if (> new-resources (resource-class:size resource-class) ) (then
          (if (== 'ls *res.failure-type*) (then
            (:= (resource-class:ls-failures resource-class) (+ &&& 1) ) )
            (else
              (:= (resource-class:bug-failures resource-class) (+ &&& 1) ) ) )
          (return () ) ) )
        (if *res.transaction-in-progress* (then
          (:= ([[] *res.saved-cycles* *res.transaction-size*
            cycle)
            ([[] *res.saved-bytes* *res.transaction-size*
            old-resources)
            ([[] *res.saved-positions* *res.transaction-size*
            (resource-class:position resource-class) )
            (:= *res.transaction-size* (+ &&& 1) ) ) )
          (:= ([[]b resource-set (resource-class:position resource-class) )
            new-resources)
          (result t) ) ) )
    (defun resource.print-schedule ()
      (msg 0)
      (loop (initial first-1 0)
        (while (< first-1 *ls.max-schedule-size*)
          (bind first-set ([[] *res.cycle:resource-set* first-1) )
            (do
              (loop (incr last-1 from (+ first-1 1) to (- *ls.max-schedule-size* 1))
                (while (resource-set:equal first-set
                  ([[] *res.cycle:resource-set* last-1)))
                (result
                  (:= last-1 (- last-1 1) )
                  (if (> last-1 first-1) (then
                    (msg (j first-1 3) ":( (j last-1 3) " " )

```

```

      (else
        (msg (j first-1 3) " " " ) )
      (resource-set:print first-set)
      (msg t)
      (:= first-1 (+ last-1 1) ) ) ) ) )
  (do
    (return () ) )
  (result t) ) )

(defun resource-set:equal ( resource-set1 resource-set2 )
  (loop (incr i from 0 to (- *res.total-classes* 1) )
    (when (!= ([[]b resource-set1 i)
      ([[]b resource-set2 i) ) )
    (do
      (return () ) )
    (result t) ) ) )

(defun resource-set:print ( resource-set )
  (sprin1
    (*hprint.structure-marker* (resource-set
      ..(loop (for resource-class in *res.all-resource-classes*
        (bind resources
          ([[]b resource-set
            (resource-class:position resource-class))))
      (splice
        (? (== 0 resources)
          () )
        ( (== 1 (resource-class:size resource-class) )
          (assert (== 1 resources) )
          '( (resource-class:name resource-class) ) )
        ( t
          '( ( (resource-class:name resource-class)
            .resources) ) ) ) ) ) ) ) ) )

(defun resource.print-failures ()
  (msg 0 t "Resource failures:" (t 20) (jc "LS" 5) (t 30) (jc "BUG" 5) t)
  (loop (for resource-class
    in (sort *res.all-resource-classes*
      (f:1 ( r1 r2 )
        (> (resource-class:ls-failures r1)
          (resource-class:ls-failures r2) ) ) )
    (when (|| (< 0 (resource-class:ls-failures resource-class) )
      (< 0 (resource-class:bug-failures resource-class) ) ) )
    (do
      (msg (resource-class:name resource-class)
        (t 20) (j (resource-class:ls-failures resource-class) 5)
        (t 30) (j (resource-class:bug-failures resource-class) 5) t)))
  (do
    (return () ) )

```

```

*****
*** Copyright (C) 1983 John R. Ellis. This file may be used for
*** educational and research purposes only, by the faculty, students,
*** and staff of Yale University only.
*****
(eval-when (compile load)
  (include list-scheduler:declarations) )

*****
***
*** (SCH.COPY-VN:SCHEDULE VN)
***
*** Attempts to schedule VN, a COPY, requeusing it if it can't be scheduled
*** because of resource conflict.
***
*** If the copy has no readers left, that means its only use was an
*** off-live use at a split which has already been scheduled above the
*** copy; the copy is now useless and can be scheduled as a NOOP.
***
*****

(defun sch.copy-vn:schedule ( vn )
  (if (== 0 (vn:readers-left vn) ) (then
    (sch.noop-vn:schedule vn) )
    (else
      (resource.start-transaction)
      (caseq (sch.copy-vn:schedule? vn)
        (success
          (resource.commit-transaction) )
        (request
          (resource.abort-transaction)
          (sch.vn:cycle:enqueue vn (+ *ls.cycle* 1) ) )
        (abort
          (resource.abort-transaction) )
        (t
          (error "Case error.") ) )
      ) ) ) )

(defun sch.copy-vn:schedule? ( vn ) (prog ()

  (let* ( (operand-vn (car (vn:operand-vns vn) ) )
    (operand-me (sch.vn:location-me operand-vn) )
    (dest-me ( ) )
    (dest-register ( ) )
    (reading-vn (sch.vn:max-height-reading-vn vn) )
    (reading-me (&& reading-vn
      (bug.vn:dest-me reading-vn) ) ) )

    *** Try scheduling the resources needed to read the operand.
    ***
    (if (! (cycle:resource-request:schedule?
      *ls.cycle*
      (if (== 'constant-generator (me:type operand-me) )
        (me:resources operand-me)
        (me:read-resources operand-me) ) ) ) )
      (then
        (return 'requeue) ) )

    *** Find a good destination register bank, and try scheduling
    *** the write into it.
    ***

```

```

(= dest-me (sch.source-me:dest-me:vn:pick-register-bank?
  operand-me
  reading-me
  vn) )
(if (! dest-me) (then
  (return 'requeue) ) )

*** At this point, all the resources of the copy have been
*** scheduled.
***
*** If the max height reader is a USE1 that wants a particular
*** register, and that register is occupied, we'll have to splice
*** a COPY between the occupant of that register and the readers
*** of the occupant, making this VN dependent on completion of
*** that newly spliced COPY. Hopefully, after that COPY is
*** scheduled, we'll have a free register to do this COPY.
***
(if (&& reading-vn
  (== 'use1 (vn:type reading-vn) )
  (== dest-me (vn:register-bank-me reading-vn) ) )
  (then
    (if-let ( (occupant-vn (cycle:me:register:occupied?
      *ls.cycle*
      dest-me
      (vn:register reading-vn) ) ) )
      (then
        (return (sch.copy-vn:blocking-vn:make-dependent vn occupant-vn))))
    (:= dest-register
      (cycle:me:vn:allocate-register
        *ls.cycle* dest-me vn (vn:register reading-vn) ) ) )
  (else
    (:= dest-register
      (cycle:me:vn:allocate-register *ls.cycle* dest-me vn ( ) ) ) ) ) )
(:= (vn:scheduled-cycle vn) *ls.cycle*)

*** Mark the operand as having been read this cycle.
***
(:= (vn:read-cycle operand-vn) *ls.cycle*)

*** Assign the register location to this VN, inserting
*** copy VNs as necessary, and decrementing the predecessor
*** counts of readers and constrained VNs; decrement the
*** read counts of the operand just read.
***
(sch.vn:me:register:cycle:assign-location
  vn
  dest-me
  dest-register
  (+ *ls.cycle* 1) )

(sch.vn:release-operands vn)
(sch.vn:release-successors vn)
(sch.vn:release-off-live vn)

*** Schedule the machine operation for the COPY.
***
(sch.copy-vn:schedule-machine-operation vn)

(return 'success) ) ) )

```



```

=====
***
*** (SCH.COPY-VN:SCHEDULE-MACHINE-OPERATION VN)
***
*** Makes a machine operation for VN, a COPY, and places it on the
*** schedule.
***
=====
(defun sch.copy-vn:schedule-machine-operation ( vn )
  (let* ( (operand-vn (car (vn:operand-vns vn) ) )
         (source-me (sch.vn:location-me operand-vn) ) )

    (push ([ *ls.schedule* *ls.cycle* )
          (.vn
            (.me:name (vn:register-bank-me vn) )
            copy
            (.me:name (vn:register-bank-me vn) )
            (.vn:register vn)
            (.vn:name vn) )
          (if (== 'constant-generator (me:type source-me) )
              (.me:name source-me)
              (.vn:name operand-vn) )
          (.me:name source-me)
          (.vn:register operand-vn)
          (.vn:name operand-vn) ) ) )
    (:= *ls.last-cycle* (max &&& *ls.cycle* ) )
    ( ) )

=====
***
*** (SCH.COPY-VN:BLOCKING-VN:MAKE-DEPENDENT VN BLOCKING-VN)
***
*** VN is a COPY that, due to a USE1, wants to put its result in a
*** register currently occupied by BLOCKING-VN; VN can't be scheduled
*** until the BLOCKING-VN's value is moved somewhere else. This procedure
*** adds in constraint edges, and possibly new COPIES, to insure that
*** BLOCKING-VN's value will be moved before scheduling of VN is attempted
*** again.
***
*** If all of BLOCKING-VN's unscheduled readers do not have VN as an
*** ancestor, then we can just wait until all of those readers are
*** scheduled, at which time the register needed by VN will be free.
*** This is accomplished by adding a new constraint edge between VN and
*** the readers.
***
*** But if one of BLOCKING-VN's readers has VN as an ancestor, then adding
*** a constraint edge will produce deadlock (the reader can't be scheduled
*** until VN is, which can't be scheduled until the reader is). So we
*** splice a new COPY between BLOCKING-VN and all of its unscheduled
*** readers, constraining VN to be scheduled after the new COPY (which
*** will have freed up the desired register). The new COPY is enqueued
*** for the current cycle.
***
*** This procedure returns 'ABORT or 'REQUEUE depending on whether VN
*** should be queued for the next cycle (usually it shouldn't -- see
*** below).
***
=====
(defun sch.copy-vn:blocking-vn:make-dependent ( vn blocking-vn )

```

```

(let* ( (unscheduled-reading-vns
        (loop (for reading-vn in (vn:reading-vns blocking-vn) )
              (when (! (vn:scheduled-cycle reading-vn) ) )
              (save reading-vn) ) )

        (unscheduled-off-live-reading-vns
        (loop (for reading-vn in (vn:off-live-reading-vns blocking-vn) )
              (when (! (vn:scheduled-cycle reading-vn) ) )
              (save reading-vn) ) )

        (all-unscheduled-reading-vns
        (append unscheduled-off-live-reading-vns unscheduled-reading-vns)))

  (if (for-every (reading-vn in all-unscheduled-reading-vns)
              (! (sch.vn:vn:descendant? vn reading-vn) ) )
      (then
        :***
        :*** The unscheduled readers of BLOCKING-VN don't have VN
        :*** as an ancestor. We make VN dependent on completion of
        :*** all these readers.

        (loop (for reading-vn in all-unscheduled-reading-vns)
              (when (|| (== 'copy (vn:type reading-vn) )
                      (== 'operation (vn:type reading-vn) ) ) )
              (initial found-one? ( ) )

              (do
                (push (vn:constraining-vns vn) reading-vn)
                (push (vn:constraining-delays vn) 0)
                (:= (vn:predecessors-left vn) (+ &&& 1) )

                (push (vn:constrained-vns reading-vn) vn)
                (:= found-one? t) )

              (result
                (if found-one?
                    'abort
                    'requeue) ) ) )

        (else :***
              :*** The unscheduled readers are dependent on VN. We've got
              :*** to splice a new COPY between BLOCKING-VN and the readers,
              :*** making VN dependent on the COPY.

              (let ( (new-copy-vn
                    (vn:create (vn:new
                              type 'copy
                              name (vn:name blocking-vn)
                              height (vn:height blocking-vn)
                              cycle (vn:register-cycle blocking-vn)
                              readers-left (length all-unscheduled-reading-vns))))

                    (vn:splice-vn blocking-vn
                                  new-copy-vn
                                  unscheduled-reading-vns
                                  unscheduled-off-live-reading-vns)

                    (:= (vn:readers-left blocking-vn)
                        (+ 1 (- &&& (vn:readers-left new-copy-vn) ) ) )

```

```

(loop (for reading-vn in unscheduled-reading-vns) (do
      ;*** ~ (sic)
      (if (not= reading-vn vn) (then
        (sch.vn:dequeue reading-vn) ) )
      (:= (vn:predecessors-left reading-vn) (+ &&& 1) ) ) )

(push (vn:constraining-vns vn) new-copy-vn)
(push (vn:constraining-delays vn) 0)
(:= (vn:predecessors-left vn) (+ &&& 1) )

(push (vn:constrained-vns new-copy-vn) vn)

(sch.vn:cycle:enqueue new-copy-vn *1s.cycle*)
'abort) ) ) )

;***=====
;***
;*** (SCH.VN:VN:DESCENDANT? VN DESCENDANT-VN)
;***
;*** Returns true if DESCENDANT-VN really is a descendent of VN. This is
;*** implemented by recursing back up through all the predecessors of
;*** DESCENDANT-VN until either we run out of predecessors, we hit a scheduled
;*** VN, or we hit VN itself. We know we can stop recursing when we hit a
;*** scheduled predecessor, since VN (and all its descendants) are unscheduled.
;***
;*** We might have to change this implementation to use bit sets or something
;*** as in common subexpression elimination to make it linear; we'll see
;*** how expensive it is.
;***
;***=====

(defun sch.vn:vn:descendant? ( vn descendant-vn )
  (? ( (= vn descendant-vn)
      t)
    ( vn:scheduled-cycle descendant-vn)
    () )
  ( t
    (|| (for-some (operand-vn in (vn:operand-vns descendant-vn) )
      (sch.vn:vn:descendant? vn operand-vn) )
      (for-some (constraining-vn in (vn:constraining-vns descendant-vn))
      (sch.vn:vn:descendant? vn constraining-vn) ) ) ) ) )

```

```

*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for
*** educational and research purposes only, by the faculty, students,
*** and staff of Yale University only.
*****
(eval-when (compile load)
  (include list-scheduler:declarations) )

=====
***
*** (SCH.OPERATION-VN:SCHEDULE VN)
***
*** Attempts to schedule VN, an operation, requeuing it if it can't
*** be scheduled for any reason.
***
*** If the operation has no readers left, that means its only use was
*** an off-live use at a split which has already been scheduled above
*** the operation; the operation is now useless and can be scheduled
*** as a NOOP.
***
=====

(defun sch.operation-vn:schedule ( vn )
  (if (&& (= 0 (vn:readers-left vn)
              (vn:name vn) )
      (then
        (sch.noop-vn:schedule vn) )
      (else
        (resource.start-transaction)
        (if (sch.operation-vn:schedule? vn) (then
          (resource.commit-transaction) )
          (else
            (resource.abort-transaction)
            (sch.vn:cycle:enqueue vn (+ *ls.cycle* 1) ) ) )
        ( ) ) ) )

(defun sch.operation-vn:schedule? ( vn ) (prog ( )

(let ( (fu-me (vn:me vn) )
      (unique-operand-vns (nodupseq (vn:operand-vns vn) ) )
      (dest-me ( ) )
      (dest-register ( ) ) )

  ;*** Try scheduling the functional unit resources.
  ;***
  (if (! (cycle:resource-request:schedule? *ls.cycle* (me:resources fu-me)))
      (then
        (return ( ) ) )
      ;*** Try scheduling the resources needed to read the operands.
      ;***
      (if (! (for-every (operand-vn in unique-operand-vns)
        (sch.operation-operand-vn:fu-me:schedule-read?
          operand-vn fu-me) ) )
          (then
            (return ( ) ) )
          ;*** Find a good destination register bank and try to schedule
          ;*** the write of the result.
          ;***
          (if (vn:name vn) (then

```

```

      (:= dest-me (sch.source-me:dest-me:vn:pick-register-bank?
                  fu-me
                  (sch.vn:max-height-dest-me vn)
                  vn) )
      (if (! dest-me) (then
        (return ( ) ) ) ) ) )

;*** At this point, we know we can schedule the operation; all the
;*** resources have already been scheduled.
;***
;*** Decrement the read counts of the operands, possibly freeing
;*** up their registers.
;***
(:= (vn:scheduled-cycle vn) *ls.cycle*)

(sch.vn:release-operands vn)

;*** Allocate a register for VN.
;***
(if dest-me (then
  (:= dest-register
    (cycle:me:vn:allocate-register *ls.cycle* dest-me vn ( ) ) ) ) )

;*** Mark the operands as having been read this cycle.
;***
(loop (for operand-vn in unique-operand-vns) (do
  (:= (vn:read-cycle operand-vn) *ls.cycle*) ) )

;*** Assign the register location to this VN, inserting
;*** copy VNs as necessary, and decrementing the predecessor
;*** counts of readers and constrained VNs.
;***
(if dest-me (then
  (sch.vn:me:register:cycle:assign-location
    vn
    dest-me
    dest-register
    (+ *ls.cycle* (+ (me:delay fu-me) 1) ) ) ) )

(sch.vn:release-successors vn)
(sch.vn:release-off-live vn)

;*** Place a machine operation for VN on the schedule.
;***
(sch.operation-vn:schedule-machine-operation vn)

(return t) ) ) )

=====
***
*** (SCH.OPERATION-OPERAND-VN:FU-ME:SCHEDULE-READ? VN FU-ME)
***
*** VN is an operand of an operation, and FU-ME is the functional unit
*** where that operation will be computed. This procedure attempts to
*** schedule the resources necessary to read the value of VN and deliver
*** it to the input of FU-ME. There are two independent considerations:
***
*** 1. The value of VN may be in either a register bank or a constant
*** generator.
***
*** 2. The value of VN may have already been read this cycle, so
*** we don't need to schedule the resources for the register bank
*** or constant generator.
***

```

```

***
*** Returns true if VN can be read this cycle, false otherwise. As a
*** side effect, the resources for reading VN and moving it to FU-ME
*** are scheduled.
***
=====
(defun sch.operation-operand-vn:fu-me:schedule-read? ( vn fu-me ) (prog ()
  (let ( (me (sch.vn:location-me vn) ) )
    (if (&& (!= *ls.cycle* (vn:read-cycle vn) )
      (! (cycle:resource-request:schedule?
        *ls.cycle*
        (if (== 'constant-generator (me:type me) )
            (me:resources me)
            (me:read-resources me) ) ) ) )
      (then
        (return () ) ) )
    (if (! (cycle:resource-request:schedule?
      *ls.cycle*
      (me:me:resources me fu-me) ) )
      (then
        (return () ) ) )
    (return t) ) ) )

***
=====
*** (SCH.SOURCE-ME:DEST-ME:VN:PICK-REGISTER-BANK? SOURCE-ME DEST-ME VN)
***
*** Attempts to find a "good" register bank connected to the outputs
*** of SOURCE-ME to which we can move a value from SOURCE-ME. The
*** register-bank ME is returned if found, otherwise () is returned.
*** As a side effect, the resources needed to write the result into the
*** register bank are scheduled.
***
*** DEST-ME is the next destination of the value after it is stored in
*** the register bank, and VN is the COPY or OPERATION VN which needs
*** the register bank.
***
*** To find a good register bank, we look at ones that:
***
***   - are connected to the output of the SOURCE-ME;
***   - are on the shortest path to DEST-ME;
***   - have a free register;
***   - can have the result moved to them from the functional unit and
***     written (i.e. there are free resources);
***
*** Of these register banks, we find the ones that will cause minimum
*** conflicts for the readers of VN. If we still have a choice, we'd
*** prefer a register bank that is DEST-ME itself (in case DEST-ME is
*** a register bank, e.g. for a USE1).
***
=====
(defun sch.source-me:dest-me:vn:pick-register-bank?
  ( source-me dest-me vn )

  (let*( (write-cycle (+ *ls.cycle*
    (if (== 'functional-unit (me:type source-me))
        (me:delay source-me)

```

```

    ) ) )
  (source-dest-delay (me:me:delay source-me dest-me) )
  (available-mes
    (loop (for me in (me:outputs source-me) )
      (when (== 'register-bank (me:type me) ) )
      (when (|| (! dest-me)
        (== me dest-me)
        (< (me:me:delay me dest-me)
          source-dest-delay)
        (&& (== 0 source-dest-delay)
          (== 0 (me:me:delay me dest-me) )
          (== 'functional-unit
            (me:type dest-me) ) ) ) )
        ;***
        ;*** We allow a "sideways" move --
        ;*** a move that gets us no closer
        ;*** to the dest -- only if the dest
        ;*** is a FU and the the value is
        ;*** already in a register connected
        ;*** to the FU.

