

APL Version of Peter Markstein's McKeeman Table Routine

NOTES: INPUT IS A VECTOR, G, WHICH REPRESENTS A GRAMMAR, ENCODED AS A CHARACTER STRING. A PRODUCTION, A←BCD WOULD BE ENCODED ABCD| FOR EXAMPLE. THE FIRST LINE COMPUTES C, THE VECTOR OF UNIQUE CHARACTERS OCCURRING IN THE GRAMMAR. S BECOMES THE VECTOR OF LHS OF PRODUCTIONS, B THE FIRST ELEMENTS ON THE RHS OF PRODUCTIONS, E THE LAST ELEMENT OF THE PRODUCTIONS. BEGINS[I;J] IS SET TO 1 IF THE JTH SYNTACTIC TYPE CAN BEGIN WITH THE ITH SYNTACTIC TYPE; LIKEWISE FOR ENDS[I;J]. THE MATRIX T, ALSO EXTERNAL TO PRECEDENCE, HOLDS THE RESULT:

$$T[I;J] = \begin{cases} 1 & \text{IF } C[I] = C[J] \\ 2 & \text{IF } C[I] > C[J] \\ 3 & \text{IF } C[I] < C[J] \\ 5 & \text{IF } C[I], C[J] \text{ IS AMBIGUOUS} \\ 0 & \text{IF } C[I], C[J] \text{ IS UNGRAMMATICAL.} \end{cases}$$

∇ PRECEDENCE

$X \leftarrow \rho C \leftarrow (\rightarrow / Q \circ \cdot < Q + \rho Z) \wedge Z \circ \cdot = Z / Z + (G \neq \rho | \rho) / G$

$T \leftarrow \text{BEGINS} \leftarrow \text{ENDS} \leftarrow (X, X) \rho 0$

$S \leftarrow C \setminus G [1 + Q + 0, K [i \bar{1} + \rho K + (G \neq \rho | \rho) / \rho G]]$

$B \leftarrow C \setminus G [2 + Q]$

$E \leftarrow C \setminus G [K - 1]$

$\text{BEGINS}[B; S] \leftarrow I + J + 1$

$\text{ENDS}[E; S] \leftarrow 1$

$\text{BEGINS} \leftarrow \text{COMPLETE BEGINS}$

$\text{ENDS} \leftarrow \text{COMPLETE ENDS}$

$\underline{XM}: T[I; J] \leftarrow C[I] \text{ NEXTTO } C[J]$

$Z \leftarrow (C[I] \text{ SMALL } C[J]) + C[I] \text{ BIGGER } C[J]$

$T[I; J] \leftarrow ((\sim K) \times T[I; J] + Z) + 5 \times K + 0 \neq T[I; J] \times Z$

$\rightarrow ((\rho C) \geq J + J + 1) / \underline{XM}$

$\rightarrow ((\rho C) \geq I + I + J + 1) / \underline{XM} \nabla$

∇ Z ← COMPLETE X; Y; Q

$Q \leftarrow + / + / Z \leftarrow X$

$\underline{LCO}: Y \leftarrow Q$

$Z \leftarrow Z \vee \cdot \wedge Z$

$\rightarrow (Y \neq Q \leftarrow + / + / Z) / \underline{LCO} \nabla$

∇ Z ← A NEXTTO B

$Z \leftarrow \vee / (0, 0, (G \neq \rho | \rho)) \wedge (0, (G = A), 0) \wedge (G = B), 0, 0 \nabla$

∇ Z ← A BIGGER B; U

$U \leftarrow (\text{BEGINS}[C \setminus B;] / C), (\sim \text{BEGINS}[B; B]) \rho B$

$Z \leftarrow 2 \times \vee / (0, 0, (G \neq \rho | \rho)) \wedge (0, (\vee / G \circ \cdot = \text{ENDS}[C \setminus A;] / C), 0) \wedge ((\vee / G \circ \cdot = U), 0, 0) \nabla$

∇ Z ← A SMALL B

$Z \leftarrow 3 \times \vee / (0, 0, (G \neq \rho | \rho)) \wedge (0, (G = A), 0) \wedge ((\vee / G \circ \cdot = \text{BEGINS}[C \setminus B;] / C), 0, 0) \nabla$