

Examples of LITTLE-generated Code

In an earlier newsletter (LITTLE news No. 20) we gave some examples of the LITTLE source code from the SETL Run-Time Library (SRTL); in this newsletter, we give some examples of the 6600 machine code produced by the current LITTLE compiler for sample parts of the SRTL. For each example, we present the LITTLE source (after macro expansion) and a COMPASS-like representation of the machine-code generated by LITTLE. Note that the use of a '+' character in the label field indicates that the instruction begins a new word.

Ex. 1--Simple Stores

The LITTLE source

T = 5000 - MAXZZYZ; TRES = T; RUNNINGBLK = T;
compiler into

	SX1	5000
	SA2	MAXZZYZ
+	IX3	X1 - X2
	BX6	X3
	SA6	T
+	SA6	TRES
	SA4	T
+	BX7	X4
	SA7	RUNNINGBLK

This code takes about 60 minor cycles (a minor cycle is 100 nanoseconds) to execute, and requires four words. The preferred code, which takes about 35 minor cycles, and three words, is

+	SA2	MAXZZYZ
	SX1	5000
+	IX6	X1 -X2
	BX7	X6
	SA6	T
+	SA7	TRES
	SA6	RUNNINGBLK

Ex 2-- Argument access and branching

The LITTLE source

```

SUBR      GETSTG (N,P);
IF ( N.EQ.0 ) GO TO A09;
GO TO B09;

```

/A09/ CALL ABORT; /B09 / ...

compilers into

CALLAD	BSS	1	.arg-list address planted .here as part of call
GETSTG	BSS	1	.entry/exit word
+	SA1	CALLAD	
	SA2	X1 + 0	
+	SA2	X2	.X2 = value of N
	MX3	0	.X3 = 0
	BX4	X2-X3	.equality test
	SB2	B0	
+	NZ	X4, L01	
	SB2	1	
L01	NE	B2, B0, L02	
+	JP	B09	
L02	NO		.full-word no-ops
+	SA0	ABORT-1	
	RJ	ABORT	
B09	...		

This code takes seven words, and executes in about 80 cycles (if ABORT is called). The preferred code, requiring three words and about 35 cycles (if ABORT called) is

GETSTG	BSS	1
	SA1	CALLAD
	SA2	X1 + B0
	SA2	X2 + B0
+	NZ	X2, B09
	RJ	ABORT

Ex. 3-- multiplication by 1

Though the LITTLE source is not available,
the following code-fragment (from routine START) was observed:

```

      SA1          T
+     SX0          1
      PX0          X0, B0      .Pack
      PX1          X1, B0
+     DX1          X1 * X0      . multiply
      UX1          X1, B0      .get integer product
      SA2          X1 + U      .ie, U(x1)
+     MX3          0
      BX6          X3
      SA6          A2
+     SA1          T
      SA2          1
+     IX3          X1 + X2
      BX6          X3
      SA6          A1

```

The source is probably

$$U(T*1) = 0; T = T + 1 ;$$

The preferred code is

```

+     SA1          T
      MX6          0
+     SA6          X1 + U
      SX7          X1 + 1
+     SA7          T

```

The longer code takes about 60 minor cycles, the shorter about 30.

Ex. 4-- simple loop

The LITTLE source

```

        TEMP = 1;
/A01/ IF (TEMP .GT. 5000)GO TO B01;
        STORAGE (TEMP) = 0;
        TEMP = TEMP + 1;
        GO TO A01;

```

which is a "memory-set" loop, compiles into

```

        SX1          1
        BX6          X1
+       SA6          TEMP
A01     SA1          TEMP
        SX2          5000
+       IX3          X2 - X1
        SB1          B0
+       PL           X3, L01
        SB1          1
L01     NE           B1, B0, B01
        SA1          TEMP
+       SB2          X1
        SA2          B2 + STORAGE
        MX3          0
+       BX6          X3
        SA6          A2
+       SA1          TEMP
        SX2          1
+       IX3          X1 + X2
        BX6          X3
        SA6          A1
+       JP           A01

```

Since loop takes more than 7 words, it doesn't fit in the stack. Since TEMP not used outside loop, it need not be stored, also TEMP used as a subscript and thus may be kept in B-register.

Thus preferred code is

	SB1	1
	SB2	5000
	SB3	1
	MX6	0
	SA6	STORAGE + B1
L	SB1	B1 + B3
	SA6	A6 + B1
	LT	B1, B2, L
+	...	

The shorter code fits in stack (main loop is a single word), and requires about three words for entire loop instead of nine for longer LITTLE generated code.

Note the perhaps the best way to handle storage-set loops is to call a storage-set function which is carefully hand coded to fit in stack and stores both $STORAGE (TEMP)$ and $STORAGE(TEMP+1)$ in single pass thru loop (such a routine is available) similar remarks apply to storage-move loops.

In summary, the example show that a relatively simple improvement in code-generation would probably reduce both code-size and execution time by a factor of two for typical LITTLE programs. Furthermore, most of these improvements could be done in a separate job step, which takes as input the code produced by LITTLE and produces improved code (on a subroutine by subroutine basis). This separate code-improver could be developed without any substantial change in the current LITTLE compiler; the only change required in the LITTLE compiler is the ability to produce symbolic, COMPASS-like output, instead of loader input modules (otherwise, the code-analyser must unpack the loader tables).