        (when (cycle:me:vn:free-register? *ls.cycle* me vn) )
        (when (cycle:resource-request:available?
          write-cycle
          (me:me:resources source-me me) ) )
        (when (cycle:resource-request:available?
          write-cycle
          (me:write-resources me) ) )
        (save-minimums me (sch.vn:me:reading-vn-conflicts vn me))))

  (register-bank-me
    (if (memq dest-me available-mes)
        dest-me
        (car available-mes) ) ) )

  (if (! register-bank-me) (then
    () )
    (else
      (cycle:resource-request:schedule
        write-cycle
        (me:me:resources source-me register-bank-me) )
      (cycle:resource-request:schedule
        write-cycle
        (me:write-resources register-bank-me) )
      register-bank-me) ) )

***
=====
*** (SCH.OPERATION-VN:SCHEDULE-MACHINE-OPERATION VN)
***
*** Places the machine operation for VN, an operation, on the schedule.
***
=====
(defun sch.operation-vn:schedule-machine-operation ( vn )
  (let ( (last-cycle
    (+ *ls.cycle* (me:delay (vn:me vn) ) ) )
    (operands
      (loop (for operand-vn in (vn:operand-vns vn) )
        (bind operand-me (sch.vn:location-me operand-vn) )
        (save

```

```

      (if (== 'constant-generator (me:type operand-me) ) (then
        '( (me:name operand-me)
          (vn:name operand-vn) ) )
        (else
          '( (me:name operand-me)
            (vn:register operand-vn)
            (vn:name operand-vn) ) ) ) ) ) )
(if (vn:register-bank-me vn) (then
  (push operands
    '( (me:name (vn:register-bank-me vn) )
      (vn:register vn)
      (vn:name vn) ) ) )
(push ( [] *ls.schedule* *ls.cycle* )
  '(vn ,(me:name (vn:me vn) ) ,(vn:operator vn) ..operands) )
(loop (incr cycle from (+ *ls.cycle* 1) to last-cycle) (do
  (push ( [] *ls.schedule* cycle)
    '(vn ,(me:name (vn:me vn) ) ) ) ) )
(:= *ls.last-cycle* (max &&& last-cycle) )
() )

```

```

*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for
*** educational and research purposes only, by the faculty, students,
*** and staff of Yale University only.
*****

```

## SPLITS AND JOINS

This module implements the part of the list-scheduler dealing with splits and joins, doing the work of SCHEDULE:SPLIT and SCHEDULE:JOIN.

To generate the proper DEFs and USEs at splits and joins, we compute of the extents (lifetimes) of VNs. That is, for each VN, we record the first cycle in which its value becomes available and the last cycle its value is used. We keep a vector that gives, for each cycle, the set of VNs that are live on entry to that cycle (whose extents include that cycle).

Partial schedules add complications. When we ask for the set of live VNs at a join, what we're really interested in is the set of live VNs on entry to the partial schedule that will be generated to precede the join. Similarly, when we ask for the set of live VNs at a split, we're really interested in the VNs that are live on exit from the partial schedule that will be inserted in the split. So we actually compute two different extent vectors, one for joins and one for splits.

For joins, a VN doesn't become live until the cycle after the cycle that wrote the value (the cycle after the end of the operation); the VN's value becomes dead at the end of the maximum of:

1. The maximum, over all readers, of the last cycle of the reader.
2. The maximum, over all off-live readers (which are jumps), of the last cycle of the jump.

For splits, a VN becomes live on entry to the cycle after the cycle in which the VN is scheduled (the cycle its operation started); the VN's value becomes dead at the end of the cycle which is the maximum of:

1. The maximum, over all readers, of the first cycle of the reader.
2. The maximum, over all off-live readers (which are jumps), of the cycle after the end of the jump.

The second computation accounts for the fact that a value that is off-live at a split remains alive until after the jump has executed, and should be reported alive at the split.

My dissertation will have to give good clear explanations of why these properly report the live VNs at splits and joins.

To efficiently compute the live extents of VNs, we use two auxiliary vectors: GEN gives for each cycle the set of VNs whose values first become live at the beginning of the cycle; and KILL gives for each cycle the set of VNs whose values become dead after the end of the cycle. The LIVE vector is computed as:

```
LIVE[ I ] := (LIVE[ I-1 ] - KILL[ I-1 ] ) + GEN[ I ]
```

```

(=eval-when (compile load)
  (include list-scheduler:declarations) )

```

```

(declare (special
  *tr.trace-number*      ;*** For debugging only.

```

```

  *sch.cycle:split-live*
  *sch.cycle:join-live*
) )

```

```

*****
***
*** (SCH.COMPUTE-LIVE-DEAD)
***
*** Computes the live-dead information for splits and joins as described
*** above. This procedure should be called after all the VNs have been
*** scheduled.
***
*****

```

```

(defun sch.compute-live-dead ()
  (:= *sch.cycle:split-live*
    (sch.gen-function:kill-function:compute-live-vector
      'sch.vn:split-gen-cycle
      'sch.vn:split-kill-cycle) )
  (:= *sch.cycle:join-live*
    (sch.gen-function:kill-function:compute-live-vector
      'sch.vn:join-gen-cycle
      'sch.vn:join-kill-cycle) )
  ( ) )

```

```

*****
***
*** (SCH.GEN-FUNCTION:KILL-FUNCTION:COMPUTE-LIVE-VECTOR
***   LIVE GEN-FUNCTION KILL-FUNCTION)
***
*** This function computes the LIVE vector, as described above.
*** GEN-FUNCTION and KILL-FUNCTION take a single argument, a VN, and
*** return the first cycle that VN is live or the last cycle it is live.
*** The new live vector is returned. The GEN-FUNCTION returns () if
*** the VN doesn't produce a value.
***
*****

```

```

(defun sch.gen-function:kill-function:compute-live-vector
  ( gen-function kill-function )

  (let ( (live (makevector (+ *ls.last-cycle* 8) ))
        (gen (makevector (+ *ls.last-cycle* 8) ))
        (kill (makevector (+ *ls.last-cycle* 8) )) )
    (vector:initialize live *vn-set.empty-set*)
    (vector:initialize gen *vn-set.empty-set*)
    (vector:initialize kill *vn-set.empty-set*)

    ;*** Set GEN and KILL for each VN.
    ;***
    (loop (for-each-vn vn)
      (bind gen-cycle (funcall gen-function vn) )
      (when gen-cycle)
      (bind kill-cycle (funcall kill-function vn) )

    (do
      (:= ( [] gen gen-cycle) (vn-set:union1 &&& vn) )
      (:= ( [] kill kill-cycle) (vn-set:union1 &&& vn) ) ) )

```

```

*** Compute LIVE from GEN and KILL
***
(= ( [] live 0) ( [] gen 0) )
(loop (incr 1 from 1 to (+ 2 *ls.last-cycle*)) (do
  (:= ( [] live 1)
    (vn-set:union
      (vn-set:difference ( [] live (- 1 1) )
        ( [] kill (- 1 1) ) )
      ( [] gen 1) ) ) ) )
live) )

=====
***
*** (SCH.VN:SPLIT-GEN-CYCLE VN)
*** (SCH.VN:JOIN-GEN-CYCLE VN)
*** (SCH.VN:SPLIT-KILL-CYCLE VN)
*** (SCH.VN:JOIN-KILL-CYCLE VN)
***
*** These functions return the first and last cycles that a VN is live, for
*** splits and joins. The -GEN- functions return () if the VN doesn't produce
*** a value.
***
=====
(defun sch.vn:split-gen-cycle ( vn )
  (& (vn:register-bank-me vn)
    (caseq (vn:type vn)
      (defl
        0)
      (operation
        (+ (vn:scheduled-cycle vn) 1) )
      (copy
        (+ (vn:scheduled-cycle vn) 1) )
      (t
        () ) ) ) )

(defun sch.vn:join-gen-cycle ( vn )
  (& (vn:register-bank-me vn)
    (let ( (cycle
      (caseq (vn:type vn)
        (defl
          0)
          (operation
            (+ (vn:scheduled-cycle vn)
              (+ 1 (me:delay (vn:me vn) ) ) ) )
            (copy
              (+ (vn:scheduled-cycle vn) 1) )
            (t
              () ) ) ) )
      (if (& (1 (vn:reading-vns vn) )
        (> cycle (sch.vn:join-kill-cycle vn) ) )
        (then
          () )
        (else
          cycle) ) ) ) )

***
*** Hairiness here: If a VN has no readers but just off-live

```

3

PS:&lt;C.S.BULLDOG.LIST-SCHEDULER.TEST&gt;SCH-SPLITS-JOINS.LSP.59

```

*** readers, and all those off-live readers are scheduled in
*** the same cycle in which VN's value is produced, then for
*** the purposes of join-live the value of VN is never generated
*** (we don't want to report its locations).

(defun sch.vn:split-kill-cycle ( vn )
  (max
    (loop (for reading-vn in (vn:reading-vns vn) )
      (reduce max 0
        (caseq (vn:type reading-vn)
          (operation (vn:scheduled-cycle reading-vn) )
          (copy (vn:scheduled-cycle reading-vn) )
          (use1 (+ 2 *ls.last-cycle*) )
          (t
            (error (list reading-vn "Case error."))))))
    (loop (for off-live-reading-vn in (vn:off-live-reading-vns vn) )
      (reduce max 0
        (+ (vn:scheduled-cycle off-live-reading-vn)
          (+ 1
            (me:delay (vn:me off-live-reading-vn)))))))))

(defun sch.vn:join-kill-cycle ( vn )
  (max
    (loop (for reading-vn in (vn:reading-vns vn) )
      (reduce max 0
        (caseq (vn:type reading-vn)
          (operation (+ (vn:scheduled-cycle reading-vn)
            (me:delay (vn:me reading-vn) ) ) )
          (copy (vn:scheduled-cycle reading-vn) )
          (use1 (+ 2 *ls.last-cycle*) )
          (t
            (error (list reading-vn "Case error."))))))
    (loop (for off-live-reading-vn in (vn:off-live-reading-vns vn) )
      (reduce max 0
        (+ (vn:scheduled-cycle off-live-reading-vn)
          (me:delay (vn:me off-live-reading-vn)))))))))

=====
***
*** (SCH.CYCLE:SPLIT-DEF CYCLE)
***
*** Constructs a DEF for a split at CYCLE (the last cycle of the jump
*** is scheduled for CYCLE).
***
=====
(defun sch.cycle:split-def ( cycle )
  'def
  ..(if *ls.trace-information*
    '( (%trace .%tr.trace-number* %cycle .(+ cycle 1) ) )
    () )
  ..(loop (for-each-vn-set-element ( [] *sch.cycle:split-lives*
    (+ cycle 1) )
    vn)
    (save '(.(vn:name vn)

```

4

```

      ,(me:name (vn:register-bank-me vn) )
      ,(vn:register vn) ) ) ) )
=====
***
***
*** (SCH.CYCLE:JOIN-USE CYCLE)
***
*** Constructs a USE for a join to the beginning of CYCLE.
***
=====
(defun sch.cycle:join-use ( cycle )
  'use
  ..(if *ls.trace-information?*
        '( (trace ,*tr.trace-number* %cycle .(+ cycle 1) ) )
        () )

  ..(loop (for-each-vn-set-element ([]) *sch.cycle:join-live* cycle) vn)
  (save '(, (vn:name vn)
          ,(me:name (vn:register-bank-me vn) )
          ,(vn:register vn) ) ) ) )
=====
***
***
*** (SCH.CYCLE:SPLIT-PARTIAL-SCHEDULE SPLIT-CYCLE)
***
*** Returns the partial schedule for a split at SPLIT-CYCLE. The partial
*** schedule consists of those operations spanning the boundary between
*** SPLIT-CYCLE and the succeeding cycle.
***
=====
(defun sch.cycle:split-partial-schedule ( split-cycle )
  (loop (incr cycle from (+ split-cycle 1) to *ls.max-schedule-size*)
    (bind partial-cycle
      (loop (for vn&oper in ([]) *ls.schedule* cycle) )
      (bind (vn . oper) vn&oper)
      (when (&& (== 'operation (vn:type vn) )
              (<= (vn:scheduled-cycle vn) split-cycle) ) )
      (save vn&oper) ) )
    (while partial-cycle)
  (save
    (if *ls.trace-information?* (then
      '( ( () . ( () trace ,*tr.trace-number* split .(+ 1 split-cycle)))
        ..partial-cycle) )
      (else
        partial-cycle) ) ) ) )
=====
***
***
*** (SCH.CYCLE:JOIN-PARTIAL-SCHEDULE JOIN-CYCLE)
***
*** Returns the partial schedule for a join at JOIN-CYCLE. The partial
*** schedule consists of those operations spanning the boundary between
*** JOIN-CYCLE and the previous cycle.
***
=====
(defun sch.cycle:join-partial-schedule ( join-cycle )
  (dreverse

```

```

(loop (decr cycle from (- join-cycle 1) to 0)
  (bind partial-cycle
    (loop (for vn&oper in ([]) *ls.schedule* cycle) )
    (bind (vn . oper) vn&oper)
    (when (&& (== 'operation (vn:type vn) )
            (<= join-cycle
                (+ (vn:scheduled-cycle vn)
                    (me:delay (vn:me vn) ) ) ) ) )
    (save vn&oper) ) )
  (while partial-cycle)
  (save
    (if *ls.trace-information?* (then
      '( ( () . ( () trace ,*tr.trace-number* join .(+ 1 join-cycle)))
        ..partial-cycle) )
      (else
        partial-cycle) ) ) ) )

```



```

*** Copyright (C) 1988 John R. Ellis. This file may be used for
*** educational and research purposes only, by the faculty, students,
*** and staff of Yale University only.
=====
LIST SCHEDULER

This module implements the main part of the list scheduler and its
interface to the outside world:

(INITIALIZE-CODE-GENERATOR)
(GENERATE-CODE LIVE-BEFORE SOURCE-RECORD-LIST LIVE-AFTER)
(SCHEDULE:LENGTH SCHEDULE)
(SCHEDULE:[] SCHEDULE I)
(SCHEDULE:JOIN SCHEDULE I)
(SCHEDULE:SPLIT SCHEDULE I JUMP-NUMBER)

These procedures are all documented in
DOCUMENTATION:CODE-GEN-INTERFACE.DOC.

For the implementation details, start with the body of GENERATE-CODE.
The list scheduler proper starts with SCH.SCHEDULE below.
=====

(eval-when (compile load)
  (include list-scheduler:declarations) )
(eval-when (compile)
  (build '(list-scheduler:heap) ) )

(declare (special
  *tr.trace-number*      ;*** For debugging only.
  ) )

(defvar *ls.max-schedule-size* 400)

(defun initialize-code-generator ()
  (vn.initialize)
  (dag.initialize)
  (sch.initialize) )

(defun generate-code ( live-before source-record-list live-after )
  (vn.initialize)
  (dag.make live-before source-record-list live-after)
  (bug.assign-mes)

  (msg 0 t "Last BUG cycle = "
    (loop (for-each-vn vn)
      (when (vn:bug-cycle vn) )
      (maximize
        (+ (vn:bug-cycle vn)
          (caseq (vn:type vn)
            (operation (me:delay (vn:me vn) ) )
            (t 0) ) ) ) )
    t)

  (sch.schedule)

  *ls.schedule*)

```

1

PS:&lt;C.S.BULLDOG.LIST-SCHEDULER.TEST&gt;SCHEDULER.LSP.163

```

(defun schedule:length ( schedule )
  (if (== schedule *ls.schedule*)
      (+ *ls.last-cycle* 1)
      (length schedule) ) )

(defun schedule:[] ( schedule i )
  (let* ( (vn&oper-list
          (if (== schedule *ls.schedule*)
              ( [] *ls.schedule* (- i 1) )
              (nth-elt schedule i) ) )
        (result
         (loop (for (vn . oper) in vn&oper-list) (save
            '(,oper ,( && vn (vn:datum vn) ) ) ) ) ) )
    (if (&& *ls.trace-information*
        (== schedule *ls.schedule*) )
        (then
          '( ( ( () trace ,*tr.trace-number* ) )
            ..result) )
        (else
          result) ) ) )

(defun schedule:join ( schedule i )
  (assert (== schedule *ls.schedule*))
  (let ( (cycle (- i 1) )
        '(, (sch.cycle:join-use cycle)
          ,(sch.cycle:join-partial-schedule cycle) ) ) ) )

(defun schedule:split ( schedule i jump-number )
  (assert (== schedule *ls.schedule*))
  (let ( (cycle (- i 1) )
        '(, (sch.cycle:split-def cycle)
          ,(sch.cycle:split-partial-schedule cycle) ) ) ) )

=====
***
*** (SCH.SCHEDULE)
***
*** This procedure is the top-level list-scheduler that actually schedules
*** VNs on the schedule.
***
=====

(declare (special
  *sch.data-ready-heap*
  ;***
  ;*** This is the list-scheduler's data-ready queue, implemented
  ;*** as a HEAP. When a VN is scheduled, the scheduling
  ;*** procedure adds any newly-data-ready VNs to the HEAP.
  ;*** VNs are kept sorted by VN:CYCLE, the earliest cycle that
  ;*** the VN could be scheduled; within the same cycle, VNs
  ;*** of larger height will take priority.

  *sch.split-live*
  *sch.join-live*
  ;***
  ;*** From SCH-SPLITS-JOINS.

```

2

```

*sch.vns-to-release-off-live*
***
*** VNs that are "off-live read" by a conditional jump cannot
*** have their readers-left count deallocated until the start
*** of the cycle AFTER the cycle containing the conditional
*** jump. This variable holds the list of the off-live VNs
*** of jumps scheduled in the current cycle; after the cycle
*** is scheduled we'll go over the list, releasing the
*** off-live VNs.
) )

(defun sch.initialize ()
  (:= *ls.schedule*          ( ) )
  (:= *ls.last-cycle*       ( ) )
  (:= *ls.cycle*            ( ) )
  (:= *sch.data-ready-heap* ( ) )
  (:= *sch.vns-to-release-off-live* ( ) )

  (:= *sch.split-live*      ( ) )
  (:= *sch.join-live*       ( ) )

  ;*** Initialize support modules.
  ;***
  (registers.initialize)
  (resource.initialize-schedule 'ls)
  ( ) )

(defun sch.schedule ()
  (sch.initialize)

  ;*** Initialize data structures.
  ;***
  (:= *ls.schedule*          (makevector *ls.max-schedule-size*) )
  (:= *ls.last-cycle*       -1)
  (:= *ls.cycle*            0)
  (:= *sch.data-ready-heap* (heap:create 'sch.vn:vn:heap-compare) )

  ;*** Initialize fields in the DAG.
  ;***
  (dag.set-counts)
  (loop (for-each-vn vn) (do
    (:= (vn:cycle vn) 0) ) )

  ;*** Mark USE1-ed registers as "avoid if possible."
  ;***
  (sch.mark-used-registers)

  ;*** Schedule all the VNs in the DAG with no predecessors (DEFS and
  ;*** pseudo-ops). This will fill the data-ready queue with "real"
  ;*** VNs.
  ;***
  (loop (for vn in *ls.entry-vns*) (do
    (sch.vn:schedule vn) ) )

  ;*** While there are more VNs to schedule, pick one out of the
  ;*** data-ready queue and try to schedule it.
  ;***
  (loop (while (> (heap:size *sch.data-ready-heap*) 0) )
    (bind vn (heap:delete *sch.data-ready-heap*) )

```

```

(do
  (assert (! (vn:scheduled-cycle vn) )
    (h vn) t "SCHEDULER: VN already scheduled.")

  (if (!= *ls.cycle* (vn:cycle vn) ) (then
    (sch.release-off-live) ) )
  (:= *ls.cycle* (vn:cycle vn) )

  (sch.vn:schedule vn) )
(result
  (sch.release-off-live) ) )

;*** Run various debugging consistency checks.
;***
(sch.check-consistency)

;*** Compute the live-dead information.
;***
(sch.compute-live-dead)

;*** Sort the operations on the schedule by source order (using
;*** VN:NUMBER -- groan). This overkill is only to keep jumps
;*** scheduled in the same cycle in source order.
;***
(loop (incr cycle from 0 to *ls.last-cycle*) (do
  (:= ([] *ls.schedule* cycle)
    (sort &&& (f:l ( vn&oper1 vn&oper2 )
      (< (vn:number (car vn&oper1) )
        (vn:number (car vn&oper2) ) ) ) ) ) ) ) ) ) )

) )

;***=====
;***
;*** (SCH.MARK-USED-REGISTERS)
;***
;*** Marks every register mentioned in a USE1 as "avoid if possible".
;*** The register allocator will avoid these registers if possible, thus
;*** reducing the amount of copying we need to do.
;***
;***=====

(defun sch.mark-used-registers ()
  (loop (for vn in *ls.exit-vns*)
    (when (== 'use (vn:type vn) ) )
  (do
    (loop (for use1-vn in (vn:operand-vns vn) )
      (when (vn:register use1-vn) )
    (do
      (vn:record-preferred-location use1-vn) ) ) ) )
  ( ) )

;***=====
;***
;*** (SCH.CHECK-CONSISTENCY)
;***
;*** Performs various debugging and consistency checks after forming a
;*** schedule.
;***
;***=====

(defun sch.check-consistency ()

```

```

;*** Make sure all the VNs were scheduled.
;***
(loop (for-each-vn vn)
      (when (not (memq (vn:type vn) '(def use) ) ) )
      (do
        (assert (vn:scheduled-cycle vn)
                (h vn) t "SCHEDULER: VN wasn't scheduled.") )

;*** Make sure registers were freed up properly.
;***
(loop (for me in *ls.register-bank-mes*) (do
      (loop (incr i from 0 to (- (me:size me) 1) )
            (bind vn (cycle:me:register:occupied?
                      (+ 1 *ls.last-cycle*) me 1) )
            (when vn)
            (do
              (assert (for-some (reading-vn in (vn:reading-vns vn) )
                            (== 'use1 (vn:type reading-vn) ) )
                      "Register " i " in bank " (me:name me) " containing: " t
                      (h vn) t
                      "wasn't freed up.") ) ) ) )
      (do
        (assert (for-some (reading-vn in (vn:reading-vns vn) )
                      (== 'use1 (vn:type reading-vn) ) )
                "Register " i " in bank " (me:name me) " containing: " t
                (h vn) t
                "wasn't freed up.") ) ) ) )
) )

```

```

;***
;*** (SCH.VN:CYCLE:ENQUEUE VN CYCLE)
;***
;*** Places VN on the data ready queue to be scheduled no earlier than
;*** CYCLE.
;***
;***
(defun sch.vn:cycle:enqueue (vn cycle)
  (:= (vn:cycle vn) (max &&& cycle) )
  (heap:insert *sch.data-ready-heap* vn)
  () )

```

```

;***
;*** (SCH.VN:DEQUEUE VN)
;***
;*** Removes VN from the data ready queue "prematurely" if it is in the
;*** queue. If it isn't, nothing is done. It will get requeued eventually
;*** later on.
;***
;***

```

```

(defun sch.vn:dequeue (vn)
  (if (== 0 (vn:predecessors-left vn) ) (then
    (heap:remove *sch.data-ready-heap* vn) ) )
  () )

```

```

;***
;*** (SCH.VN:VN:HEAP-COMPARE VN1 VN2)
;***
;*** Comparison function used by the list-scheduler priority queue. VN1
;*** takes priority over VN2 if either its :CYCLE is earlier or, if the

```

```

;*** :CYCLES are the same, if its :HEIGHT is larger; or if the :HEIGHTs
;*** are the same, if :BUG-CYCLE is earlier.
;***
;***
;***

```

```

(defun sch.vn:vn:heap-compare (vn1 vn2)
  (? ( (< (vn:cycle vn1) (vn:cycle vn2) )
      t)
      ( (== (vn:cycle vn1) (vn:cycle vn2) )
        (? ( (> (vn:height vn1) (vn:height vn2) )
            t)
            ( (== (vn:height vn1) (vn:height vn2) )
              (if (&& (vn:bug-cycle vn1)
                  (vn:bug-cycle vn2) )
                (then
                  (< (vn:bug-cycle vn1) (vn:bug-cycle vn2) ) )
                (else
                  t ) ) ) )
            ( t ) ) ) )
      ( t ) ) )

```

```

;***
;*** (SCH.VN:SCHEDULE VN)
;***
;*** This procedure tries to schedule VN by dispatching to the appropriate
;*** procedure according to VN's type. Each such procedure tries to schedule
;*** VN in the current cycle. If it succeeds, any new data-ready successor
;*** VNs are enqueued in the data-ready heap. If it fails, it enqueues VN
;*** back in the heap for a later cycle.
;***
;***

```

```

(defun sch.vn:schedule (vn)
  (caseq (vn:type vn)
    (pseudo-op (sch.pseudo-op-vn:schedule vn) )
    (def (sch.def-vn:schedule vn) )
    (operation (sch.operation-vn:schedule vn) )
    (copy (sch.copy-vn:schedule vn) )
    (use1 (sch.use1-vn:schedule vn) )
    (t (error (list vn "Case error.") ) ) ) )

```

```

;***
;*** (SCH.PSEUDO-OP-VN:SCHEDULE VN)
;***
;*** Schedules a pseudo-op VN by placing its pseudo-operation on the
;*** schedule.
;***
;***

```

```

(defun sch.pseudo-op-vn:schedule (vn)
  (:= (vn:scheduled-cycle vn) *ls.cycle*)
  (push ( [] *ls.schedule* *ls.cycle* )
        '(.vn () ,(vn:oper vn) ) )

```

```

(:= *ls.last-cycle* (max &?? *ls.cycle*))
() )

=====
***
*** (SCH.DEF-VN:SCHEDULE VN)
***
*** Schedules a DEF VN merely by scheduling all of its readers which are
*** DEF1 VNs.
***
=====

(defun sch.def-vn:schedule ( vn )
  (:= (vn:scheduled-cycle vn) *ls.cycle*)

  (loop (for reading-vn in (vn:reading-vns vn) ) (do
    (sch.def1-vn:schedule reading-vn) ) )
  () )

=====
***
*** (SCH.DEF1-VN:SCHEDULE VN)
***
*** VN is a DEF1. If VN is a constant DEF, then all of its successors
*** are released into the scheduling queue.
***
*** If VN is a register bank, then a register is allocated and assigned
*** to to it, and all the successors are released into the scheduling
*** queue. If this is a loop trace and VN's name is written on this
*** trace, then we assume the name is an induction variable; we set the
*** locations of all the USE1s of the same name that don't currently
*** have a register location to be the register just allocated to this
*** DEF VN. This will reduce copying of the induction variables, and
*** any copying that must be done will be integrated into the body of
*** the loop (presumably filling "holes" in the schedule).
***
=====

(defun sch.def1-vn:schedule ( vn )
  (let ( (dest-me (vn:me vn) )
        (dest-register ( ) ) )

    (:= (vn:scheduled-cycle vn) *ls.cycle*)

    (if (== 'constant-generator (me:type dest-me) ) (then
      (sch.vn:splice-copies vn) )

      (else
        (:= dest-register
          (cycle:me:vn:allocate-register
            0 dest-me vn (vn:register vn) ) )

        (if (& *ls.loop-trace*
          (< 0 (dag.name:trace-write-count (vn:name vn) ) ) )
          (then
            (if-let ( (use-vn (dag.name:use-vn (vn:name vn) ) ) ) (then
              (loop (for use1-vn in (vn:operand-vns use-vn) )
                (when (! (vn:register use1-vn) ) )
                (do
                  (:= (vn:register-bank-me use1-vn) dest-me)
                  (:= (vn:me use1-vn) dest-me)
                )
              )
            )
          )
        )
      )
  )
)

```

```

(:= (vn:likely-mes use1-vn) '(dest-me) )
(:= (vn:register use1-vn) dest-register)
(vn:record-preferred-location use1-vn))))))

(sch.vn:me:register:cycle:assign-location
  vn
  dest-me
  dest-register
  () ) )

(sch.vn:release-successors vn)
() )

=====
***
*** (SCH.USE1-VN:SCHEDULE VN)
***
*** VN is a USE1. It is "scheduled" just by insuring that if it requires
*** a specific location then the value is actually there.
***
=====

(defun sch.use1-vn:schedule ( vn )
  (:= (vn:scheduled-cycle vn) *ls.cycle*)
  (if (vn:register-bank-me vn) (then
    (let ( (operand-vn (car (vn:operand-vns vn) ) ) )
      (assert (& (== (vn:register-bank-me vn)
                    (vn:register-bank-me operand-vn) )
              (== (vn:register vn)
                  (vn:register operand-vn) ) ) )
      (h vn) t
      "A value didn't end up in its USE location." ) ) )
  () )

=====
***
*** (SCH.NOOP-VN:SCHEDULE VN)
***
*** VN is either a COPY or an OPERATION which we've decided shouldn't
*** produce any machine operations. But because the bookkeeper interface
*** needs a representative for each source operation on the schedule, we
*** need to schedule a NOOP machine operation.
***
*** Currently the only NOOPs arise from some operation that produces
*** a value whose only use is as an off-live variable at a split. If
*** the operation is scheduled below the split, it can be scheduled as
*** NOOP that doesn't consume any resources.
***
=====

(defun sch.noop-vn:schedule ( vn )
  (:= (vn:scheduled-cycle vn) *ls.cycle*)

  (sch.vn:release-operands vn)
  (sch.vn:release-successors vn)
  (sch.vn:release-off-live vn)

  (push ([] *ls.schedule* *ls.cycle*)
    (vn
      ()
      noop
    )
  )
)

```

```

    ..(vn:oper vn) )
    (:= *ls.last-cycle* (max &&& *ls.cycle*))
    () )

```

```

=====
***
*** (SCH.VN:RELEASE-OPERANDS VN)
***
*** For each operand of VN, this procedure decrements the readers-left
*** count. If the count reaches 0, the VN's register is deallocated.
***
=====

```

```

(defun sch.vn:release-operands ( vn )
  (loop (for operand-vn in (vn:operand-vns vn) ) (do
    (sch.vn:read-release operand-vn) ) )
  () )

```

```

=====
***
*** (SCH.VN:RELEASE-OFF-LIVE VN)
*** (SCH.RELEASE-OFF-LIVE)
***
*** SCH.VN:RELEASE-OFF-LIVE performs a function similar to
*** SCH.VN:RELEASE-OPERANDS. But the readers-left count of the off-live
*** VNs can't be decremented until after the end of the current cycle
*** (because an off-live VN's value must stay alive until after the jump
*** completes). So we just remember this VN, and at the end of this
*** cycle SCH.RELEASE-OFF-LIVE will be called to do the actual release.
***
=====

```

```

(defun sch.vn:release-off-live ( vn )
  (if (vn:off-live-vns vn) (then
    (push *sch.vns-to-release-off-live* vn) ) )
  () )

```

```

(defun sch.release-off-live ()
  (loop (for vn in *sch.vns-to-release-off-live* ) (do
    (loop (for off-live-vn in (vn:off-live-vns vn) ) (do
      (sch.vn:read-release off-live-vn) ) ) )
    (:= *sch.vns-to-release-off-live* () )
  () )

```

```

=====
***
*** (SCH.VN:READ-RELEASE VN)
***
*** Decrements the readers-left count of VN, and if it reaches 0, then
*** its register is deallocated.
***
*** A VN's register cannot be reused by somebody else until the maximum,
*** over all the readers, of the last cycle of the reader, and the
*** maximum, over all the off-live readers, of the cycle following the
*** end of the off-live reader. This guarrantess that the a read register
*** is "alive" for the duration of all the operations that read it.
***
=====

```

```

(defun sch.vn:read-release ( vn )

```

```

  (:= (vn:readers-left vn) (- &&& 1) )
  (if (&& (<= (vn:readers-left vn) 0)
    (vn:register-bank-me vn) )
  (then
    (let ( (cycle-free
      (max
        (loop (for reading-vn in (vn:reading-vns vn) )
          (reduce max 0
            (caseq (vn:type reading-vn)
              (operation
                (+ (vn:scheduled-cycle reading-vn)
                  (me:delay (vn:me reading-vn) ) ) )
              (copy
                (vn:scheduled-cycle reading-vn) )
              (t
                (error (list reading-vn "Case error."))))))
        (loop (for off-live-reading-vn
          in (vn:off-live-reading-vns vn) )
          (reduce max 0
            (+ (vn:scheduled-cycle off-live-reading-vn)
              (+ 1
                (me:delay (vn:me off-live-reading-vn)))))))))
      (cycle:me:register:deallocate cycle-free
        (vn:register-bank-me vn)
        (vn:register vn) ) ) ) )
    () )

```

```

=====
***
*** (SCH.VN:RELEASE-SUCCESSORS VN)
***
*** Decrements the predecessors-left count of each successor of VN. If
*** it reaches 0, the successor is enqueued for scheduling. A reading
*** successor is enqueued for the cycle that the value of VN becomes
*** available; a constrained successor is enqueued for the current cycle
*** plus the delay of the constraint.
***
=====

```

```

(defun sch.vn:release-successors ( vn )
  (let ( (available-cycle
    (caseq (vn:type vn)
      (operation copy
        (vn:register-cycle vn) )
      (t
        *ls.cycle* ) ) )
    (loop (for reading-vn in (vn:reading-vns vn) ) (do
      (:= (vn:cycle reading-vn) (max &&& available-cycle) )
      (if (>= 0 (:= (vn:predecessors-left reading-vn) (- &&& 1) ) )
        (then
          (sch.vn:cycle:enqueue reading-vn (vn:cycle reading-vn) ) ) ) )
      (loop (for constrained-vn in (vn:constrained-vns vn) ) (do
        (:= (vn:cycle constrained-vn)
          (max &&&
            (cycle:delays:constrained-cycle
              *ls.cycle*

```

```

(vn:delay vn)
(vn:delay constrained-vn)
(vn:constrained-vn:delay vn constrained-vn))))
(if (>= 0 (:= (vn:predecessors-left constrained-vn) (- &&& 1) ) )
  (then
    (sch.vn:cycle:enqueue constrained-vn
      (vn:cycle constrained-vn))))))
( ) )

=====
***
*** (SCH.VN:ME:REGISTER:CYCLE:ASSIGN-LOCATION VN ME REGISTER CYCLE)
***
*** Assigns a register location to VN that is available at time CYCLE
*** and later. Updates the live-dead table to signify a new live VN
*** starting at CYCLE. Copies are spliced between VN and all the readers
*** of VN that can't read its value directly.
***
=====
(defun sch.vn:me:register:cycle:assign-location ( vn me register cycle )
  (:= (vn:register-bank-me vn) me)
  (:= (vn:register vn) register)
  (:= (vn:register-cycle vn) cycle)

  (sch.vn:splice-copies vn)
  ( ) )

=====
***
*** (SCH.VN:SPLICE-COPIES VN)
***
*** Splices one or more copies between VN and those readers that can't
*** read the current location of VN directly. Each reader is examined
*** in turn to see if it can read the current location of VN. All the
*** readers that can't are grouped according to the first ME on the path
*** from the location of VN to the desired destination of the reader
*** (a functional unit or register bank). Between VN and each such group
*** is spliced a new COPY VN.
***
=====
(defun sch.vn:splice-copies ( vn )
  (let ( (first-on-path&reading-vns-table ( ) ) )

    (loop (for reading-vn in (vn:reading-vns vn)
      (bind first-on-path (sch.vn:reading-vn:needs-copy?
        vn reading-vn) )

      (when first-on-path)

      (do
        (if-let ( (me&reading-vns (assoc first-on-path
          first-on-path&reading-vns-table)))
          (then
            (push (cdr me&reading-vns) reading-vn) )
          (else
            (push first-on-path&reading-vns-table
              '(,first-on-path ,reading-vn) ) ) ) ) )

    (loop (for ( ( ) . reading-vns) in first-on-path&reading-vns-table)
      (bind copy-vn (vn:create (vn:new
        type 'copy

```

```

name (vn:name vn)
height (vn:height vn)
cycle *ls.cycle*
readers-left (length reading-vns)
predecessors-left 1) ) )
(do
  (vn:splice-vn vn copy-vn reading-vns ( )
    (:= (vn:readers-left vn) (+ 1 (- &&& (length reading-vns) ) ) ) ) )
  ( ) )

=====
***
*** (SCH.VN:READING-VN:NEEDS-COPY? VN READING-VN)
***
*** Returns non-() if READING-VN can't read the value of VN from its current
*** location, false o.w. The non-() value returned is usually the first ME
*** on the path from the location of VN to wherever READING-VN wants it.
***
=====
(defun sch.vn:reading-vn:needs-copy? ( vn reading-vn )
  (let ( (source-me (sch.vn:location-me vn) )

    ( ?
      *** READING-VN is an operation. A COPY is needed if the
      *** location of VN is not one of the direct inputs of
      *** the functional unit assigned READING-VN, or if
      *** the current location of VN conflicts with the other
      *** operands of READING-VN.
      ***
      ( (& (== 'operation (vn:type reading-vn) )
        (|| (! (memq source-me (me:inputs (vn:me reading-vn) ) ) )
          (sch.vn:me:reading-vn:conflict?
            vn source-me reading-vn) ) )
        (me:me:first-on-path source-me (vn:me reading-vn) ) )

      *** READING-VN is a USE1. A COPY is needed if VN's
      *** location is a constant generator, or if VN's register
      *** location is not the same as that required by the
      *** USE1.
      ***
      ( (== 'use1 (vn:type reading-vn) )
        (?( (== 'constant-generator (me:type source-me) )
          (if (! (vn:register-bank-me vn) ) (then
            t)
            (else
              (me:me:first-on-path source-me
                (vn:register-bank-me reading-vn))))))

        ( (& (vn:register-bank-me reading-vn)
          (|| (!== source-me (vn:register-bank-me reading-vn) )
            (!== (vn:register vn) (vn:register reading-vn))))
          (me:me:first-on-path source-me
            (vn:register-bank-me reading-vn) ) )

        ( t
          ( ) ) ) ) )

      ( t
        ( ) ) ) ) )

```

```

=====
***
*** (SCH.VN:ME:READING-VN-CONFLICTS VN ME)
***
*** Returns the number of readers of VN that conflict with using ME
*** for the result of VN. ME is either the destination register bank
*** to be used for VN or else the constant generator assigned to VN.
***
*** A reader of VN, RVN conflicts with the choice of ME for storing VN
*** if it isn't possible for RVN to read out all of its operands in the
*** same cycle (e.g. because there aren't enough ports).
***
*** Currently, we assume that the only possible resource that could cause
*** conflict is the number of available ports (i.e. all ports of a
*** register bank are connected to the identical set of destinations).
*** So checking to see if ME conflicts with previously assigned operands
*** is merely a matter of counting the number of ME ports required by
*** all the operands.
***
=====
(defun sch.vn:me:reading-vn-conflicts (vn me)
  (loop (for reading-vn in (vn:reading-vns vn))
        (when (sch.vn:me:reading-vn:conflict? vn me reading-vn)
              (reduce + 0 1) ) ) )

(defun sch.vn:me:reading-vn:conflict? (vn me reading-vn)
  (if (== 'constant-generator (me:type me)) (then
    (loop (for operand-vn in (vn:operand-vns reading-vn))
          (when (== vn operand-vn)
                (when (== me (vn:me operand-vn) )
                      (do
                        (return t) )
                        (result () ) ) )
          )
      )
      (else
        (let ( (read-ports (me:read-ports me) ) ) ;*** Compiler bug.
              (> (sch.vn:me:reading-vn:required-ports vn me reading-vn)
                  read-ports) ) ) ) )

(defun sch.vn:me:reading-vn:required-ports (vn me reading-vn)
  (loop (for operand-vn in (vn:operand-vns reading-vn))
        (when (!= vn operand-vn)
              (when (== me (vn:register-bank-me operand-vn) )
                    (reduce + 1 1) ) )
        )

=====
***
*** (SCH.VN:MAX-HEIGHT-READING-VN VN)
***
*** Returns the reader of VN of maximum height.
***
=====
(defun sch.vn:max-height-reading-vn (vn)
  (loop (for reading-vn in (vn:reading-vns vn))
        (maximize reading-vn (vn:height reading-vn) ) ) )
=====

```

```

=====
***
*** (SCH.VN:MAX-HEIGHT-DEST-ME VN)
***
*** Finds the reader of VN that has maximum height, and then returns
*** the destination ME that VN must deliver its result to for it to be
*** used by the reader. The ME returned depends on the type of
*** READING-VN:
***
*** OPERATION: the functional unit assigned to the VN.
*** COPY: wherever the max height reader of VN wants its value.
*** USE1: any register bank assigned to the use.
***
*** Returning () means that we don't have any particular preference for
*** a destination ME; this should only occur if VN is a USE1 and doesn't
*** have a required location.
***
*** GROAN: We use an unexported function from BUG.
***
=====
(defun sch.vn:max-height-dest-me (vn)
  (bug.vn:dest-me
   (loop (for reading-vn in (vn:reading-vns vn))
         (maximize reading-vn (vn:height reading-vn) ) ) ) )

=====
***
*** (SCH.VN:LOCATION-ME VN)
***
*** Returns the current ME location of the value of VN which must be
*** read to get the value of VN. Normally, this is the register bank
*** ME, unless VN is a DEF1 assigned a constant generator, in which case
*** the c.g. ME is returned.
***
=====
(defun sch.vn:location-me (vn)
  (if (&& (vn:me vn)
      (== 'constant-generator (me:type (vn:me vn) ) ) )
      (vn:me vn)
      (vn:register-bank-me vn) ) )

(defun sch.print-schedule ()
  (msg 0 t "Schedule:" t)
  (loop (incr i from 0 to *ls.last-cycle*) (do
    (msg (j i 8) " : "
      (h (loop (for (vn . oper) in ([]) *ls.schedule* i) ) (save
        oper) )
        t) ) )
    ) )
) )

```

```

:***-----*****
:*** Copyright (C) 1988 John R. Ellis. This file may be used for *****
:*** educational and research purposes only, by the faculty, students, *****
:*** and staff of Yale University only. *****
:***-----*****

```

```

: A pipelined sequential model used for comparing with ELI-MODEL.LSP.
: One operation can be initiated every cycle, and the register port can
: read and write a value every cycle.
:-----

```

```

(def-machine-model

```

```

    :*** Register bank R
    :***
    (def-me (name      r)
           (type      register-bank)
           (size      #.( * 8 84) )
           (inputs    (fu1 fu3 fu4 r c) )
           (read-ports 8)
           (write-ports 1)
           (read-resources ( (r-read) ) )
           (write-resources ( (r-write) ) ) )
    (def-resource-class r-read 8)
    (def-resource-class r-write 1)

    :*** Constant generator C
    :***
    (def-me (name      c)
           (type      constant-generator)
           (resources ( (constant) ) )
           (constraint-function sm.immediate-constant? ) )
    (def-resource-class constant 1)

    :*** 1-cycle functional unit FU1
    :***
    (def-me (name      fu1)
           (type      functional-unit)
           (delay      0)
           (inputs    (r c) )
           (resources ( (cycle) ) )
           (operators (float fix inot idiv isub leq imax imin iadd ineg ior
                        ige ilt ile ine iexp isul igt iand iabs bitrev
                        if-true if-false if-ilt if-igt if-leq if-ine if-ile
                        if-ige if-leq if-ine if-ile if-ige) ) )

    :*** 8-cycle functional unit FU3
    :***
    (def-me (name      fu3)
           (type      functional-unit)
           (delay      2)
           (inputs    (r c) )
           (resources ( (cycle) ) )
           (operators (fsub fadd fneg fmin fmax fabs fdiv cos sin sqrt flt
                        fgt feq fne file fge) ) )

    :*** 4-cycle functional unit FU4
    :***
    (def-me (name      fu4)

```

```

           (type      functional-unit)
           (delay      3)
           (inputs    (r c) )
           (resources ( (cycle) ) )
           (operators (vbase ivload fvload ipload fpload ivstore fvstore
                        ipstore fpstore faul isul) ) )

```

```

    (def-resource-class cycle 1) )

```

```

(defun sm.immediate-constant? ( constant )

```

```

  (? ( inump constant)
     (&& (< constant 2047)
          (> constant -2048) ) )
  ( = 0 constant)
  t)
  ( (consp constant)
    (== 'address (car constant) ) )
  t ) ) )

```

```

(defun load-constant? ( constant )

```

```

  (if (sm.immediate-constant? constant) (then
     (!= 'load *ls.immediate-constant-actions*) )
    (else
     t) ) )

```



```

*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for *****
*** educational and research purposes only, by the faculty, students, *****
*** and staff of Yale University only. *****
*****
=====
:
: The ELI model used by Rutt's CG experiments.
:
:
=====
(alias ^^^ 'atomconcat)

(defun em.cluster ( i (unique-name unique-fields) )
  (let ( (ci (atomconcat 'c i) )
        (ci-3 (atomconcat 'c (mod (- i 3) 8) ) )
        (ci-1 (atomconcat 'c (mod (- i 1) 8) ) )
        (ci+1 (atomconcat 'c (mod (+ i 1) 8) ) )
        (ci+3 (atomconcat 'c (mod (+ i 3) 8) ) ) )
    '(

      ;*** Register bank CiR
      ;***
      (def-me (name .(^^ ci 'r) )
              (type register-bank)
              (size 64)
              (inputs (.(^^ ci-3 'r)
                      .(^^ ci-1 'r)
                      .(^^ ci 'r)
                      .(^^ ci+1 'r)
                      .(^^ ci+3 'r)
                      .(^^ ci 'c)
                      .(^^ ci 'm)
                      .(^^ ci '+)
                      .(^^ ci '=)
                      .(^^ ci unique-name) ) )
              (read-ports 8)
              (write-ports 3)
              (read-resources (.(^^ ci 'r-read) ) )
              (write-resources (.(^^ ci 'r-write) ) ) )
      (def-resource-class .(^^ ci 'r-read) 8)
      (def-resource-class .(^^ ci 'r-write) 8)

      ;*** Integer adder Ci+
      ;***
      (def-me (name .(^^ ci '+) )
              (type functional-unit)
              (delay 0)
              (inputs (.(^^ ci 'r)
                      .(^^ ci 'c) ) )
              (resources (.(^^ ci '+) ) )
              (operators (float fix inot idiv isub leq imax imin
                          iadd ior ige ilt ile ine iexp inul igt
                          iand iabs bitrev) ) )
      (def-resource-class .(^^ ci '+) 1)

      ;*** Integer tester Ci=
      ;***
      (def-me (name .(^^ ci '=) )
              (type functional-unit)
              (delay 0)
              (inputs (.(^^ ci 'r)

```

```

                      .(^^ ci 'c) ) )
              (resources (.(^^ ci '=) ) )
              (operators (if-true if-false if-ilt if-igt if-ieq if-ine
                          if-ile if-ige if-ileq if-ine if-ile if-ige)))
      (def-resource-class .(^^ ci '=) 1)

      ;*** Memory CiM
      ;***
      (def-me (name .(^^ ci 'm) )
              (type functional-unit)
              (delay 2)
              (inputs (.(^^ ci 'r)
                      .(^^ ci 'c) ) )
              (resources (.(^^ ci 'm) ) )
              (operators (fvbase ivload fvload ipload ipload ivstore
                          fvstore ipstore fpstore) ) )
      (def-resource-class .(^^ ci 'm) 1)

      ;*** Constant generator CiC
      ;***
      (def-me (name .(^^ ci 'c) )
              (type constant-generator)
              (resources (.(^^ ci 'c) ) )
              (constraint-function
                (lambda ( constant )
                  (?( (inusp constant)
                     (&& (< constant 2047)
                          (> constant -2048) ) )
                    ( = 0 constant)
                    t)
                  ( (consp constant)
                    (== 'address (car constant) ) )
                  t)
                ( ) ) ) )
      (def-resource-class .(^^ ci 'c) 1)

      ;*** Unique functional-unit
      ;***
      (def-me (name .(^^ ci unique-name) )
              (type functional-unit)
              (inputs (.(^^ ci 'r)
                      .(^^ ci 'c) ) )
              (resources (.(^^ ci unique-name) ) )
              .,unique-fields)
      (def-resource-class .(^^ ci unique-name) 1)

      ;*** Bus connections
      ;***
      (def-bus .(^^ ci-3 'r) .(^^ ci 'r) 1)
      (def-bus .(^^ ci+3 'r) .(^^ ci 'r) 1)
      (def-bus .(^^ ci-1 'r) .(^^ ci 'r) 1)
      (def-bus .(^^ ci+1 'r) .(^^ ci 'r) 1)
    ) ) )

(= em.floating-test
 '(f=
  ( (delay 0)
    (operators (if-flt if-fgt if-feq if-fne if-file if-fge) ) ) ) )

(= em.table-lookup

```

```

'(t
  ( (delay 0)
    (operators (fdiv idiv cos sin sqrt) ) ) ) )
(:= em.floating-multiply
  '(f+
    ( (delay 3)
      (operators (fmul imul) ) ) ) )
(:= em.floating-add
  '(f+
    ( (delay 2)
      (operators (fsub fadd fain fmax fabs) ) ) ) )
(eval '(def-machine-model
  ..(em.cluster 0 em.floating-test)
  ..(em.cluster 1 em.table-lookup)
  ..(em.cluster 2 em.floating-multiply)
  ..(em.cluster 3 em.floating-add)
  ..(em.cluster 4 em.floating-test)
  ..(em.cluster 5 em.table-lookup)
  ..(em.cluster 6 em.floating-multiply)
  ..(em.cluster 7 em.floating-add)
  ) )

```

```

*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for
*** educational and research purposes only, by the faculty, students,
*** and staff of Yale University only.
*****
=====
: SHORTEST PATH
:
: This module provides functions for constructing and accessing the
: shortest paths between machine elements.
:
: (SHORTEST-PATH.INITIALIZE)
:   Initializes this module.
:
: (SHORTEST-PATH.COMPUTE)
:   Constructs the shortest-path tables using a transitive closure
:   algorithm.
:
: (SHORTEST-PATH.PRINT)
:   Prints out the table of shortest-path delays between all MEs.
:
: (ME:ME:DELAY ME1 ME2)
:   Returns the delay in getting a value from the outputs of ME1 to
:   the inputs of ME2. The delay is the number of cycles required for
:   the move and is currently equivalent to the number of register banks
:   along the path between the outputs of ME1 and the inputs of ME2.
:   Either ME1 or ME2 may be (), in which case the result is 0.
:
: (ME:ME:FIRST-ON-PATH ME1 ME2)
:   Returns the list of MEs that are first on the shortest path from
:   ME1 to ME2; () if ME1 = ME2.
:
:=====
(eval-when (compile load)
  (include list-scheduler:declarations) )

(declare (special
  *sp.me:me:path*
  *sp.me:me:delay*
) )

(defun shortest-path.initialize ()
  (:= *sp.me:me:path* ())
  (:= *sp.me:me:delay* ())
  ())

(defun shortest-path.compute ()
  (let* ((size (me-universe:size))
        (path (array:new '(,size ,size) ))
        (delay (array:new '(,size ,size) )))
    (:= *sp.me:me:path* path)
    (:= *sp.me:me:delay* delay)

    (loop (for-each-me me) (do
      (loop (for output-me in (me:outputs me) ) (do
        (:= ([])a delay (me:number me) (me:number output-me) )
        0)
        (push ([])a path (me:number me) (me:number output-me) )

```

```

output-me) ) ) )
(msg 0)
(loop (incr k from 0 to (- size 1) )
  (bind me-k (me-universe:number:me k) )
  (when (== 'register-bank (me:type me-k) ) )
  (do
    (msg " " k)

    (loop (incr i from 0 to (- size 1) )
      (when (!= 1 k) )
      (bind i-k-delay ([])a delay i k) )
      (when i-k-delay)

      (do
        (loop (incr j from 0 to (- size 1) )
          (when (!= j k) )
          (bind i-j-delay ([])a delay i j)
            k-j-delay ([])a delay k j) )
          (when k-j-delay)
          (bind i-k-j-delay (+ i (+ i-k-delay k-j-delay) ) )

          (do
            (?( (| (| (| i-j-delay)
              (< i-k-j-delay i-j-delay) )
              (:= ([])a delay i j) i-k-j-delay)
              (:= ([])a path i j) ([])a path i k) ) )
              (:= i-k-j-delay i-j-delay)
              (:= ([])a path i j)
              (unionq &&& ([])a path i k))))))))))

) ) )

(defun me:me:delay ( me1 me2 )
  (?( (| me1)
    0)
    ( (| me2)
    0)
    ( t
    ([])a *sp.me:me:delay* (me:number me1) (me:number me2) ) ) ) )

(defun me:me:first-on-path ( me1 me2 )
  (?( (| me1)
    () )
    ( (| me2)
    () )
    ( t
    ([])a *sp.me:me:path* (me:number me1) (me:number me2) ) ) ) )

(defun shortest-path.print ()
  (let ( (size (me-universe:size) ) )
    (loop (incr first-j from 0 to (- size 1) by 13)
      (bind last-j (min (+ first-j 12) (- size 1) ) )
      (do
        (msg 0 t)
        (let ( (names
              (loop (incr j from first-j to last-j) (save
                (sp.name:split (me:name (me-universe:number:me j) )
                t) ) ) ) )
          (msg 0 0)

```

```

(loop (for (top-name bottom-name) in names) (do
  (msg "| " (jc top-name -4) ) ) )
(msg 0 0)
(loop (for (top-name bottom-name) in names) (do
  (msg "| " (jc bottom-name -4) ) ) )

(msg 0)
(loop (incr i from 0 to (- size 1) )
  (bind me-i (me-universe:number:me i)
    (left-name right-name) (sp.name:split (me:name me-i) ()))
  (do
    (if (eqstr "" right-name) (then
      (msg (c left-name) (t i) ) )
      (else
        (msg (c left-name) "://" (c right-name) (t i) ) ) )
    (loop (incr j from first-j to last-j)
      (bind me-j (me-universe:number:me j)
        val (me:me:delay me-i me-j) )
      (do
        (if val (then
          (if (inump val) (then
            (msg (j val -4) " ") )
            (else
              (msg (f val 4 2) " ") ) ) ) )
          (else
            (msg (jc "_" -5) ) ) ) ) )
        (msg t) ) ) ) )
  ) )
) )

```

```

***=====
***
*** (SP.NAME:SPLIT NAME SPLIT-<=8?)
***
*** Auxiliary function used by LS.DELAY-TABLE:PRINT to produce pretty
*** truncated labels for the margins.
***
***=====

```

```

(defun sp.name:split ( name split-<=8? )
  (?( (<= (stlength name) 4)
    '(,name "" ) )
  ( (<= (stlength name) 8)
    (if split-<=8?
      '(,(substring name 1 4)
        ,(substring name 5) )
      '(,name "" ) ) )
  ( t
    '(,(substring name 1 4)
      ,(substring name -4) ) ) ) )

```

```

*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for
*** educational and research purposes only, by the faculty, students,
*** and staff of Yale University only.
*****
=====
: ELI MACHINE-MODEL
:
: A simple-minded, unrealistic ELI with 4 memory clusters, 4 floating
: clusters, almost all parameters hardwired.
:
=====

(alias '^^ 'atomconcat)

(defun em.make-memory-cluster (i n)
  (let ( (m1 (atomconcat 'm 1))
        (m1-1 (atomconcat 'm (mod (- 1 1) n)))
        (m1+1 (atomconcat 'm (mod (+ 1 1) n)))
        (f1 (atomconcat 'f 1))
        (f1-1 (atomconcat 'f (mod (- 1 1) n))) )
    ( (def-me (name .(^^ m1 'r) )
          (type register-bank)
          (size 32)
          (inputs (.(^^ m1 '+)
                  .(^^ m1 'm)
                  .(^^ m1 'c)
                  .(^^ m1 'r)
                  .(^^ m1-1 'r)
                  .(^^ m1+1 'r)
                  .(^^ f1-1 'r)
                  .(^^ f1 'r) ) )
          (read-ports 5)
          (write-ports 4)
          (read-resources (.(^^ m1 'r-read) ) )
          (write-resources (.(^^ m1 'r-write) ) ) )

      (def-resource-class .(^^ m1 'r-read) 5)
      (def-resource-class .(^^ m1 'r-write) 4)

      (def-me (name .(^^ m1 '+) )
          (type functional-unit)
          (delay 0)
          (inputs (.(^^ m1 'r)
                  .(^^ m1 'c) ) )
          (resources (.(^^ m1 '+) ) )
          (operators (if-ieq if-ine if-igt if-ige if-ile if-ilt
                     if-ie0mod if-true if-false iabs iadd iand
                     idiv ieq iexp ie0mod ige igt ile ilt imax
                     imin imul ine ineg inot ior isub bitrev) ) )

      (def-resource-class .(^^ m1 '+) 1)

      (def-me (name .(^^ m1 'm) )
          (type functional-unit)
          (delay 2)
          (inputs (.(^^ m1 'r) ) )
          (resources (.(^^ m1 'm) ) )
          (operators (vbase ivload fvload ipload fpload ivstore
                     fvstore ipstore fpstore) ) )
    )
  )
)

```

1

PS:&lt;C.S.BULLDOG.LIST-SCHEDULER.TEST&gt;SIMPLE-ELI-MODEL.LSP.3

```

(def-resource-class .(^^ m1 'm) 1)

(def-me (name .(^^ m1 'c) )
  (type constant-generator)
  (resources (.(^^ m1 'c) ) )
  (constraint-function
    (lambda (constant)
      (|| (consp constant)
          (inusp constant) ) ) ) )

(def-resource-class .(^^ m1 'c) 1)

(def-bus .(^^ m1-1 'r) .(^^ m1 'r) 1)
(def-bus .(^^ m1+1 'r) .(^^ m1 'r) 1)
(def-bus .(^^ f1-1 'r) .(^^ m1 'r) 1)
(def-bus .(^^ f1 'r) .(^^ m1 'r) 1)
)

(defun em.make-floating-cluster (i n)
  (let ( (m1 (atomconcat 'm 1))
        (m1+1 (atomconcat 'm (mod (+ 1 1) n)))
        (f1 (atomconcat 'f 1)) )
    ( (def-me (name .(^^ f1 'r) )
          (type register-bank)
          (size 32)
          (inputs (.(^^ f1 '+)
                  .(^^ f1 '=)
                  .(^^ f1 '*')
                  .(^^ f1 'r)
                  .(^^ m1 'r)
                  .(^^ m1+1 'r) ) )
          (read-ports 4)
          (write-ports 3)
          (read-resources (.(^^ f1 'r-read) ) )
          (write-resources (.(^^ f1 'r-write) ) ) )

      (def-resource-class .(^^ f1 'r-read) 4)
      (def-resource-class .(^^ f1 'r-write) 3)

      (def-me (name .(^^ f1 '+) )
          (type functional-unit)
          (delay 1)
          (inputs (.(^^ f1 'r) ) )
          (resources (.(^^ f1 '+) ) )
          (operators (fabs fadd feq fgt fgt fix file float flt
                     fmax fmin fne fneg fsub) ) )

      (def-me (name .(^^ f1 '=) )
          (type functional-unit)
          (delay 0)
          (inputs (.(^^ f1 'r) ) )
          (resources (.(^^ f1 '+) ) )
          (operators (if-feq if-fne if-fgt if-fge if-file if-flt) ) )

      (def-resource-class .(^^ f1 '+) 1)

      (def-me (name .(^^ f1 '*') )
          (type functional-unit)
          (delay 2)
          (inputs (.(^^ f1 'r) ) )
          (resources (.(^^ f1 '*') ) )
    )
  )
)

```

2

```
(operators (faul rdiv fexp cos sin sqrt tan) ) )
```

```
(def-resource-class ,(^^ fi '+) 1)
```

```
(def-bus ,(^^ m1 'r) ,(^^ fi 'r) 1)  
(def-bus ,(^^ m1+1 'r) ,(^^ fi 'r) 1)  
) )
```

```
(eval
```

```
  (def-machine-model
```

```
    ..(loop (incr i from 0 to 8) (splice  
      (em.make-memory-cluster i 4) ) )
```

```
    ..(loop (incr i from 0 to 8) (splice  
      (em.make-floating-cluster i 4) ) )  
  ) )
```

```

*****
*** Copyright (C) 1983 John R. Ellis. This file may be used for *****
*** educational and research purposes only, by the faculty, students, *****
*** and staff of Yale University only. *****
*****
=====
: MACHINE MODEL SIMULATOR
:
: This module implements the program simulator for the list-scheduler machine
: model.
:
: A program is a list of instructions. An instruction is a list of
: operations. An operation has the form:
:
: (me operator . operands)
:
: ME is the name of some machine element, OPERATOR is the name of the
: operator to be executed, and OPERANDS are the operand specifiers.
:
: An operand has one of two forms:
:
: (register-bank-me register name) to access a register;
:
: (constant-generator-me name) to generate a constant;
:
: NAME is the name of the value in the register or being generated. For
: registers, it is just a hint about what the symbolic contents of the
: register should be (the simulator remembers and checks those hints).
: For constant generators, NAME is the value that should be generated.
:
: A value name can be:
:
: a simple variable, e.g. X
:
: a number, e.g. 3.0
:
: a VBASE expression, e.g. (VBASE V)
:
: an address expression of the form (ADDRESS x) where "x" is any of the
: above name forms.
:
: A sample operation:
:
: (MO+ IADD (MOR 2 X) (MOR 4 (VBASE V)) (MOC 1))
:
: A continuation operation (a "pusher") for a multi-cycle pipeline is
: represented as:
:
: (FS+)
:
: A pseudo-operation doesn't a () ME:
:
: ( () LABEL L38)
: (NIL DCL %F%I FLOAT 256 ((1 256)) NIL)
:
:-----
: (SIMULATE PROGRAM ACTUALS)
: Runs the simulator on PROGRAM with actual parameters ACTUALS (specified
: the same as for the NADDR interpreter).
:
: (SIM.PRINT-PROGRAM-VARIABLES VAR-SPECS LOOK-IN-REGISTERS?)
: VAR-SPECS is a list of program variables; prints out each variable

```

```

: in turn. Each spec in VAR-SPECS is either an atom (for a vector
: or scalar) or a triple (VAR REAL IMAG) where VAR is the tiny lisp
: name of a complex variable and REAL and IMAG are the NADDR variables
: used to implement the real and imaginary parts. If
: LOOK-IN-REGISTERS? is true, then the contents of any registers
: containing the variables are printed out as well; otherwise only
: memory is examined.
:
: *SIM.OPERATION-COUNT*
: *SIM.INSTRUCTION-COUNT*
: Total number of operations and instructions executed, excluding
: pseudo-ops and instructions consisting only of pseudo-ops.
:
: (SVAR 'VAR1 'VAR2 ...)
: Interactive syntax for printing out variables, equivalent to:
:
: (SIM.PRINT-PROGRAM-VARIABLES '(VAR1 VAR2 ...) T)
:
: which will look both in memory and in registers for the variables.
:
: #R ME-NAME REGISTER
: Interactive syntax for getting the contents of a register.
:
:-----
: (eval-when (compile load)
: (include list-scheduler:declarations) )
: (eval-when (compile)
: (build '(utilities:options) ) )
:
: (def-option *sim.trace-instructions?* () list-scheduler: "
: If non-(), then instructions will be printed out as they are executed.
: If T, tracing begins with the first cycle. If some number, then tracing
: begins with that cycle.
: ")
:
: (def-option *sim.step-instructions?* () list-scheduler: "
: This is just like *SIM.TRACE-INSTRUCTIONS?*, except that a break-point
: is invoked after each instruction traced.
: ")
:
: (def-option *sim.trace-values?* () list-scheduler: "
: If non-(), then values being stored in registers and memory are printed
: out after the cycle in which the store actually occurs (i.e. at the
: end of pipe execution.
: ")
:
: (def-option *sim.break-label* () list-scheduler: "
: If non-(), then the simulator will stop with a breakpoint if a label of
: of this name is reached.
: ")
:
: (declare (special
: *sim.pc* ;*** These globals are mainly for debugging --
: *sim.running?* ;*** when a breakpoint occurs, you can look at
: *sim.instr* ;*** the current cycle, instruction, and
: *sim.oper* ;*** operation.
: *sim.cycle*

```

```

*sim.start-pc*      ;*** The pc to start execution at (where the
                    ;*** DEF-BLOCK is).

*sim.significant-instruction?*
                    ;*** True if the current instruction
                    ;*** contained a non-pseudo-op.
*sim.operation-count* ;*** # operations initiated this cycle.
*sim.copy-count*    ;*** # COPYs executed this cycle.

*sim.def-block*
) )

(defun simulate ( program actuals )
  (registers.initialize)

  (sim.initialize-memory)
  (sim.initialize-variables)
  (sim.initialize-labels)
  (sim.initialize-pipelines)

  (:= *sim.def-block* ( ) )
  (sim.program:declare program)

  (if *sim.def-block* (then
    (sim.formals:actuals:bind (oper:part *sim.def-block* 'in-variables)
      actuals) ) )

  (sim.program:execute *sim.start-pc*)

  (sim.print-program-variables (oper:part *sim.def-block* 'out-variables)
    ( ) )

  ( ) )

(defun-inline sim.trace? (function ( trace-var )
  (& trace-var
    (|| (! (inup trace-var) )
      (<= trace-var *sim.cycle* ) ) ) ) )

;*****
;***
;*** MEMORY
;***
;*** The simulator memory is represented as a vector, of course, from
;*** which space for variables, vectors, and constants is allocated.
;*** Memory is initialized to all zeroes, so that we can index past the
;*** end of an array.
;***
;*** (SIM.INITIALIZE-MEMORY)
;***   Initializes the memory to a default size.
;***
;*** (SIM.MEMORY:[ ] ADDR)
;***   Returns the contents of the ADDRth word.
;***
;*** (SIM.MEMORY:ALLOCATE LENGTH)
;***   Allocates a chunk of LENGTH words from memory, returning the
;***   address of the first word. Memory is grown if its current size
;***   is too small.
;***
;*** (SIM.ADDRESS:COERCE ADDR)
;***   Coerces a memory address to be an integer between 0 and the size

```

```

;*** of memory - 1. We do this because it's perfectly ok for compacted
;*** programs to generate bogus indices when they're doing nested
;*** vector references (A (L I) J) (as in SOLVE). It would be better
;*** to detect vector-out-of-bounds and return 0, but we don't yet
;*** have the hooks in generated code to do that easily (no vector
;*** names).
;***
;*****
(declare (special
  *sim.memory*
  *sim.next-free-address*
) )

(defun sim.initialize-memory ( )
  (:= *sim.memory* (makevector 1000) )
  (vector:initialize *sim.memory* 0)
  (:= *sim.next-free-address* 0)
  ( ) )

(defmacro sim.memory:[ ] ( addr )
  '([ ] *sim.memory* ,addr) )

(defun sim.memory:allocate ( length )
  (let ( (base *sim.next-free-address*)
    (max-length (vectorlength *sim.memory*)) )
    (:= *sim.next-free-address* (+ &&& length) )
    (if (> *sim.next-free-address* max-length) (then
      (:= *sim.memory* (vector:copy &&& (* 2 max-length) ) )
      (loop (incr 1 from max-length to (- (vectorlength *sim.memory*) 1))
        (do
          (:= ([ ] *sim.memory* i) 0) ) ) ) )
      base) )

(defun sim.address:coerce ( addr )
  (if (&& (inup addr)
    (>= addr 0)
    (< addr *sim.next-free-address*))
    (then
      addr)
    (else
      0) ) )

;*****
;***
;*** VARIABLES
;***
;*** Information about variable names and constants is kept in a hash
;*** table mapping the name or constant onto a record of type SIM.VD which
;*** describes the variable or constant. The only variables actually
;*** stored in memory are input and output variables and vectors.
;*** Constants are stored in memory only if they are contained in an
;*** (ADDRESS constant) expression.
;***
;*** (SIM.INITIALIZE-VARIABLES)
;***   Forgets all previous variable information.
;***

```



```

*** (SIM.VARIABLE:DECLARE NAME LENGTH DIMENSIONS INITIAL-VALUES )
*** Declares NAME (a constant or variable), where LENGTH, DIMENSIONS,
*** and INITIAL-VALUES are as specified on a DCL pseudo-op. Space
*** is allocated in memory for the variable or constant and
*** initialized to the initial values.
***
*** (SIM.VARIABLE:INITIALIZE VARIABLE INITIAL-LIST )
*** VARIABLE should already have been declared. Initializes its
*** memory with the values taken from INITIAL-LIST; if the list is
*** () the variable is initialized to all zeroes.
***
*** (SIM.VARIABLE:VD          VARIABLE)
*** (SIM.VARIABLE:BASE      VARIABLE)
*** (SIM.VARIABLE:LENGTH    VARIABLE)
*** (SIM.VARIABLE:DIMENSIONS VARIABLE)
*** (SIM.VARIABLE:RANK      VARIABLE)
*** These functions return the declared information about a variable.
*** An error is raised if the variable isn't declared.
***
=====
(declare (special
  *sim.variable:vd* ;*** Hash table mapping names to SIM.VD records.
) )

(def-struct sim.vd ;*** A descriptor for a variable.
  name ;*** Its name.
  base ;*** The address of its first word.
  length ;*** Its length.
  dimensions ;*** Its dimensions list, () if a scalar.
  rank ;*** The number of dimensions.
)

(defun sim.initialize-variables ()
  (:= *sim.variable:vd* (hash-table:create 'eql () ( ) ) )
  ( ) )

(defun sim.variable:vd ( variable )
  (|| (||h *sim.variable:vd* variable)
    (error (list variable "SIMULATOR: Variable isn't declared.")) ) )

(defun sim.variable:base ( variable )
  (sim.vd:base (sim.variable:vd variable) ) )

(defun sim.variable:length ( variable )
  (sim.vd:length (sim.variable:vd variable) ) )

(defun sim.variable:dimensions ( variable )
  (sim.vd:dimensions (sim.variable:vd variable) ) )

(defun sim.variable:rank ( variable )
  (sim.vd:rank (sim.variable:vd variable) ) )

(defun sim.variable:declare ( name length dimensions initial-values )
  (if-let ( (nd (||h *sim.variable:vd* name) ) ) (then
    (assert (& (== length (sim.vd:length nd) )
              (== dimensions (sim.vd:dimensions nd) ) )
      (h name) " " (h length) " " (h dimensions) " " t
      (h nd) t
    )
  )
  )

```

```

"SIMULATOR: Name previously declared with other attributes.") )
(else
  (let ( (nd (sim.vd:new
            name      name
            length    length
            dimensions dimensions
            rank      (length dimensions)
            base      (sim.memory:allocate length) ) ) )
    (:= (||h *sim.variable:vd* name) nd)
    (sim.variable:initialize name initial-values) ) )
  ( ) )

(defun sim.variable:initialize ( variable initial-list )
  (?( (! initial-list)
    (:= initial-list '(0) ) )
    ( (! (consp initial-list) )
      (:= initial-list (list initial-list) ) ) )

  (loop (incr i from (sim.variable:base variable)
        to (+ -i (+ (sim.variable:base variable)
                    (sim.variable:length variable) ) ) )
    (initial value
      0
      rest-initial-list initial-list)

  (do
    (if (! rest-initial-list)
      (:= rest-initial-list initial-list) )
    (:= value (pop rest-initial-list) )
    (:= (sim.memory:[] i) value) ) )
  ( ) )

=====
***
*** LABELS
***
*** Program labels are represented as Lisp symbols. A hash table maps the
*** symbols onto the corresponding program locations. A program location
*** is just some tail of the list of instructions representing the program.
***
*** (SIM.LABELS:INITIALIZE)
*** Forgets all information about labels.
***
*** (SIM.LABEL:DECLARE LABEL PC)
*** Declares LABEL to be at program location PC.
***
*** (SIM.LABEL:PC LABEL)
*** Returns the program location of a label.
***
=====
(declare (special
  *sim.label:pc*
) )

(defun sim.initialize-labels ()
  (:= *sim.label:pc* (hash-table:create () ( ) ( ) ) )
  ( ) )

(defun sim.label:declare ( label pc )

```

```

(assert (not (andp (h *sim.label:pc* label)
                  (h label) t "SIMULATOR: Label already declared."))
        (:= (andp (h *sim.label:pc* label) pc)
            ()
            )
)

(defun sim.label:pc (label)
  (if (andp (h *sim.label:pc* label)
            (error (list label "SIMULATOR: Label not declared.")))
      )
)

;=====
;***
;*** OPERANDS
;***
;*** See above for how operands are represented in the program.
;***
;*** (SIM.OPERAND:DECLARE OPERAND)
;*** Declares any constant mentioned in the OPERAND, allocating and
;*** initializing a memory location to contain it. Only constants
;*** within (ADDRESS ...) expressions in the operand need be so
;*** allocated.
;***
;*** (SIM.OPERAND:VALUE OPERAND)
;*** Returns the value of an operand, which is the contents of the
;*** specified register or the constant generated. For register
;*** operands, the name mentioned in the operand is checked against
;*** the symbolic name contained in the register; if they're not the
;*** same, an error is raised.
;***
;=====

(defun sim.operand:declare (operand)
  (let ((cg name) operand)
    (if (& (= 2 (length operand))
        (= 'constant-generator (me:type (name:me cg)))
        (consp name)
        (= 'address (car name)))
      (then
        (let ((address item) name)
          (if (litatom item)
              (sim.variable:declare item 1 () ())
              (numberp item)
              (sim.variable:declare item 1 () item)
              (& (consp item)
                  (= 'vbase (car item)))
              (let ((vd (sim.variable:vd (cadr item)))
                    (assert (< 0 (sim.vd:rank vd))
                           (h item) t "SIMULATOR: Invalid VBASE name.")
                    (sim.variable:declare item 1 () (sim.vd:base vd))))
                (t
                 (error (list name "SIMULATOR: Invalid operand."))))
            )
          )
      )
  )

(defun sim.operand:value (operand)
  (let* ((me-name . args) operand)
    (me (name:me me-name))

    (caseq (me:type me)
      (constant-generator
       (let ((constant) args)
         (if (consp constant) (then

```

```

(caseq (car constant)
  ((vbase address)
   (sim.variable:base (cadr constant)))
  (t
   (error (list operand "Case error."))))
  (else
   constant)))
(register-bank
 (let ((register name) args)
   (if (not (name (me:value-names me) register)) (then
     (msg (h operand) t
          "SIMULATOR: Name of value in register is different "
          "from name of operand.")
       (break-point simulator)))
     (me:values me) register)))
(t
 (error (list operand "SIMULATOR: Invalid operand."))))

;=====
;***
;*** PIPELINES
;***
;*** The results from functional unit pipelines are not written into their
;*** destinations until after the cycle from which they emerge from the
;*** pipe. To implement this behavior, we keep two priority queues, one
;*** for registers and one for memory. Entries in each queue are
;*** destination/value pairs, keyed by the cycle in which the value is
;*** to be stored into the destination. When an operation is executed,
;*** it's result is entered into the appropriate priority queue with the
;*** cycle it should be stored. At the end of each instruction cycle,
;*** the simulator stores away all the results queued for that cycle.
;***
;*** (SIM.INITIALIZE-PIPELINES)
;*** Initializes the queues representing the pipelines.
;***
;*** (SIM.DELAY:REGISTER-OPERAND:VALUE:STORE DELAY OPERAND VALUE)
;*** Remembers the fact that VALUE should be stored into the register
;*** destination described by OPERAND (see above) at the end of the
;*** cycle which is DELAY cycles after the current cycle.
;***
;*** (SIM.DELAY:ADDRESS:VALUE:STORE DELAY ADDRESS VALUE)
;*** Remembers the fact that VALUE should be stored in memory location
;*** ADDRESS at the end of the cycle which is DELAY cycles after
;*** the current cycle.
;***
;*** (SIM.STORE-STORED-RESULTS)
;*** Removes all entries from the two priority queues for the current
;*** cycle and stores their values in the given destinations.
;***
;=====

(declare (special
  *sim.saved-register-cycles*
  *sim.saved-register-values*
  *sim.saved-register-operands*
  *sim.next-saved-register-slot*

  *sim.saved-memory-cycles*
  *sim.saved-memory-values*
  *sim.saved-memory-addresses*
  *sim.next-saved-memory-slot*
  )
)

```

```

(defun sim.initialize-pipelines ()
  (:= *sim.saved-register-cycles* (makevector 100) )
  (:= *sim.saved-register-values* (makevector 100) )
  (:= *sim.saved-register-operands* (makevector 100) )
  (:= *sim.next-saved-register-slot* 0)

  (:= *sim.saved-memory-cycles* (makevector 100) )
  (:= *sim.saved-memory-values* (makevector 100) )
  (:= *sim.saved-memory-addresses* (makevector 100) )
  (:= *sim.next-saved-memory-slot* 0)
  () )

(defun sim.delay:register-operand:value:store ( delay operand value )
  (let ( (cycle (+ *sim.cycle* delay) ) )
    (:= ( [] *sim.saved-register-cycles* *sim.next-saved-register-slot* )
        cycle)
    (:= ( [] *sim.saved-register-operands* *sim.next-saved-register-slot* )
        operand)
    (:= ( [] *sim.saved-register-values* *sim.next-saved-register-slot* )
        value)
    (:= *sim.next-saved-register-slot* (+ &2 1) )
    () ) )

(defun sim.delay:address:value:store ( delay address value )
  (let ( (cycle (+ *sim.cycle* delay) ) )
    (:= ( [] *sim.saved-memory-cycles* *sim.next-saved-memory-slot* )
        cycle)
    (:= ( [] *sim.saved-memory-addresses* *sim.next-saved-memory-slot* )
        address)
    (:= ( [] *sim.saved-memory-values* *sim.next-saved-memory-slot* )
        value)
    (:= *sim.next-saved-memory-slot* (+ &2 1) )
    () ) )

(defun sim.store-saved-results ()
  (loop (initial next-slot 0)
    (incr i from 0 to (- *sim.next-saved-register-slot* 1) )
    (bind cycle ( [] *sim.saved-register-cycles* 1)
      operand ( [] *sim.saved-register-operands* 1)
      value ( [] *sim.saved-register-values* 1) )

    (do
      (if (== cycle *sim.cycle*) (then
        (let* ( ( (me-name register name) operand)
              ( (me (name:me me-name) ) )
              (:= ( [] (me:values me) register) value)
              (:= ( [] (me:value-names me) register) name)

              (if (sim.trace? *sim.trace-values?) (then
                (msg 0 (j operand -20) " = " value t) ) ) ) )
          (else
            (:= ( [] *sim.saved-register-cycles* next-slot) cycle)
            (:= ( [] *sim.saved-register-operands* next-slot) operand)
            (:= ( [] *sim.saved-register-values* next-slot) value)
            (:= next-slot (+ &2 1) ) ) ) )
      (result
        (:= *sim.next-saved-register-slot* next-slot) ) )
  )

```

```

(loop (initial next-slot 0)
  (incr i from 0 to (- *sim.next-saved-memory-slot* 1) )
  (bind cycle ( [] *sim.saved-memory-cycles* 1)
    address ( [] *sim.saved-memory-addresses* 1)
    value ( [] *sim.saved-memory-values* 1) )

  (do
    (if (== cycle *sim.cycle*) (then
      (:= (sim.memory:[] address) value)

      (if (sim.trace? *sim.trace-values?) (then
        (msg 0 (j address -20) " = " value t) ) ) )
      (else
        (:= ( [] *sim.saved-memory-cycles* next-slot) cycle)
        (:= ( [] *sim.saved-memory-addresses* next-slot) address)
        (:= ( [] *sim.saved-memory-values* next-slot) value)
        (:= next-slot (+ &2 1) ) ) ) )
    (result
      (:= *sim.next-saved-memory-slot* next-slot) ) )
  () )

;*****
;***
;*** (SIM.PROGRAM:DECLARE PROGRAM)
;***
;*** Scans all the operations in PROGRAM, declaring all the operands, allocating
;*** and initializing memory for constants, input and output variables, and
;*** vectors. Labels are found and declared, and the DEF-BLOCK is put
;*** in *SIM.DEF-BLOCK*.
;***
;*****

(defun sim.program:declare ( program )
  (:= *sim.largest-trace* 1)
  (:= *sim.start-pc* program)

  (loop (initial *sim.pc* program)
    (while *sim.pc*)
    (bind instr (car *sim.pc*) )

    (do
      (sim.instr:declare-pass-1 instr)
      (next *sim.pc* (cdr *sim.pc*) ) )

    (loop (initial *sim.pc* program)
      (while *sim.pc*)
      (bind instr (car *sim.pc*) )

      (do
        (sim.instr:declare-pass-2 instr)
        (next *sim.pc* (cdr *sim.pc*) ) )

      () )

  (defun sim.instr:declare-pass-1 ( instr )
    (loop (for oper in instr) (do
      (sim.oper:declare-pass-1 oper) )
    () )

  (defun sim.oper:declare-pass-1 ( oper )

```

```

(let* ( ( (me-name . naddr-oper) oper)
        ( (operator . operands) naddr-oper) )
  (? ( ! me-name)
    (caseq operator
      (def-block
        (:= *sim.def-block* '(operator . operands) )
        (:= *sim.start-pc* *sim.pc*) )
      (label
        (sim.label:declare (car operands) *sim.pc*) )
      (dcl
        (sim.variable:declare
          (oper:part naddr-oper 'variable)
          (oper:part naddr-oper 'length)
          (oper:part naddr-oper 'dimensions)
          (oper:part naddr-oper 'initial-list) ) )
      (trace
        (let ( ( (number . ()) operands)
              (:= *sim.largest-trace* (max && number))))))
    ) ) )

```

```

(defun sim.instr:declare-pass-2 ( instr )
  (loop (for oper in instr) (do
    (sim.oper:declare-pass-2 oper) )
  ) )

```

```

(defun sim.oper:declare-pass-2 ( oper )
  (let* ( ( (me-name . naddr-oper) oper)
          ( (operator . operands) naddr-oper) )
    (? ( ! me-name)
      () )

    ( ! naddr-oper)
    () )

    ( == 'copy operator)
    (sim.operand:declare (cadr operands) ) )

    ( t
      (caseq (operator:group operator)
        ( (one-in-one-out two-in-one-out vload)
          (loop (for operand in (cdr operands) ) (do
            (sim.operand:declare operand) ) ) )
        (vstore
          (loop (for operand in operands) (do
            (sim.operand:declare operand) ) ) )
        (if-compare
          (sim.operand:declare (car operands) )
          (sim.operand:declare (cadr operands) ) )
        (if-boolean
          (sim.operand:declare (car operands) ) ) ) )
    ) ) )

```

```

:***=====
:***
:*** (SIM.PROGRAM:EXECUTE START-PC)
:***
:*** This is the main execute loop, executing instructions until we run
:*** off the end of the program or we hit a STDP.
:***

```

```

:***=====
:***
(defun sim.program:execute ( start-pc )
  (sim.initialize-execution-records)

  (loop (incr *sim.cycle* from 0)
    (initial *sim.running?* t
            pc start-pc)
    (while *sim.running?*)
    (while pc)
    (do
      (:= *sim.pc* ())
      (sim.instr:execute (car pc) )
      (:= pc (|| *sim.pc* (cdr pc) ) ) ) )
  ) )

```

```

:***=====
:***
:*** (SIM.INSTR:EXECUTE *SIM.INSTR*)
:***
:*** Executes a single instruction.
:***
:***=====

```

```

(defun sim.instr:execute ( *sim.instr* )
  (:= *sim.significant-instruction?* () )
  (:= *sim.operation-count* 0)
  (:= *sim.copy-count* 0)

  (if (|| (sim.trace? *sim.trace-instructions?*)
          (sim.trace? *sim.step-instructions?*) )
    (then
      (msg 0 (j *sim.cycle* 4) " : " (h *sim.instr*) t)
      (if (sim.trace? *sim.step-instructions?*) (then
        (break-point single-step) ) ) ) )

  (loop (for oper in *sim.instr*) (do
    (sim.oper:execute oper) ) )
  (sim.store-saved-results)

  (sim.record-instruction-execution)
  () )

```

```

:***=====
:***
:*** (SIM.OPER:EXECUTE *SIM.OPER*)
:***
:*** Executes a single operation by dispatching according to type:
:*** pseudo-operation, register-bank operation (a copy), or functional-unit
:*** operation.
:***
:***=====

```

```

(defun sim.oper:execute ( *sim.oper* )
  (let* ( ( (me-name . ()) *sim.oper* )
          (if ( ! me-name) (then
            (sim.pseudo-op:execute *sim.oper*) )
          (else
            (let ( (me (name:me me-name) ) )
              (caseq (me:type me)
                (register-bank

```

```

      (sim.register-bank-me:oper:execute me *sim.oper*) )
    (functional-unit
      (sim.functional-unit-me:oper:execute me *sim.oper*) )
    (t
      (error (list *sim.oper*
                  "SIMULATOR: Invalid ME in machine operation."))))
  ) )
) )

```

```

=====
***
*** (SIM.PSEUDO-OP:EXECUTE OPER)
***
*** Executes a pseudo-operation.
***
=====

```

```

(defun sim.pseudo-op:execute ( oper )
  (let ( ( () operator . operands) oper )
    (caseq operator
      (label
        (if (== (car operands) *sim.break-label*) (then
              (break-point simulator-label-break) ) ) )
      (trace
        (let ( ( (number type . ()) operands) )
          (:= *sim.current-trace* number)
          (:= *sim.current-trace-type* type) ) )
      (goto
        (if (! *sim.pc*) (then
              (:= *sim.pc* (sim.label:pc (car operands) ) ) ) ) )
      (stop
        (:= *sim.running?* () ) )
      (esc
        (eval '(progn .operands) ) ) )
    ) ) )
) )

```

```

=====
***
*** (SIM.REGISTER-BANK-ME:OPER:EXECUTE ME OPER)
***
*** Executes a register bank operation for register bank ME. The only
*** defined operation for register banks is COPY.
***
=====

```

```

(defun sim.register-bank-me:oper:execute ( me ( () copy dest source) )
  (assert (== 'copy copy)
    (h *sim.oper* "SIMULATOR: Unknown register-bank operator.") )
  (sim.delay:register-operand:value:store
    0
    dest
    (sim.operand:value source) )

  (:= *sim.copy-count* (+ &&& 1) )
  (:= *sim.significant-instruction?* t)
  )
)

```

```

=====
***
*** (SIM.FUNCTIONAL-UNIT-ME:OPER:EXECUTE ME OPER)

```

```

***
*** Executes a functional unit operation for functional unit ME. We
*** cheat here by invoking the private interfaces to the NADDR interpreter
*** to do most of our work. This is the right approach, but unfortunately
*** those interfaces are currently private and undocumented. Gross.
***
=====

```

```

(defun sim.functional-unit-me:oper:execute ( me ( () operator . operands) )
  (:= *sim.significant-instruction?* t)

  (if (! operator) (then
    () )
    (else
      (:= *sim.operation-count* (+ &&& 1) )

      (let* ( (naddr.operator (operator:naddr.operator operator) )
              (delay (me:delay me) )
              (function (naddr.operator:execute-function
                          naddr.operator) ) )

        (caseq (naddr.operator:group naddr.operator)
          (one-in-one-out
            (sim.delay:register-operand:value:store
              delay
              (car operands)
              (funcall function
                (sim.operand:value (cadr operands))))))

          (two-in-one-out
            (sim.delay:register-operand:value:store
              delay
              (car operands)
              (funcall function
                (sim.operand:value (cadr operands) )
                (sim.operand:value (caddr operands) ) ) ) )

          (vload
            (caseq operator
              ( (ivload fvload)
                (let ( ( (dest vector index) operands) )
                  (sim.delay:register-operand:value:store
                    delay
                    dest
                    (sim.memory:[]
                      (sim.address:coerce
                        (+ (sim.operand:value vector)
                          (sim.operand:value index))))))

                  ( (ipload fpload)
                    (let ( ( (dest index) operands) )
                      (sim.delay:register-operand:value:store
                        delay
                        dest
                        (sim.memory:[]
                          (sim.address:coerce
                            (sim.operand:value index))))))

                    (t
                      (error (list *sim.oper* "Case error.") ) ) ) )

                (vstore
                  (caseq operator
                    ( (ivstore fvstore)

```

```

      (let ( ( (vector index source) operands) )
        (sim.delay:address:value:store
          0 ;*** correct.
          (+ (sim.operand:value vector)
             (sim.operand:value index) )
            (sim.operand:value source) ) ) )

      ( (ipstore fpstore)
        (let ( ( (index source) operands) )
          (sim.delay:address:value:store
            0 ;*** correct.
            (sim.operand:value index)
            (sim.operand:value source) ) ) )

      (t
        (error (list *sim.oper* "Case error.") ) ) ) )

      (if-compare
        (if (! *sim.pc*) (then
          (let* ( ( (operand1 operand2 label1 label2) operands)
                (label (funcall function
                          (sim.operand:value operand1)
                          (sim.operand:value operand2)
                          label1
                          label2) ) )
            (if label (then
              (:= *sim.pc* (sim.label:pc label)))))))))

      (if-boolean
        (if (! *sim.pc*) (then
          (let* ( ( (operand1 label1 label2) operands)
                (label (funcall function
                          (sim.operand:value operand1)
                          label1
                          label2) ) )
            (if label (then
              (:= *sim.pc* (sim.label:pc label)))))))))

      (t
        (error (list *sim.oper* "Case error.") ) ) ) ) )
  ) )

```

```

;***
;***
;*** (SIM.FORMALS:ACTUALS:BIND FORMALS ACTUALS)
;***
;*** Binds the actual parameters to the formal parameters of a program.
;*** FORMALS is the list of formal parameters as given in DEF-BLOCK.
;*** ACTUALS is the list of actual parameters as given to INTERPRET.
;***
;***
=====
(defun sim.formals:actuals:bind ( formals actuals )
  (assert (listp formals) )
  (assert (listp actuals) )

  (loop (for formal in formals)
        (for actual in actuals)
        (do
          (? ;*** A non-complex scalar or array?
            ;***

```

```

  ( (litatom formal)
    (sim.variable:initialize formal actual) )

    ;*** A complex scalar?
    ;***
    ( (& (consp formal)
      (== 3 (length formal) )
      (for-every (var in formal) (litatom var) )
      (consp actual)
      (== 2 (length actual) )
      (numberp (car actual) )
      (numberp (cadr actual) ) )
      (sim.variable:initialize (cadr formal) (car actual) )
      (sim.variable:initialize (caddr formal) (cadr actual) ) )

    ;*** A complex array?
    ;***
    ( (& (consp formal)
      (== 3 (length formal) )
      (for-every (var in formal) (litatom var) )
      (consp actual)
      (for-every (pair in actual)
        (& (consp pair)
          (numberp (car pair) )
          (numberp (cadr pair) ) ) ) )
      (sim.variable:initialize (cadr formal)
        (loop (for (real imag) in actual) (save real) ) )
      (sim.variable:initialize (caddr formal)
        (loop (for (real imag) in actual) (save imag) ) ) )

    ( t
      (error (list formal actual
        "SIMULATOR: Invalid formal/actual pair.") ) ) ) ) )
  ) )

```

```

;***
;***
;*** EXECUTION RECORDS
;***
;*** A SIM.ER records the execution statistics of a trace (or the entire
;*** program).
;***
;***
=====
(def-struct sim.er ;***
  (number 0) ;*** Number of the trace
  ;***
  (entries 0) ;*** # times the trace was entered.
  ;***
  (instructions 0) ;*** # instructions executed.
  ;***
  (operations 0) ;*** # operations initiating a functional unit.
  ;***
  (copies 0) ;*** # COPY operations.
  ;***
  (ps-instructions 0) ;*** The fields keep the same counts for split- and
  (ps-operations 0) ;*** and join- partial schedules; these counts are
  (pj-instructions 0) ;*** included in the totals above.
  (pj-operations 0) ;***
  ) ;***
;***
;***
=====

```

```

***
*** (SIM.INITIALIZE-EXECUTION-RECORDS)
***   Initializes the statistics recording.
***
*** (SIM.RECORD-INSTRUCTION-EXECUTION)
***   This is called after executing an instruction. It just bumps the
***   execution totals for the current trace, using the global variables
***   *SIM.OPERATION-COUNT* and *SIM.COPY-COUNT*.
***
*** (SIMULATOR.PRINT-EXECUTION-STATISTICS)
***   Dumps out all the execution statistics, including a per-trace profile.
***
=====
(declare (special
  *sim.all-ers*           ;*** All execution records.
  *sim.trace:er*         ;*** Vector mapping trace numbers to
                        ;*** execution records.
  *sim.largest-trace*    ;*** Largest trace number in the program.
  *sim.current-trace*    ;*** Current trace number.
  *sim.last-trace*       ;*** Previous trace number.
  *sim.current-trace-type* ;*** SPLIT, JOIN, or ().

) )

(defun sim.initialize-execution-records ()
  (:= *sim.current-trace* 1)
  (:= *sim.last-trace* ())
  (:= *sim.current-trace-type* ())

  (:= *sim.trace:er* (makevector (+ 1 *sim.largest-trace*)))

  (:= *sim.all-ers* ())
  (loop (incr i from 1 to *sim.largest-trace*) (do
    (:= ([] *sim.trace:er* i) (sim.er:new number i))
    (push *sim.all-ers* ([] *sim.trace:er* i))))
  ())

(defun sim.record-instruction-execution ()
  (if *sim.significant-instruction?* (then
    (let ( (er ([] *sim.trace:er* *sim.current-trace*)) )
      (:= (sim.er:instructions er) (+ &&& 1))
      (:= (sim.er:operations er) (+ &&& *sim.operation-count*))
      (:= (sim.er:copies er) (+ &&& *sim.copy-count*))

      (if (:= *sim.current-trace* *sim.last-trace*) (then
        (:= (sim.er:entries er) (+ &&& 1))))
      (:= *sim.last-trace* *sim.current-trace*)

      (case *sim.current-trace-type*
        (split
         (:= (sim.er:ps-instructions er)
            (+ &&& 1))
          (:= (sim.er:ps-operations er)
            (+ &&& *sim.operation-count*)) )
        (join
         (:= (sim.er:pj-instructions er)
            (+ &&& 1))
          (:= (sim.er:pj-operations er)
            (+ &&& *sim.operation-count*)) )
        ( ) )
    )
  )
)

```

```

( )
  (error (list *sim.current-trace-type* "Case error.*****)))

(defun simulator.print-execution-statistics ()
  (let ( (sorted-ers (sort *sim.all-ers*
                          (f:l (er1 er2)
                              (> (sim.er:instructions er1)
                                   (sim.er:instructions er2))))) )
    (t-er (sim.er:new number "Tot")))

    (loop (for er in *sim.all-ers*) (do
      (:= (sim.er:instructions t-er) (+ &&& (sim.er:instructions er)))
      (:= (sim.er:operations t-er) (+ &&& (sim.er:operations er)))
      (:= (sim.er:copies t-er) (+ &&& (sim.er:copies er)))
      (:= (sim.er:entries t-er) (+ &&& (sim.er:entries er)))
      (:= (sim.er:ps-instructions t-er) (+ &&& (sim.er:ps-instructions er)))
      (:= (sim.er:ps-operations t-er) (+ &&& (sim.er:ps-operations er)))
      (:= (sim.er:pj-instructions t-er) (+ &&& (sim.er:pj-instructions er)))
      (:= (sim.er:pj-operations t-er) (+ &&& (sim.er:pj-operations er)))
    ))

    (msg 0 t "EXECUTION PROFILE" t)
    (sim.er:print-heading)
    (msg (jc "-" 70 #/-) t)
    (sim.er:print t-er t-er)
    (msg (jc "-" 70 #/-) t)

    (loop (for er in sorted-ers)
      (incr i from 1)
      (when (> (sim.er:instructions er) 0) )

      (do
        (sim.er:print er t-er (if (== 0 (mod i 3)) #/ #/)))
      (msg (jc "-" 70 #/-) t)
    )
  )

(defun sim.er:print-heading ()
  (msg (jc "Tr" 3)
    (jc "Instrs" 7)
    (jc "#" 4)
    (jc "Ops" 7)
    (jc "#" 4)
    (jc "Copies" 7)
    (jc "#" 4)
    (jc "PS-I" 6)
    (jc "#I" 3)
    (jc "PJ-I" 6)
    (jc "#I" 3)
    (jc "Entries" 6)
    (jc "I//E" 7)
    (jc "O//I" 5)
    (jc "C//I" 5)
    t)
  ())

(defun sim.er:print (er t-er pc)
  (msg (jc (sim.er:number er)

```

```

(jc (sim.er:instructions er) 7 pc)
(jc (%r (sim.er:instructions er) (sim.er:instructions t-er) ) 4 pc)
(jc (sim.er:operations er) 7 pc)
(jc (%r (sim.er:operations er) (sim.er:operations t-er) ) 4 pc)
(jc (sim.er:copies er) 7 pc)
(jc (%r (sim.er:copies er) (sim.er:copies t-er) ) 4 pc)
(jc (sim.er:ps-instructions er) 8 pc)
(jc (%r (sim.er:ps-instructions er) (sim.er:instructions t-er) ) 3 pc)
(jc (sim.er:pj-instructions er) 8 pc)
(jc (%r (sim.er:pj-instructions er) (sim.er:instructions t-er) ) 8 pc)
(jc (sim.er:entries er) 8 pc)
(f //f (sim.er:instructions er) (sim.er:entries er) ) 7 i pc)
(f //f (sim.er:operations er) (sim.er:instructions er) ) 5 i pc)
(f //f (sim.er:copies er) (sim.er:instructions er) ) 5 i pc)
t)
) )

(defun % ( x y )
  (if (= y 0)
    0
    (// (+ x 100.0) y) ) )

(defun %r ( x y )
  (round (% x y) ) )

(defun //f ( x y )
  (if (= y 0)
    0
    (// (flonum x) y) ) )

;*****
;***
;*** VARIABLE PRINTING (see above).
;***
;*****

(defun sim.print-program-variables ( var-specs look-in-registers? )
  (loop (for var-spec in var-specs) (do
    (sim.var-spec:print var-spec look-in-registers?) ) )
  ) )

(defmacro svar var-specs
  '(sim.print-program-variables ',var-specs t) )

(defun sim.var-spec:print ( var-spec look-in-registers? )
  (? (& look-in-registers?
    (|| (! ([]h *sim.variable:vd* var-spec) )
      (= 0 (sim.variable:rank var-spec) ) ) )
    (sim.var:find-registers var-spec) )
  ( (consp var-spec)
    (sim.var:real:imag:print
      (car var-spec) (cadr var-spec) (caddr var-spec) ) )
  ( t
    (sim.var:real:imag:print var-spec var-spec ( ) ) ) )
  ) )

```

```

(defun sim.var:find-registers ( var )
  (loop (for me in *ls.register-bank-mes*
    (initial found-one? ( ) )
    (do
      (loop (incr i from 0 to (- (me:size me) 1) )
        (when (= var ([] (me:value-names me) i) ) )
        (do
          (if (! found-one?) (then
            (msg 0 "Variable " var " = " ) )
            (msg (t 20) " = " ([] (me:values me) i) " = "
              (t 35) '(,(me:name me) .i) t)
            (:= found-one? t) ) ) )
        (result
          (if (! found-one?) (then
            (msg 0 "Variable " var " not found.") ) ) ) )
      ) )
  ) )

(defun sim.var:real:imag:print ( var real imag )

  (let ( (real-base (sim.variable:base real) )
    (imag-base (if imag (sim.variable:base imag) ( ) ) ) )

    (if (! (sim.variable:dimensions real) ) (then
      (msg 0 "Variable " var " = "
        (sim.real:imag:offset:value real-base imag-base 0) t) )

      (else
        ;*** an array, print out each element of the array, labeled
        ;*** with its index.
        ;***
        (let ( (dimensions ( ) ) )

          ;*** First make a list of triples of the form (LOWER UPPER
          ;*** SIZE) corresponding to the dimensions of the array.
          ;***
          (loop (for (lower upper)
            in (reverse (sim.variable:dimensions real) ) )
            (initial size 1)
            (do
              (push dimensions '(,lower ,upper ,size) )
              (:= size (+ size (+ 1 (- upper lower) ) ) ) ) ) )

          ;*** Now print each element of the array, labeled with
          ;*** the index.
          ;***
          (msg "Variable " var " = " t)
          (loop (incr i from 0)
            (bind value (sim.real:imag:offset:value
              real-base imag-base i)
              last-i (- i 1) )
            (initial first-value '(unique-garbage)
              first-i ( )
              final-i (- (sim.variable:length real) i) )
            (when (|| (> i final-i)
              (& (!= value first-value) ) ) )
            (do
              (if first-i (then
                (msg " " )
                (if (> last-i first-i) (then
                  (sim.index:print first-i var dimensions)
                  (msg " :") ) )
                (sim.index:print last-i var dimensions)
              )
            )
          )
        )
      )
    )
  )

```



```

      (msg (t 18) " = " first-value t) ) )
    (if (> i final-1) (then
      (return () ) ) )
    (:= first-i i)
    (:= first-value value) ) ) ) )
  ( ) ) )

```

```

(defun sim.index:print ( i var dimensions )
  (loop (for (lower upper size) in dimensions)
    (initial remainder i)
    (do
      (msg " " (+ lower (/ remainder size) ) )
      (:= remainder (\ remainder size) ) ) ) )

```

```

(defun sim.real:imag:offset:value ( real-base imag-base offset )
  (if imag-base (then
    '( (sim.memory:[] (+ real-base offset) )
      (sim.memory:[] (+ imag-base offset) ) ) )
    (else
      (sim.memory:[] (+ real-base offset) ) ) ) )

```

```

(def-sharp-sharp r
  '(sim.sharp-sharp-r ',(read) ',(read) ) )

```

```

(defun sim.sharp-sharp-r ( me-name register )
  (let ( (me (name:me me-name) ) )
    '( ( [] (me:values me) register)
      ( [] (me:value-names me) register) ) ) )

```

```

***-----*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for *****
*** educational and research purposes only, by the faculty, students, *****
*** and staff of Yale University only. *****
***-----*****
=====
: TEST BED
:
: This module implements a test bed that allows the use of "trace
: snapshots" to simulate the compiler interface without actually requiring
: the compiler to be present.
:
: To use it, first create a trace snapshot as described in
: GENERATE-CODE-HOOK. Then, if the snapshot was for trace 3 of program
: PRIME, load the file PRIME-3.TRACE.
:
: (TLS NAME) [Test List Scheduler]
: Invokes the test bed on the variable "name" that was loaded from the
: file name.TRACE.
:
: The interface functions START-TRACE, PREDECESSORS, NAME:DATYPE, and
: NAME:RANK are defined here that simulate the real ones.
:
: To get back to the "real" interface, call (TEST-BED.UNDO)
:
:=====
(declare (special
  *tb.symbol-table* ;*** A symbol table (list) of tuples of the form:
                    ;***
                    ;*** (name datatype rank)

  *tb.index:datum* ;*** A vector filled in by PREDECESSORS that
                    ;*** maps integer trace positions of operations
                    ;*** onto the client-supplied datum passed
                    ;*** into PREDECESSORS.

  *tb.index:result* ;*** A vector mapping trace positions of
                    ;*** operations onto the original result
                    ;*** returned by PREDECESSORS (in the real
                    ;*** compiler).

  *tb.current-index* ;*** The trace position of the next expected
                    ;*** operation expected by PREDECESSORS. We
                    ;*** assume the client calls PREDECESSORS in
                    ;*** order.

  *tb.saved-names&defs* ;*** List of pairs of function names and
                    ;*** definitions for the original functions
                    ;*** simulated by this interface.

) )

(if (not (boundp '*tb.saved-names&defs*')) (then
  (:= *tb.saved-names&defs*
    (loop (for name in '(start-trace predecessors name:datatype name:rank))
      (save
        '(,name ,(fundef name) ) ) ) ) ) )

:=====
:***
:*** (TLS NAME)

```

1

PS:&lt;C.S.BULLDOG.LIST-SCHEDULER.TEST&gt;TEST-BED.LSP.17

```

:***
:***-----*****
(defun tls
  ( (symbol-table live-before source-record-list live-after
    predecessors-results)
    &optional (print? t) )

  ;*** Remember the symbol table.
  ;***
  (:= *tb.symbol-table* symbol-table)

  ;*** Initialize the vectors mapping trace positions.
  ;***
  (:= *tb.index:datum* (makevector (length source-record-list) ) )
  (:= *tb.index:result* (makevector (length source-record-list) ) )

  (vector:initialize *tb.index:result* predecessors-results)

  ;*** Initialize the starting index to 0, 1 if there is a DEF
  ;*** in the first position (the client doesn't call PREDECESSORS
  ;*** on DEF.
  ;***
  (:= *tb.current-index*
    (if (== 'def (operator (car (car source-record-list) ) ) )
      1
      0) )

  ;*** Now invoke whatever we want to invoke.
  ;***
  (generate-code live-before source-record-list live-after)

  (if print? (then
    (sch.print-schedule) ) )
  () )

:=====
:***
:*** (START-TRACE)
:***
:=====
(defun start-trace ()
  () )

:=====
:***
:*** (PREDECESSORS SOURCE-OPER TRACE-DIRECTION DATUM)
:***
:=====
(defun predecessors ( source-oper trace-direction datum )
  (let* ( (result
    (loop (for (index . reasons) in ([] *tb.index:result*
      *tb.current-index*)
      (save
        '(, ( ( [] *tb.index:datum* index)
          ,reasons) ) ) ) )
    (:= ( [] *tb.index:datum* *tb.current-index* datum)
      (:= *tb.current-index* (+ &@& 1) )

```

2

```

result) )

:*****
:***
:*** (NAME:DATATYPE NAME)
:***
:*****

(defun name:datatype ( name )
  (let ( ( ( () datatype rank) (assoc name *tb.symbol-table*) ) )
    (assert datatype "NAME:DATATYPE: Name not found: " (h name) )
    datatype ) )

(defun name:rank ( name )
  (let ( ( ( () datatype rank) (assoc name *tb.symbol-table*) ) )
    (assert rank "NAME:DATATYPE: Name not found: " (h name) )
    rank ) )

:*****
:***
:*** (TEST-BED.UNDO)
:***
:*****

(defun test-bed.undo ()
  (loop (for (name def) in *tb.saved-names&defs*) (do
    (:= (fundef name) def) ) )
  ) )

```

```

*****
*** Copyright (C) 1988 John R. Ellis. This file may be used for
*** educational and research purposes only, by the faculty, students,
*** and staff of Yale University only.
*****

```

#### VALUE NODES

A VN (Value Node) represents one value (and possibly an operation) in the expression DAG for a trace.

```

=====
(def-struct vn
  number      : A unique number; guaranteed to produce source
               : ordering of nodes in the trace.
  type        : One of {DEF, DEF1, OPERATION, USE, USE1, COPY,
               : PSEUDO-CP}
  constant?   : True if a constant DEF.
  datatype    : Datatype of this value, one of {INTEGER, FLOAT}
  datum       : The opaque bookkeepr datum handed to us by
               : GENERATE-CODE.
  name        : The variable name written by this node; for
               : constant DEFs, the value of the constant.
  oper        : Source MADDR operation this VN represents.
  operand-vns : List of VNs that are read by this VN.
  constraining-vns : List of VNs that must precede this VN, but which
               : are not values read by this node.
  constraining-delays : A list of numbers corresponding to
               : :CONSTRAINT-VNS; each number specifies how many
               : cycles later this VN must be done after the
               : corresponding constraining VN.
  off-live-vns : List of VNs "conditionally read" by this VN
               : on the off-trace edge (if it is a jump).
  reading-vns  : List of VNs that read the value of this VN.
  constrained-vns : List of VNs that are constrained by this VN (via
               : :CONSTRAINING-VNS).
  off-live-reading-vns : List of VNs (jumps) that "off-live read" this
               : VN.
  depth        : A lower bound on the earliest this VN could be
               : scheduled.
  height       : A lower bound on the "time" of this VN from the
               : exit of the DAG, analagous to the :DEPTH.
  me           : The functional unit or constant generator where this
               : VN will be computed.
)
=====

```

```

likely-mes    : List of MES, the current best guess as to where
               : this VN will be located or computed.
bug-cycle     : The cycle in which BUG guessed this VN could be
               : be scheduled.
cycle         : In list scheduling, the current earliest cycle this
               : VN can be scheduled.
scheduled-cycle : In list scheduling, the cycle this VN was actually
               : scheduled.
register-bank-me : The register location of this VN and the time
register       : it is available. For USEs, this is the desired
register-cycle : register location (if given).
read-cycle    : The last cycle in which the value of this VN
               : was read. This tells us the cycle in which
               : the value is available on a register-bank or
               : constant-generator "port".
readers-left  : The number of successors of this VN that read the
               : value (including off-live readers) that haven't
               : been scheduled yet.
predecessors-left : The number of predecessors of this VN (operands
               : and constraining VNs) that haven't been scheduled
               : yet.
locations     : For USEs and DEFs only, a list of pairs (ME
               : REGISTER) specifying the locations of the USE/DEF.
)
=====

```

VNs form their own OBJECT-UNIVERSE (see OBJECT-UNIVERSE.LSP).

```

(VN:INITIALIZE)
  Initializes this module, forgetting about all previous VNs.
(VN:CREATE VN)
  Records VN in the VN object universe. The idiom for making a
  new VN is:
      (VN:CREATE (VN:NEW ...) )
(VN:OPERATOR VN)
(VN:GROUP VN)
  Accessing functions for the MADDR operation of the VN.
(VN:DELAY VN)
  For operation VNs, the delay of the assigned ME (in :ME); 0 for
  all other VN types.
(VN:CONSTRAINED-VN:DELAY VN CONSTRAINED-VN)
  VN is on of the :CONSTRAINING-VNS of CONSTRAINEDS-VN. Returns the
  constraining delay between the two (if there are several delays
  between the two, the largest is returned).
(VN:INSERT-VNS VN PREDECESSOR-VNS)
  VN should have exactly one operand. The list PREDECESSOR-VNS is
  spliced between that operand and VN, so that each VN in
  PREDECESSOR-VNS now reads the operand, and VN now reads all of the

```

```

predecessors.
(VN:SPlice-VN VN NEW-VN READING-VNS OFF-LIVE-READING-VNS)
READING-VNS and OFF-LIVE-READING-VNS should be subsets of the
corresponding readers of VN. NEW-VN is spliced so that it is a
reader of VN, with READING-VNS now reading NEW-VN instead of VN.
All the VNs in OFF-LIVE-READING-VNS are made to off-live read NEW-VN
instead of VN.

(VN:DELETE VN)
Deletes a VN from the DAG and VN universe. A VN can be deleted
only if it has no readers and no off-live readers.

(VN:PRINT VN)
Prints VN in a pretty debugging format.

$#VN I
Interactive syntax for accessing VN with number I.
=====

(visible-fields vn number type name oper)

(eval-when (compile load)
  (include list-scheduler:declarations) )

(def-object-universe *ls.vn-universe*
  (object-name      vn)
  (mapping-type     numbered-objects)
  (object-number-function
    (lambda (object)
      (vn:number object) ) )
  (set-object-number-function
    (lambda (object number)
      (:= (vn:number object) number) ) )
  (initial-size 200) )

(def-sharp-sharp vn
  '(vn-universe:number:vn ,(read) ) )

(defun vn.initialize ()
  (vn-universe:initialize)
  () )

(defun vn:create ( vn )
  (if (! vn) (then
    (:= vn (vn:new) ) )
    (else
      (assert (vn:is vn) ) ) )
  (vn-universe:add vn)
  vn)

(defun vn:operator ( vn )
  (assert (vn:is vn) )
  (oper:operator (vn:oper vn) ) )

```

```

(defun vn:group ( vn )
  (assert (vn:is vn) )
  (oper:group (vn:oper vn) ) )

(defun vn:delay ( vn )
  (assert (vn:is vn) )
  (caseq (vn:type vn)
    (operation
      (assert (vn:me vn) )
      (me:delay (vn:me vn) ) )
    (t
      () ) ) )

(defun vn:constrained-vn:delay ( vn constrained-vn )
  (assert (vn:is vn) )
  (loop (for constraining-vn in (vn:constraining-vns constrained-vn) )
    (for constraining-delay in (vn:constraining-delays constrained-vn) )
    (when (== constraining-vn vn) )
    (maximize constraining-delay) ) )

(defun vn:insert-vns ( vn predecessor-vns )
  (assert (== 1 (length (vn:operand-vns vn) ) ) )
  (let ( (operand-vn (car (vn:operand-vns vn) ) ) )
    (:= (vn:reading-vns operand-vn) (top-level-removeq vn &&&) )

    (loop (for predecessor-vn in predecessor-vns) (do
      (push (vn:reading-vns operand-vn) predecessor-vn)
      (:= (vn:operand-vns predecessor-vn) '(,operand-vn) )
      (:= (vn:reading-vns predecessor-vn) '(,vn) ) ) )

    (:= (vn:operand-vns vn) predecessor-vns)
    () ) )

(defun vn:splice-vn ( vn new-vn reading-vns off-live-reading-vns )
  (assert (subset? reading-vns (vn:reading-vns vn) ) )
  (assert (subset? off-live-reading-vns (vn:off-live-reading-vns vn) ) )
  (assert (|| reading-vns off-live-reading-vns) )

  (:= (vn:reading-vns vn) (set-diffq &&& reading-vns) )
  (push (vn:reading-vns vn) new-vn)

  (:= (vn:operand-vns new-vn) '(,vn) )
  (:= (vn:reading-vns new-vn) reading-vns)

  (loop (for reading-vn in reading-vns) (do
    (:= (vn:operand-vns reading-vn)
      (top-level-substq new-vn vn &&&) ) ) )

  (:= (vn:off-live-reading-vns vn) (set-diffq &&& off-live-reading-vns) )
  (:= (vn:off-live-reading-vns new-vn) off-live-reading-vns)

  (loop (for off-live-reading-vn in off-live-reading-vns) (do
    (:= (vn:off-live-vns off-live-reading-vn)
      (top-level-substq new-vn vn &&&) ) ) )
  () )

```

```

(defun vn:delete ( vn )
  (assert (&& (! (vn:reading-vns      vn) )
            (! (vn:off-live-reading-vns vn) )
            (! (vn:off-live-vns      vn) ) ) )

  ;** Remove references to this VN from any VNs constrained by this
  ;** one: ugh, messy.
  ;**
  (loop (for constrained-vn in (vn:constrained-vns vn) ) (do
    (loop (for constraining-vn in (vn:constraining-vns constrained-vn)
      (for delay in (vn:constraining-delays constrained-vn)
        (when (!= constraining-vn vn) )
        (initial new-constraining-vns ()
          new-constraining-delays () )
        (do
          (push new-constraining-vns constraining-vn)
          (push new-constraining-delays delay) )
        (result
          (:= (vn:constraining-vns constrained-vn)
              (dreverse new-constraining-vns) )
          (:= (vn:constraining-delays constrained-vn)
              (dreverse new-constraining-delays) ) ) ) ) )

  ;** Now removes references to this VN from constraining VNs.
  ;**
  (loop (for constraining-vn in (vn:constraining-vns vn) ) (do
    (:= (vn:constrained-vns constraining-vn)
        (top-level-removeq vn &&&) ) ) )

  ;** Remove all references to VN from its operands.
  ;**
  (loop (for operand-vn in (vn:operand-vns vn) ) (do
    (:= (vn:reading-vns operand-vn)
        (top-level-removeq vn &&&) ) ) )

  (vn-universe:delete vn)
  () )

(defun vn:print ( vn )
  (assert (vn:is vn) )

  (msg 0 (vn:number vn) " : " )
  (caseq (vn:type vn)
    (def (msg "def " (vn:name vn) " " ) )
    (def1 (msg "def1 " (vn:name vn) " " ) )
    (use (msg "use " (vn:name vn) " " ) )
    (use1 (msg "use1 " (vn:name vn) " " ) )
    (operation (msg (vn:oper vn) " " ) )
    (pseudo-op (msg (vn:oper vn) " " ) )
    (copy (msg "copy " (vn:oper vn) " " ) )
    (t (error (list vn "VN:PRINT: Unknown VN:TYPE.") ) ) )
  (if (vn:constant? vn) (then
    (msg "c " ) ) )
  (if (vn:datatype vn) (then
    (msg (vn:datatype vn) " " ) ) )
  (msg t)

  (if (vn:operand-vns vn) (then
    (msg " operand-vns: " )
    (loop (for operand-vn in (vn:operand-vns vn) ) (do
      (msg (vn:number operand-vn) " " ) ) ) ) )

```

```

(msg t) ) )

(if (vn:constraining-vns vn) (then
  (msg " constraining-vns: " )
  (loop (for constraining-vn in (vn:constraining-vns vn) )
    (for constraining-delay in (vn:constraining-delays vn) )
    (do
      (msg (vn:number constraining-vn) "@" constraining-delay " " ) ) ) )
  (msg t) ) )

(if (vn:off-live-vns vn) (then
  (msg " off-live-vns: " )
  (loop (for off-live-vn in (vn:off-live-vns vn) ) (do
    (msg (vn:number off-live-vn) " " ) ) )
  (msg t) ) )

(if (vn:constrained-vns vn) (then
  (msg " constrained-vns: " )
  (loop (for constrained-vn in (vn:constrained-vns vn) ) (do
    (msg (vn:number constrained-vn) " " ) ) )
  (msg t) ) )

(if (vn:depth vn) (then
  (msg " depth: " (vn:depth vn) t) ) )

(if (vn:cycle vn) (then
  (msg " cycle: " (vn:cycle vn) t) ) )

(if (vn:bug-cycle vn) (then
  (msg " bug-cycle: " (vn:bug-cycle vn) t) ) )

(if (vn:scheduled-cycle vn) (then
  (msg " scheduled-cycle: " (vn:scheduled-cycle vn) t) ) )

(if (vn:me vn) (then
  (msg " me: "
    (me:name (vn:me vn) ) t) ) )

(if (vn:likely-mes vn) (then
  (msg " likely-mes: " (h (me-list:name-list (vn:likely-mes vn)))
    t) ) )

(if (vn:locations vn) (then
  (msg " locations: "
    (h (loop (for (me register) in (vn:locations vn) ) (save
      '(. (me:name me) ,register) ) ) )
    t) ) )

(if (vn:register-bank-me vn) (then
  (msg " register: "
    (me:name (vn:register-bank-me vn) ) " " (vn:register vn)
    t) ) )
  () )

```

```

*****
*** Copyright (C) 1983 John R. Ellis. This file may be used for
*** educational and research purposes only, by the faculty, students,
*** and staff of Yale University only.
*****

```

## WRITE-VNS

This module provides the interface functions for printing out VN DAGs. See the similar modules in DRAWING: for details about how to get DAGs printed out.

### (WRITE-VNS &OPTIONAL (FILENAME "IN.DAG") DAG-COMMANDS)

Writes a description of the current VN DAG to FILENAME. DAG-COMMANDS is an optional list of assignment statements that will be recorded with the DAG and EVALed by the DAG-drawing software; these assignments control the formatting of the DAG. See the variable \*VV.DEFAULT-DAG-COMMANDS\* below.

```

(eval-when (compile eval)
  (build '(drawing:dag-input-node) ) )

```

```

(defvar *vv.default-dag-commands*
  ( ( := *dag:box-height*           48)
    ( := *dag:node-width*          100)
    ( := *dag:line-width*          25)
    ( := *dag:box-width*           80)
    ( := *dag:minimum-level-separation* 30)
    ( := *dag:ideal-text-lines-per-node* 4)
    ( := *dag:fold-text-lines?*    () )
    ( := *dag:critical-threshold*  0)
    ( := *dag:remove-non-critical-nodes&edges?*
      ( ) )
    ( := *dag:shade-critical-nodes?*
      ( ) )
    ( := *dag:text-x-margin*       8)
    ( := *dag:make-slides?*       ( ) ) ) )

```

```

(defun write-vns ( &optional (filename "in.dag") dag-commands)
  (let ( (filename (file-list:filename
                    (file-list:merge filename '(() () () "DAG") ) ) ) ) )
    (iota ( (file filename "out newversion") ) (without file
      (msg (h '(. *vv.default-dag-commands* ..dag-commands)
              100000 100000) t)
      (msg (h (vv.dag:graph) 100000 100000) t)
      (filename file) ) ) ) )

```

```

(defun vv.dag:graph ()
  (let ( (*vv.all-nodes* () ) )
    (loop (for-each-vn vn)
      (when (vv.vn:visible? vn)
        (do
          (push *vv.all-nodes*
            (caseq (vn:type vn)
              ( (def use)
                (vv.def-use-vn:create-node vn) )
              ( (defl use1)
                (vv.defl-use1-vn:create-node vn) )
              (operation

```

```

          (vv.operation-vn:create-node vn) )
            (copy
              (vv.copy-vn:create-node vn) ) ) ) ) )
    (result
      (dreverse *vv.all-nodes*) ) ) ) )

```

```

(defun vv.def-use-vn:create-node ( vn )
  (dag:input-node:new
    name (vn:number vn)
    label '(.(string-msg (vn:number vn) " " (vn:name vn) )
            ,(me-list:name-list (vn:likely-mes vn) ) )
    style 'shading2
    child-edge-styles
      (vv.vn:child-edge-styles vn) ) )

```

```

(defun vv.defl-use1-vn:create-node ( vn )
  (dag:input-node:new
    name (vn:number vn)
    label '(.(string-msg (vn:number vn) " " (vn:name vn) )
            ,(vv.vn:location vn) )
    style 'shading2
    child-edge-styles
      (vv.vn:child-edge-styles vn) ) )

```

```

(defun vv.operation-vn:create-node ( vn )
  (dag:input-node:new
    name (vn:number vn)
    label '(.(string-msg (vn:number vn) " " (me:name (vn:me vn) ) )
            ,(vn:oper vn)
            ,(vv.vn:location vn) )
    child-edge-styles
      (vv.create-chain
        (me:delay (vn:me vn) )
        'node
        '(.(me:name (vn:me vn) ) )
        ( )
        'normal
        (vv.vn:child-edge-styles vn) ) ) ) )

```

```

(defun vv.copy-vn:create-node ( vn )
  (dag:input-node:new
    name (vn:number vn)
    label '(.(string-msg (vn:number vn) " " (vv.vn:me-name vn) )
            copy
            ,(vv.vn:location vn) )
    style 'shaded
    child-edge-styles
      (vv.vn:child-edge-styles vn) ) )

```

```

(defun vv.vn:me-name ( vn )
  (?( (vn:register-bank-me vn)
    (me:name (vn:register-bank-me vn) ) )
    ( (vn:me vn)
      (me:name (vn:me vn) ) )
    (t
      ==) ) )

```

```

(defun wv.vn:location ( vn )
  (? ( (vn:register-bank-me vn)
      (string-msg (me:name (vn:register-bank-me vn)) " "
                  (vn:register vn) ) )
    ( (vn:me vn)
      (me:name (vn:me vn) ) )
    (t
     "" ) ) )

(defun wv.vn:child-edge-styles ( vn )
  (nconc
   (loop (for reading-vn in (nodupseq (vn:reading-vns vn) ) )
         (when (wv.vn:visible? reading-vn) )
         (save
          (wv.vn:child-vn:style:create-edge vn reading-vn 'thick) ) )
   (loop (for constrained-vn in (nodupseq (vn:constrained-vns vn) ) )
         (when (&& (> (vn:constrained-vn:delay vn constrained-vn) 0)
                (wv.vn:visible? constrained-vn) ) )
         (save
          (wv.vn:child-vn:style:create-edge vn constrained-vn 'shaded))))))

(defun wv.vn:child-vn:style:create-edge ( vn child-vn style )
  (let ( (next-cycle (wv.vn:next-cycle vn) )
        (child-cycle (wv.vn:cycle child-vn) ) )
    (if (== next-cycle child-cycle) (then
      '(.(vn:number child-vn) ,style) )
      (else
       (let ( (closest-pred-vn (wv.vn:closest-pred-vn child-vn) )
             (if (!= vn closest-pred-vn) (then
              '(.(vn:number child-vn) ,style) )
              (else
               (car (wv.create-chain
                    (- child-cycle next-cycle)
                    'line
                    ()
                    style
                    style
                    '(.(vn:number child-vn) ,style))))))))))

(defun wv.vn:closest-pred-vn ( vn )
  (let ( (closest-operand-vn
        (loop (for operand-vn in (vn:operand-vns vn) )
              (maximize operand-vn (wv.vn:next-cycle operand-vn) ) ) )
        (closest-constraining-vn
        (loop (for constraining-vn in (vn:constraining-vns vn) )
              (for constraining-delay in (vn:constraining-delays vn) )
              (when (> constraining-delay 0) )
              (maximize constraining-vn
                        (wv.vn:next-cycle constraining-vn) ) ) ) )
    (if (&& closest-constraining-vn
          (> (wv.vn:next-cycle closest-constraining-vn)
              (wv.vn:next-cycle closest-operand-vn) ) )
        (closest-constraining-vn
         closest-operand-vn) )
    (closest-operand-vn) ) )

```

```

(defun wv.create-chain
  ( length type label box-style edge-style child-edge-styles )
  (loop (initial prev-child-edge-styles child-edge-styles)
        (decr i from length to 1)
        (bind node (dag:input-node:new
                    name (gensym)
                    type type
                    label label
                    style box-style
                    child-edge-styles
                    prev-child-edge-styles) )
        (do
         (push *wv.all-nodes* node)
         (:= prev-child-edge-styles '(.(dag:input-node:name node)
                                     ,edge-style) ) ) )
        (result prev-child-edge-styles) ) )

(defun wv.vn:next-cycle ( vn )
  (+ 1 (+ (wv.vn:cycle vn)
          (caseq (vn:type vn)
                 (operation (me:delay (vn:me vn) ) )
                 (t 0) ) ) ) )

(defun wv.vn:cycle ( vn )
  (caseq (vn:type vn)
         (def
          (if (== 'def1 (vn:type (car (vn:reading-vns vn) ) ) )
              -2
              -1) )
         (def1
          -1)
         ( (copy operation)
           (vn:cycle vn) )
         (use1
          (vn:cycle vn) )
         (use
          (if (== 'use1 (vn:type (car (vn:operand-vns vn) ) ) )
              (+ 1 (vn:cycle vn) )
              (vn:cycle vn) ) )
         (t
          (error (list vn "Case error.") ) ) ) )

(defun wv.vn:visible? ( vn )
  (caseq (vn:type vn)
         (pseudo-op
          () )
         (def
          (for-some (reading-vn in (vn:reading-vns vn) )
                    (&& (!= 'use (vn:type reading-vn) )
                        (wv.vn:visible? reading-vn) ) ) )
         (def1
          (for-some (reading-vn in (vn:reading-vns vn) )
                    (!= 'use1 (vn:type reading-vn) ) )
          ( operation copy
            t
            ( use1 use
              (wv.vn:visible? (car (vn:operand-vns vn) ) ) ) )
          (t
           (error (list vn "Case error.") ) ) ) )

```



